## Pathfinding with Trees

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## Pathfinding



# We are given <br> > a map 

## Pathfinding



# We are given <br> > a map <br> > a start point 

## Pathfinding



# We are given <br> > a map <br> > a start point <br> > a goal point 

## Dijkstra's Algorithm



A simple solution:
Look at all points neighbouring the start point

## Dijkstra's Algorithm



A simple solution:
Look at all points neighbouring the start point Continue until the goal is found

## Dijkstra's Algorithm



A simple solution:
Look at all points neighbouring the start point Continue until the goal is found
Follow the arrows back to the start

## Dijkstra's Algorithm -> Tree Cache



Save the search tree for a single root

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Save the search tree for a single root
When searching:

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Save the search tree for a single root
When searching:
> traverse the tree from start to root

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When searching:
> traverse the tree from start to root
> traverse the tree from goal to root

## Dijkstra's Algorithm -> Tree Cache



Save the search tree for a single root
When searching:
> traverse the tree from start to root
traverse the tree from goal to root
> invert the second part
> concatenate parts at the root

## Tree Cache: Example 2



Paths can have redundant parts

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Paths can have redundant parts
The redundant part can be easily removed by looking for the lowest common ancestor (LCA) of the two nodes

## LCA: Naive implementation



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## LCA: Faster implementation



## LCA: Faster implementation



A depth-first traversal and corresponding levels:

| A | B | A | C | D | F | D | G | D | H | D | C | E | C | A |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 0 | 1 | 2 | 3 | 2 | 3 | 2 | 3 | 2 | 1 | 2 | 1 | 0 |

## LCA: Faster implementation



A depth-first traversal and corresponding levels:

| A | B | A | C | D | F | D | G | D | H | D | C | E | C | A |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 0 | 1 | 2 | 3 | 2 | 3 | 2 | 3 | 2 | 1 | 2 | 1 | 0 |

## Range minimum query implementation

| A | B | A | C | D | F | D | G | D | H | D | C | E | C | A |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 0 | 1 | 2 | 3 | 2 | 3 | 2 | 3 | 2 | 1 | 2 | 1 | 0 |

## Range minimum query implementation

| A | B | A | C | D | F | D | G | D | H | D | C | E | C | A |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 0 | 1 | 2 | 3 | 2 | 3 | 2 | 3 | 2 | 1 | 2 | 1 | 0 |

The range minimum query can be solved with a few lookups by pre-calculating the minimum for all sub-ranges of length $2^{i}$ :

| A | A | A | C | D | D | D | D | D | D | C | C | C | A | 2 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A | A | A | C | D | D | D | D | C | C | C | A |  |  | 4 |
| A | A | A | C | C | C | C | A |  |  |  |  |  |  | 8 |

## Range minimum query implementation

| A | B | A | C | D | F | D | G | D | H | D | C | E | C | A |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 0 | 1 | 2 | 3 | 2 | 3 | 2 | 3 | 2 | 1 | 2 | 1 | 0 |

The range minimum query can be solved with a few lookups by pre-calculating the minimum for all sub-ranges of length $2^{i}$ :

| A | A | A | C | D | D | D | D | D | D | C | C | C | A | 2 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A | A | A | C | D | D | D | D | C | C | C | A |  |  | 4 |
| A | A | A | C | C | C | C | A |  |  |  |  |  |  | 8 |

## Tree Cache: Example 3



## Paths can be bad without redundant parts

## Tree Cache improvement: More than one tree



Generate more than one tree, store distance as well as parent

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Generate more than one tree, store distance as well as parent
When searching:
> calculate path distance for each tree
> choose tree with shortest path
construct path

## Tree Cache: Example 4



Tree Cache generated path far too long without redundancy

## Using Multiple Trees



## Using multiple trees generates better path with redundancy

## Both improvements combined



Generate more than one tree, store distance as well as parent, calculate LCA information for each tree

## Both improvements combined



Generate more than one tree, store distance as well as parent, calculate LCA
information for each tree When searching:
> look up LCA and calculate path distance for each tree
> choose tree with shortest path
construct path

## Experiments



Maps from the Gridbased Path Planning Competition Most maps $512 \times 512$ tiles, the largest 768x768

## Results



## Results



| map size:768x768 | time per tree | memory per tree |
| :---: | :---: | :---: |
| Trees | 230 ms | 10 MB |
| Trees + LCA | 400 ms | 110 MB |

## Results: Root placement




## Conclusions

## Tree Cache finds potentially bad paths very fast

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Tree Cache finds potentially bad paths very fast
Path quality can be improved significantly
Combination of improvements big amounts of memory
Good results using only a few trees

## Future Work

## advanced root placement strategies

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# advanced root placement strategies <br> compress tree information 

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# advanced root placement strategies <br> compress tree information <br> adapt to directed graphs 

