

Value Compression of Pattern Databases

Nathan R. Sturtevant, *University of Denver*; Ariel Felner, *BGU*; Malte Helmert, *University of Basel*

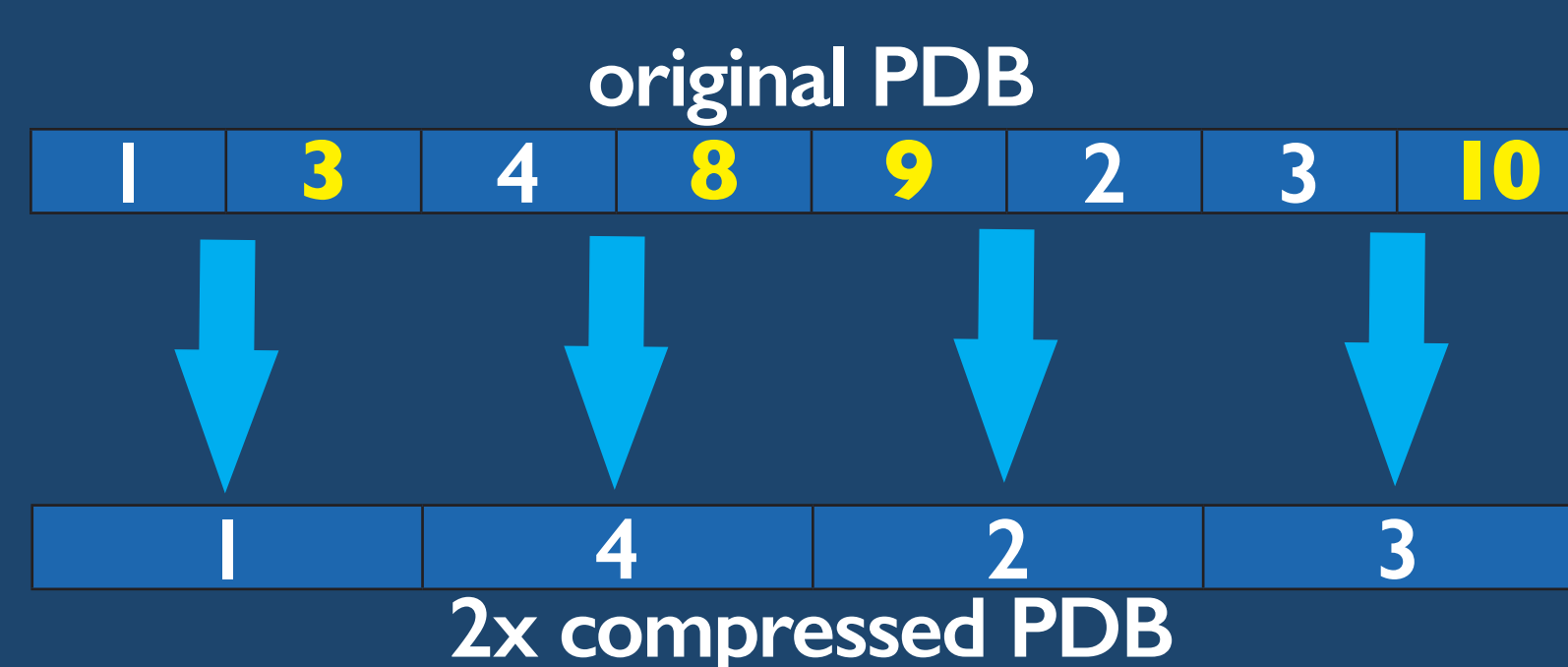
Pattern Database Compression

What is a Pattern Database (PDB)?

A **PDB** is a heuristic that estimates the distance to the goal for search algorithms such as A*. It computes and stores the distances in an abstract state space and stores them in a table.

Entry Compression (Felner et al, 2007) compresses the PDB by **combining** entries and storing the minimum. Thus, there are fewer entries in the compressed PDB.

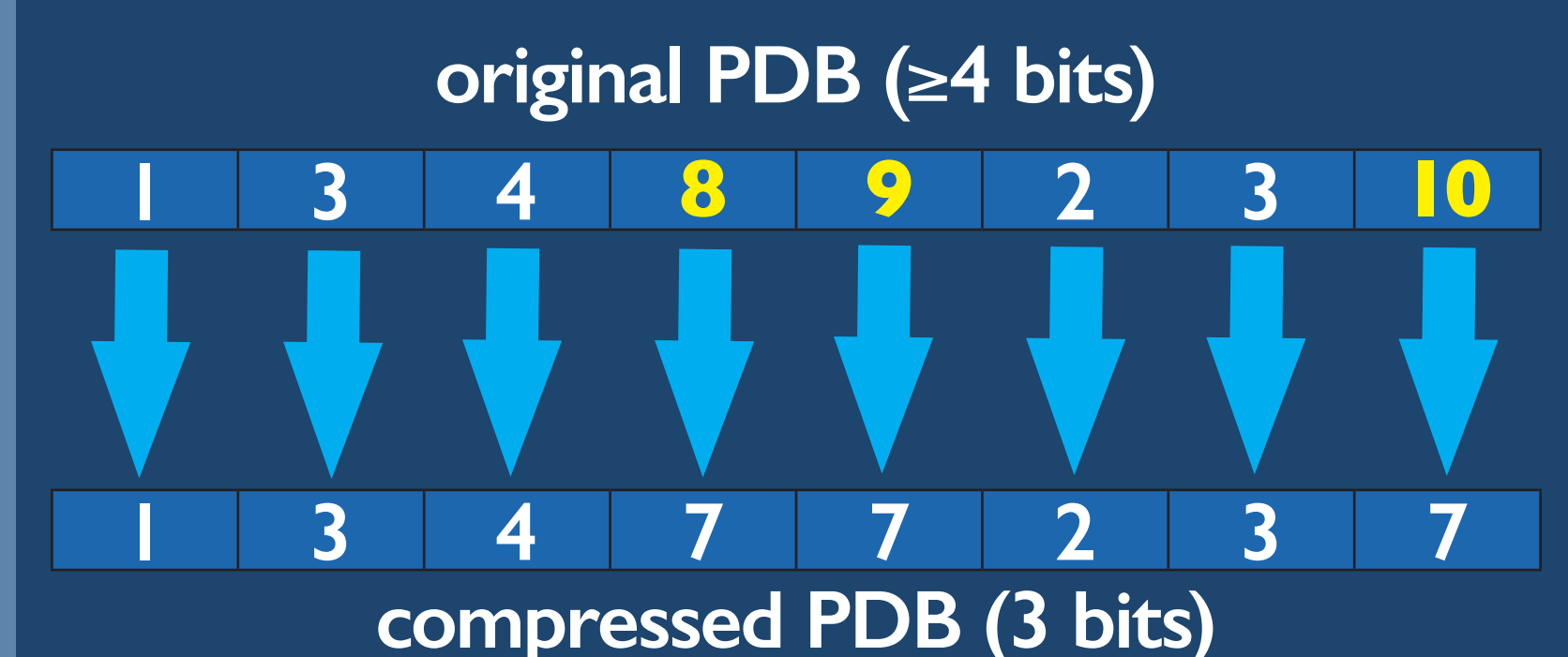
Entry Compression (EC)



The minimum value is stored so that the heuristic is still admissible (non-overestimating).

Value Compression (VC) [new]

VC keeps all entries in the PDB, but reduces the number of bits per entry. This reduces the ranges of values that can be stored.



Value Compression: No wasted bits

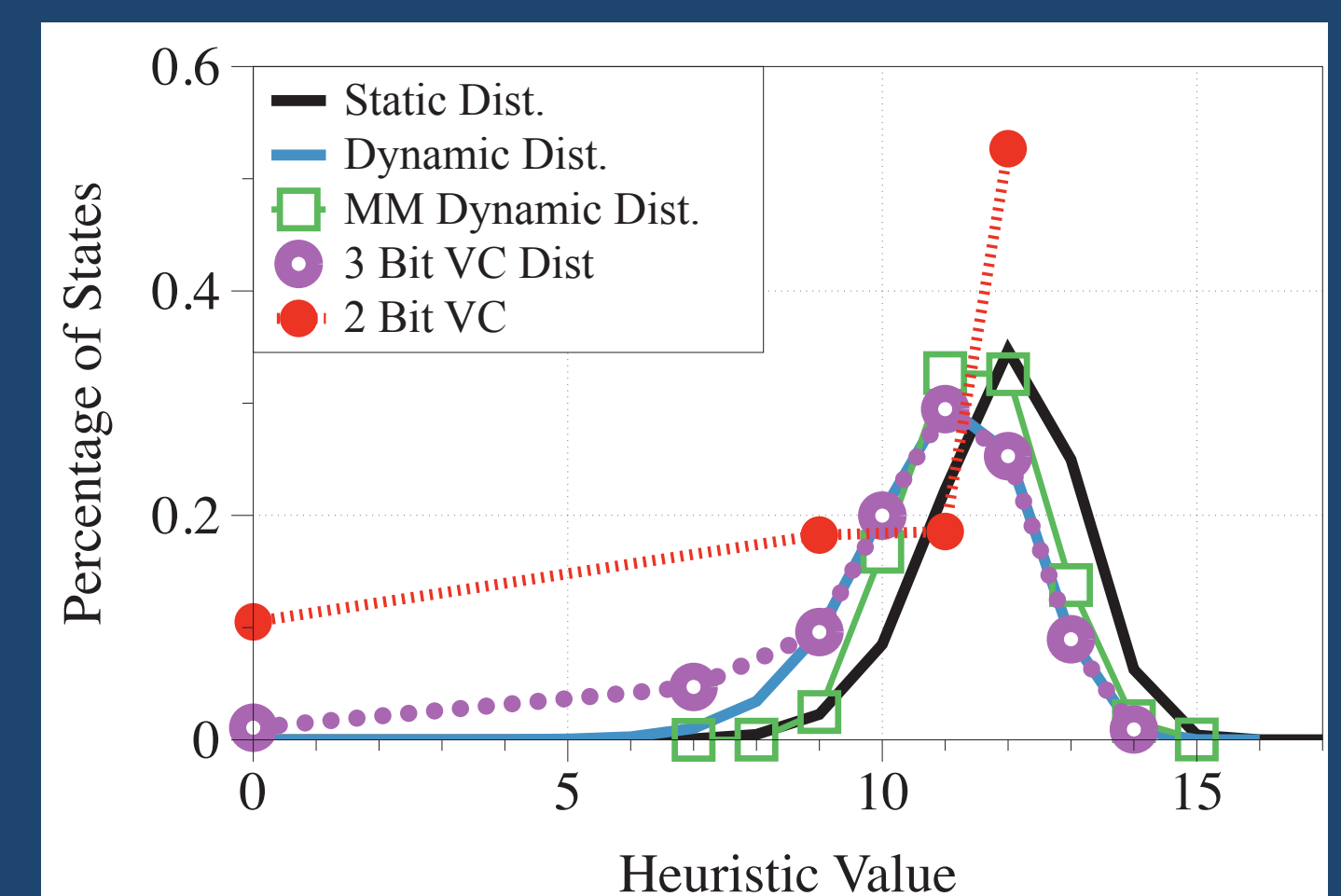
Value Compression Details

- Value Compression (VC) divides the heuristic values into ranges $R_1 \dots R_n$. Ranges are stored instead of values.
- The minimum value of the range is used during search to maintain admissibility.
- Dynamic programming is used to optimize the ranges and maximize the average h-value in the compressed PDB.

Example of EC vs VC in Top Spin

	1.760MB		880MB		440MB
D	Total	VC ₂	VC ₂	EC ₂	VC ₄
0	1	1	12	2	10,188,753
1	11	11	22	40	4
2	94	94	186	340	
3	731	731	1,430	2,596	
4	5,353	5,353	10,340	18,736	
5	37,275	37,275	70,894	127,756	
6	245,468	245,468	457,304	813,700	
7	1,508,099	1,508,099	2,722,458	4,724,408	
8	8,391,721	8,391,721	14,408,820	23,870,392	
9	40,012,497	40,012,497	63,502,746	97,318,252	
10	150,000,765	150,000,765	212,692,340	290,434,356	
11	393,482,172	393,482,172	478,114,034	553,276,900	
12	612,084,904	612,084,904	601,419,722	549,750,508	
13	440,655,534	440,655,534	328,304,534	217,340,348	
14	110,437,757	110,437,757	59,883,892	26,009,144	
15	7,389,524	7,460,178	2,721,910	634,464	
16	70,633	70,654	11,924	616	
17	21		2		
Avg.	11.90	11.90	11.90	11.59	11.38

Runtime Heuristic Distribution



Experimental Results

General Observations

- Conventional wisdom says the low values in the PDB are more important than the high values. So, removing low values from the PDB should hurt performance.
- In practice, the most common heuristic values in the runtime distribution (looked up during search) must be preserved.
- Bidirectional Pathmax (BPMX) is crucial for local propagation of heuristic values and recovery of lost information.
- VC is most effective when the number of values in the PDB is just larger than the nearest power of two.
- EC can be effectively combined with VC.**

Experimental Setup

The paper contains experiments on many domains and algorithms. We report on 18-4 Top Spin here.

We compare compression factors in IDA* with BPMX as well as combinations of EC and VC. Combinations with VC have the best performance (bold).

Memory	EC	VC	VC-bits	Nodes	Time
1	1	1	8	3.88M	15.29
0.5	1	2	4	3.88M	15.32
0.375	1	2.66	3	4.03M	15.44
0.25	1	4	2	10.39M	33.63
0.5	2	1	8	7.11M	27.70
0.25	2	2	4	7.11M	27.88
0.1875	2	2.66	3	7.37M	28.44
0.125	2	4	2	30.43M	80.04
0.25	4	1	8	13.75M	51.06
0.125	4	2	4	13.74M	50.97
0.094	4	2.66	3	14.31M	51.52
0.0625	4	4	2	30.48M	77.68

Bidirectional Search: MM

- The MM algorithm guarantees that the search frontiers meet in the middle.
- Small heuristic values aren't used.**

Memory	EC	VC	VC-bits	Nodes	Time
1	1	1	8	2.07M	30.60
0.5	2	1	8	3.55M	55.61
0.5	1	2	4	2.07M	30.93
0.25	4	1	8	5.42M	83.87
0.25	2	2	4	3.55M	55.55
0.25	1	4	2	4.63M	66.46
0.125	4	2	4	5.42M	84.06
0.125	2	4	2	5.19M	79.17

