# Additive Pattern Databases for Decoupled Search 

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## Setting \& Motivation

- optimal classical planning as heuristic search
- state of the art: abstraction heuristics
- successful alternative to explicit search: decoupled search


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- successful alternative to explicit search: decoupled search
- goal: abstraction heuristics for decoupled search


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- pattern database (PDB) heuristics:
- project variables to a subset
- store perfect heuristic values of abstraction


## Decoupled Search in a Nutshell

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- branch over center states and actions, handle leaves separately
- decoupled state $s^{\mathcal{F}}$ :
- center state
- pricing function: cost of reachable leaf states
- $\rightarrow$ represents exponentially many (explicit) member states


## Heuristics for Decoupled Search So Far

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buy-leaves compilation
- compile prices of $s^{\mathcal{F}}$ into new task
- evaluate $h$ on compiled task
- problems:
- impractical for abstraction-based heuristics
- pattern selection based on original task


## Alternative Definition for Computing Decoupled Heuristics

contribution: explicit decoupled heuristic

$$
h_{\mathcal{F}, \mathrm{ex}}\left(s^{\mathcal{F}}\right)=\min _{s \in\left[s^{\mathcal{F}}\right]} \operatorname{price}\left(s^{\mathcal{F}}, s\right)+h(s)
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- "best" use of given explicit heuristic
- problem: exponentially many member states


## Single PDBs for Decoupled Search

reminder: explicit decoupled heuristic

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contribution: decoupled PDB

$$
\operatorname{dPDB}\left(h^{P}, s^{\mathcal{F}}\right)=\min _{s^{P} \in S^{P}} \operatorname{price}\left(s^{\mathcal{F}}, s^{P}\right)+h^{P}\left(s^{P}\right)
$$

## Combining Multiple PDBs

## explicit search

- given: $\mathcal{H}=\left\{H_{1}, \ldots, H_{n}\right\}$ with $H_{i}$ additive set of (PDB) heuristics (e.g., disjoint PDBs, cost-partitioned PDBs, etc.)
- canonical combination:

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- how to transfer to decoupled search?


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## properties

- information-lossy: use different minimizing member state for each PDB
- inadmissible: may count prices of leaves multiple times in different heuristics


## Explicit Decoupled Canonical Heuristic (1)

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## Explicit Decoupled Canonical Heuristic (2)

## complexity <br> computing $h_{\mathcal{F}, \text { ex }}^{\mathcal{H}}$ is NP-complete

- practical implementation via branch-and-bound
- incremental computation of member states allows pruning
- worst case: enumeration of all exponentially many member states


## Polynomial-time Approximations (1)

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- admissible, but lossy approximation


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## leaf-disjoint (LD) PDBs

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each PDB affects at most one leaf

- minimize sum of prices and heuristic separately for each set of affected leaves
- heuristic value equals $h_{\mathcal{F}, \mathrm{ex}}^{\mathcal{H}}$


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- alternative way of computing explicit heuristics for decoupled search
- efficient computation of PDBs for decoupled search
- admissible combination of sets of additive PDBs NP-complete
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- efficient computation of PDBs for decoupled search
- admissible combination of sets of additive PDBs NP-complete
- practical implementation and polynomial-time approximations
- future work:
- many results independent of type of heuristic: use different abstractions
- integrate cost partitioning into decoupled search: leaf price partitioning

