

# Negated Occurrences of Predicates in PDDL Axiom Bodies

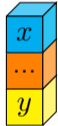
Gabriele Röger   Claudia Grundke

University of Basel

September 24, 2024

# PDDL Axioms

$$\mathit{above}(x, y) \leftarrow \mathit{on}(x, y) \vee \exists z (\mathit{on}(x, z) \wedge \mathit{above}(z, y))$$



# PDDL Axioms

derived predicate      basic predicate

$$\mathit{above}(x, y) \leftarrow \mathit{on}(x, y) \vee \exists z (\mathit{on}(x, z) \wedge \mathit{above}(z, y))$$



# PDDL Axioms

derived predicate

basic predicate

$$\mathit{above}(x, y) \leftarrow \mathit{on}(x, y) \vee \exists z (\mathit{on}(x, z) \wedge \mathit{above}(z, y))$$



$$\mathit{illegal}() \leftarrow \exists x \mathit{above}(x, x)$$



# Axiom Formalisms

$$A(x) \leftarrow \forall y \neg B(y) \wedge \neg C(x)$$

Stratified  
Axioms<sup>[1]</sup>

Semipositive  
Axioms<sup>[2]</sup>

$$A(x) \leftarrow \forall y B(y) \wedge \neg C(x)$$

Stratified  
Datalog

$$A(x) \leftarrow \exists y \neg B(y) \wedge \neg C(x)$$

Semipositive  
Datalog

$$A(x) \leftarrow \exists y B(y) \wedge \neg C(x)$$

[1] Sylvie Thiébaux, Jörg Hoffmann, Bernhard Nebel (2005). In Defense of PDDL Axioms. AIJ.

[2] Stefan Edelkamp, Jörg Hoffmann (2004). PDDL2.2: The Language for the Classical Part of the 4th International Planning Competition. Technical Report.

# Axiom Formalisms

$$A(x) \leftarrow \forall y \neg B(y) \wedge \neg C(x)$$

Stratified  
Axioms<sup>[1]</sup>



Semipositive  
Axioms<sup>[2]</sup>

$$A(x) \leftarrow \forall y B(y) \wedge \neg C(x)$$

Stratified  
Datalog

$$A(x) \leftarrow \exists y \neg B(y) \wedge \neg C(x)$$



Semipositive  
Datalog

$$A(x) \leftarrow \exists y B(y) \wedge \neg C(x)$$

[1] Sylvie Thiébaux, Jörg Hoffmann, Bernhard Nebel (2005). In Defense of PDDL Axioms. AIJ.

[2] Stefan Edelkamp, Jörg Hoffmann (2004). PDDL2.2: The Language for the Classical Part of the 4th International Planning Competition. Technical Report.

# Linear Order

Even number of objects?

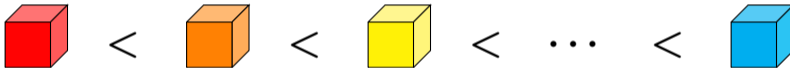


...



# Linear Order

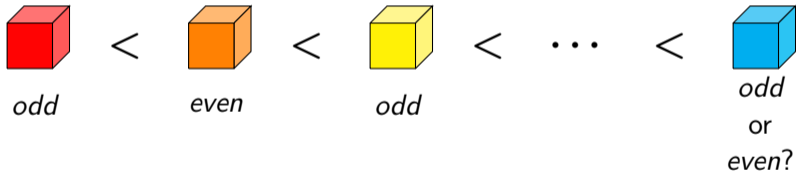
Even number of objects?



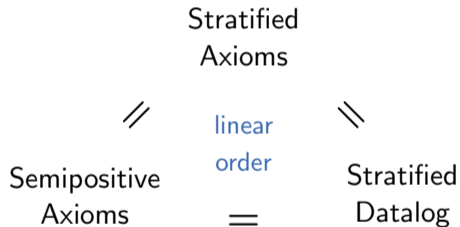


# Linear Order

Even number of objects?



# Capturing PTIME



With linear orders all three axiom formalisms can express any polynomial-time algorithm that can decide if a state is legal.  
(Immerman-Vardi Theorem)

## In General

$$A(x) \leftarrow \forall y \neg B(y) \wedge \neg C(x)$$

Stratified  
Axioms

?

?

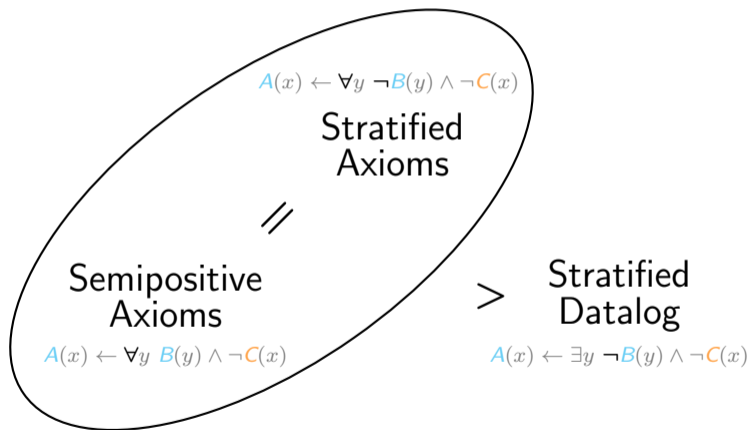
Semipositive  
Axioms

$$A(x) \leftarrow \forall y B(y) \wedge \neg C(x)$$

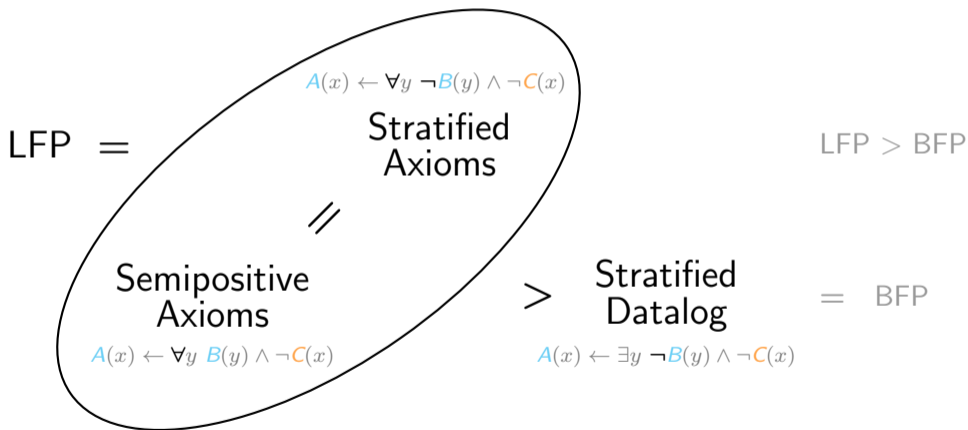
Stratified  
Datalog

$$A(x) \leftarrow \exists y \neg B(y) \wedge \neg C(x)$$

## In General



## In General



# Summary

With linear order:

Stratified Axioms = Semipositive Axioms = Stratified Datalog

In General:

Stratified Axioms = Semipositive Axioms > Stratified Datalog



Link to our ICAPS 2024 paper  
[Formal Representations of  
Classical Planning Domains](#)



Link to our PuK 2024 paper  
[Negated Occurrences of  
Predicates in PDDL Axiom Bodies](#)

# Relationships to Fixed Point Logics

Libkin, 2004 (Corollaries 10.8 & 10.13):

$\text{LFP-sim} = \text{LFP} = \text{LFP}_0$

Trivial:  $\text{AP}_0 \leq \text{AP}$

New:  $\text{LFP}_0 \leq \text{AP}_0$ ,  $\text{AP} \leq \text{LFP-sim}$

**Thus:**  $\text{AP} = \text{AP}_0$

Ebbinghaus and Flum, 1995 (Theorems 7.7.2 & 8.1.1):

$\text{Dat} = \text{BFP}$ ,  $\text{BFP} < \text{LFP}$

**Thus:**  $\text{Dat} < \text{AP}$ ,  $\text{Dat} < \text{AP}_0$

$\text{AP}$  = stratified axioms

$\text{AP}_0$  = semipositive axioms

$\text{Dat}$  = stratified Datalog

$\text{LFP}$  = least fixed-point logic

$\text{LFP-sim}$  =  $\text{LFP}$  with simultaneous fixed points

$\text{LFP}_0$  = first-order logic extended with single fixed point

$\text{BFP}$  = bounded  $\text{LFP}$

---

Libkin, L. (2004). Elements of Finite Model Theory. Springer Berlin, Heidelberg.

Ebbinghaus, H.-D. and Flum, J. (1995). Finite Model Theory. Springer-Verlag.