Generalized Potential Heuristics for Classical Planning

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## Introduction

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Mativat	tion		

- Problems of a classical planning domain share a common structure
- Generalized planning tries to find a solution for a whole domain

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In this '	Work		

- For some domains a solution that solves the whole domain can be described easily
- We try to learn these solutions from small instances

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# Background

Generalized Potential Heuristics

Learning the Heuristic

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### Representing Progress with Heuristic Functions

Descending & dead-end avoiding heuristics (Seipp et al., 2016)

• descending:

all alive states have an improving successor

• dead-end avoiding:

all improving successors of alive states are solvable

A state is *alive* if it is reachable, solvable and not a goal.

Descending, dead-end avoiding heuristics

- guide greedy search to a goal
- use at most  $h(s_0) h(s_g)$  steps
- encode a *measure of progress* (Parmar, 2002)

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Descrip	tion Logic	S	

#### Description Logic $\mathcal{SOI}$ with Role Value Maps

- Primitive concepts
  - represent set of objects with some property
- Primitive roles
  - represent relation between objects

#### Complex concepts

- ⊥,⊤
- ¬*C*
- $C_1 \sqcup C_2$ ,  $C_1 \sqcap C_2$
- ∀*R*.*C*, ∃*R*.*C*
- $R_1 = R_2$
- $\{a_1,\ldots,a_n\}$

Complex roles

- R<sup>-1</sup>
- R<sup>+</sup>
- $R_1 \circ R_2$

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### Description Logic for Planning Domain

Description logic for a planning domain

• Interpretation for every state of any instance

Concepts and roles

- Primitive concept for each unary predicate
  - Example: *clear*
- Primitive role for each binary predicate
  - Example: on
- Primitive concepts and roles for predicates in the goal
  - Example: on<sub>G</sub>

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### Generalized Potential Heuristics

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Generali	zed Poter	ntial Heuristics	

#### Definition (Generalized Potential Heuristic)

Linear combination of features well-defined over all instances:

$$h(s) = \sum_{f \in \mathcal{F}} w(f) \cdot f(s)$$

- We use two types of features based on description logics:
  - cardinality features |C|
  - distance features (see paper)

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Exampl	e: Clearing	g a Block	

• Consider the subset of Blocksworld problems where the goal is to clear a given subset of blocks

Descending and Dead-end Avoiding Generalized Potential Heuristic

$$h(s) = 2 \cdot |C_1| + |C_2|$$

- C<sub>1</sub> ≡ ∃on<sup>+</sup>.clear<sub>G</sub>:
   "Set of blocks above some block that needs to be cleared"
- C<sub>2</sub> ≡ holding:
   "Set of blocks being held"



Existence of Descending and Dead-End Avoiding Heuristics

- We prove that descending, dead-end avoiding generalized heuristics exist for a number of standard domains:
  - Blocksworld
  - Gripper
  - Spanner
  - Miconic
  - Logistics
- Greedy search solves all instances in linear time
- ightarrow The challenge: can we obtain these heuristics automatically?

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## Learning the Heuristic

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Overview of our inductive approach:

- **(**) Fully expand small instances to generate training set S.
- Generate set of generalized features *F* with all features under a certain syntactic complexity.
- Compute simplest potential heuristic on *F* that is descending and dead-end avoiding on states in *S*.
  - If no such h exists, try with larger set  $\mathcal{F}$ .
  - If it does exist, test *h* on unseen instances.

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С	omputing t	he Weig	hts			
	Mixed Intege	r Linear Pr	rogram			
		$\min_{w} \sum_{f \in \mathcal{F}} [v]$	$w_f  eq 0] \mathcal{K}(f)$	subject t	0	
	$\bigvee_{s' \in succ(s)}$	$h(s') + 1 \le 1$	$\leq h(s)$	for alive	states <i>s</i>	
		$h(s') \ge$	$\geq h(s)$	where s	itions ( <i>s</i> , <i>s</i> ′) is alive unsolvable	

 $\bullet$  Solutions map to heuristics that are descending and dead-end avoiding on all states in  $\mathcal{S}...$ 

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С	omputing	the We	ights				
	Mixed Integ	er Linear	Program				
		$\min_{w} \sum_{f \in .}$	$\sum_{\mathcal{F}} [w_f \neq 0] \mathcal{K}(f)$	su	bject to		
	$\bigvee_{s'\in succ(s)}$		$1 \leq h(s)$	fo	r alive states <i>s</i>		
		h(s'	$h') \geq h(s)$	wł	r transitions ( <i>s</i> nere <i>s</i> is alive nd <i>s</i> ′ is unsolva		

- $\bullet$  Solutions map to heuristics that are descending and dead-end avoiding on all states in  $\mathcal{S}...$
- ... and have minimum complexity.

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Results			

- Our approach learns generalized heuristics on standard domains such as Gripper, Miconic, Spanner, VisitAll.
- We have (manually) checked that they are descending and dead-end avoiding on all instances of the domain.
  - Exception VisitAll: No linear solution possible
- Steepest-ascent hill-climbing with these heuristics solves any instances of these domains in linear time.
- Other domains such as Blocksworld appear to need better feature exploration strategies.

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# Conclusion

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Contrib	utions		

- General descending and dead-end avoiding heuristics exist for several planning domains.
- These solve any instance in linear time.
- We can learn them automatically from a suitable logical model and small instances.

#### Discussion and Future Work

- The learned heuristic can be easily interpreted.
- The learned heuristic has only inductive guarantees, but
  - We have shown how it can be refined in an online fashion whenever it doesn't generalize correctly.
  - One could attempt to prove the correctness of the heuristic deductively with an automatic theorem prover.
- Better feature generation methods are necessary to scale up to more complex problems.

## **Bonus Slides**

#### Example: unrestricted Blocksworld instances

$$h_{\mathsf{bw}}(s) = -4|C_6| - |\textit{holding}| - 2|\textit{ontable}| - 2|C_7|,$$

- C<sub>1</sub>: ontable<sub>G</sub> □ ontable
   Blocks that are correctly placed on the table
- C<sub>2</sub>: (∃on<sub>G</sub>.⊤) ⊓ (on = on<sub>G</sub>)
   Blocks that are placed on their target block
- C<sub>3</sub>: ¬(ontable<sub>G</sub> ⊔ ∃on<sub>G</sub>.⊤) Blocks that are not mentioned in the goal
- $C_4$ :  $C_1 \sqcup C_2 \sqcup C_3$ Blocks where block (or table) below is consistent with the goal
- C<sub>5</sub>: ∀on<sup>-1</sup><sub>G</sub>.(on = on<sub>G</sub>)
   Blocks where the block above is consistent with the goal
- C<sub>6</sub>: C<sub>4</sub> □ ∀on<sup>+</sup>.(C<sub>4</sub> □ C<sub>5</sub>) Blocks that are well-placed.
- C<sub>7</sub>: holding □ ∃on<sub>G</sub>.(clear □ C<sub>6</sub>)
   Blocks held while their target block is clear and well-placed.

	G	М	S	v
# of training instances	8	12	11	9
# of iterations	2.0	2.7	1.0	1.7
$ \mathcal{F} $	469	2105	904	330
# of MIP variables	2017	7273	3381	1039
# of MIP constraints	2238	7331	3370	1190
Complexity of h	8 (18)	6 (14)	8 (20)	5 (8)
# of features in $h$	5	4	5	3
Total time	8h	32m	178s	87s
Total MIP time	7.4h	26m	6.8s	2.1s