Discrete Structures

C1. Introduction to Graphs

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Graphs and Directed Graphs

Graphs

Graphs (of various kinds) are ubiquitous in Computer Science and its applications.

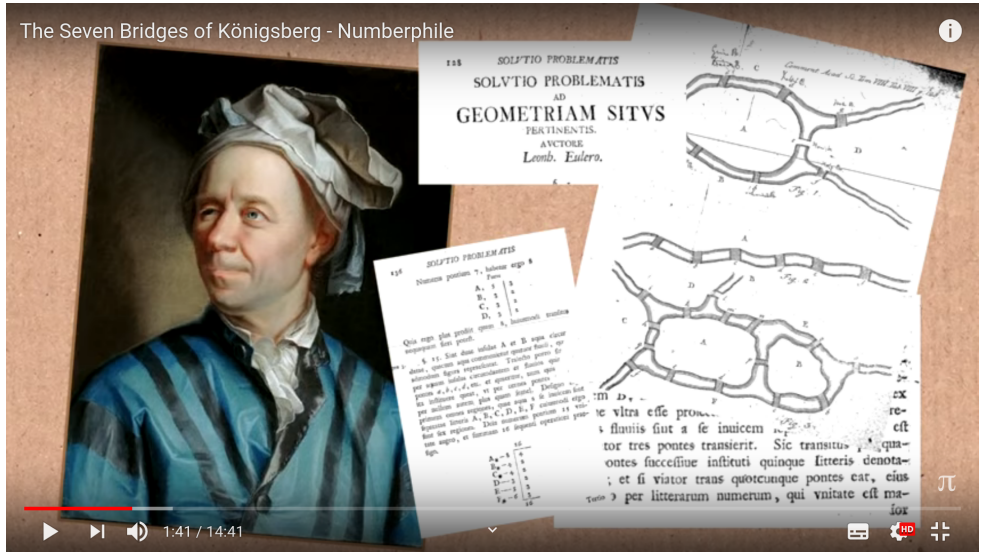
Some examples:

- Boolean circuits in hardware design
- control flow graphs in compilers
- pathfinding in video games
- computer networks
- neural networks
- social networks

Graph Theory

- Graph theory was founded in 1736 by Leonhard Euler's study of the Seven Bridges of Königsberg problem.
- It remains one of the main areas of discrete mathematics to this day.

More on Euler and the Seven Bridges of Königsberg:



- The Seven Bridges of Königsberg – Numberphile.
  `https://youtu.be/W18FDEA1jRQ`

Graphs and Directed Graphs – Definitions

Definition (Graph)

A graph (also: undirected graph) is a pair $G = \langle V, E \rangle$, where

- V is a finite set called the set of vertices, and
- $E \subseteq \{\{u, v\} \subseteq V \mid u \neq v\}$ is called the set of edges.

German: Graph, ungerichteter Graph, Knoten, Kanten

Graphs and Directed Graphs – Definitions

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Definition (Directed Graph)

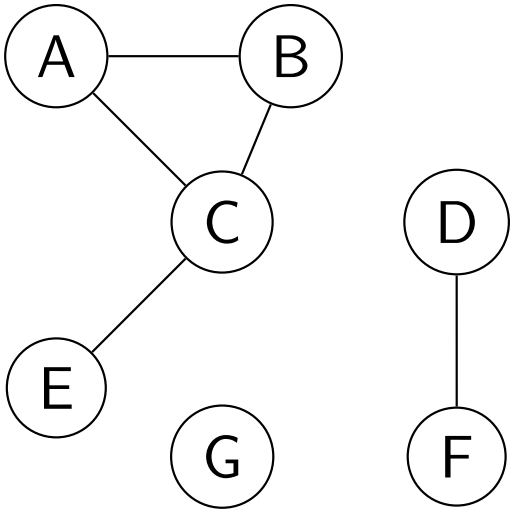
A directed graph (also: digraph) is a pair $G = \langle N, A \rangle$, where

- N is a finite set called the set of nodes, and
- $A \subseteq N \times N$ is called the set of arcs.

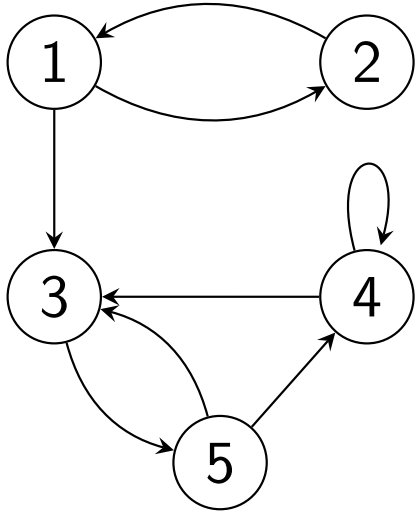
German: gerichteter Graph, Digraph, Knoten, Kanten/Pfeile

Graphs and Directed Graphs – Pictorially

often described pictorially:



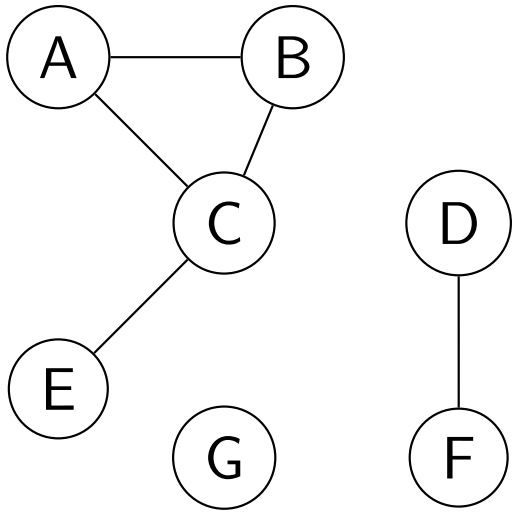
graph $\langle V, E \rangle$



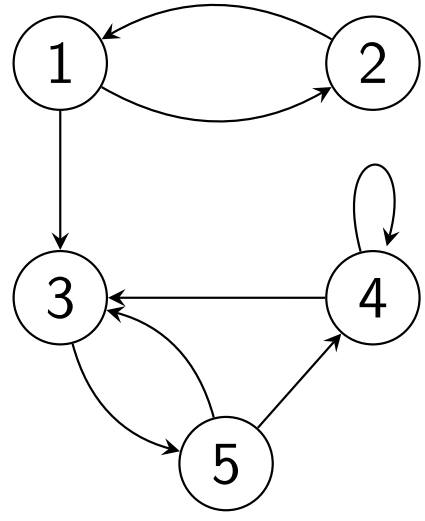
directed graph $\langle N, A \rangle$

Graphs and Directed Graphs – Pictorially

often described pictorially:



graph $\langle V, E \rangle$



directed graph $\langle N, A \rangle$

- $V = \{\mathsf{A}, \mathsf{B}, \mathsf{C}, \mathsf{D}, \mathsf{E}, \mathsf{F}, \mathsf{G}\}$
- $E = \{\{\mathsf{A}, \mathsf{B}\}, \{\mathsf{A}, \mathsf{C}\}, \{\mathsf{B}, \mathsf{C}\}, \{\mathsf{C}, \mathsf{E}\}, \{\mathsf{D}, \mathsf{F}\}\}$

- $N = \{1, 2, 3, 4, 5\}$
- $A = \{\langle 1, 2\rangle, \langle 1, 3\rangle, \langle 2, 1\rangle, \langle 3, 5\rangle, \langle 4, 3\rangle, \langle 4, 4\rangle, \langle 5, 3\rangle, \langle 5, 4\rangle\}$

Relationship to Relations

graphs vs. directed graphs:

- edges are sets of two elements, arcs are pairs
- arcs can be self-loops $\langle v, v \rangle$; edges cannot (why not?)

(di-)graphs vs. relations:

- A directed graph $\langle N, A \rangle$ is essentially identical to
  (= contains the same information as)
  an arbitrary relation R_A over the finite set N:
  $u \, R_A \, v$ iff $\langle u, v \rangle \in A$
- A graph $\langle V, E \rangle$ is essentially identical to
  an irreflexive symmetric relation R_E over the finite set V:
  $u \, R_E \, v$ iff $\{u, v\} \in E$

Other Kinds of Graphs

many variations exist, for example:

- self-loops may be allowed in edges ("non-simple" graphs)
- labeled graphs: additional information associated with vertices and/or edges
- weighted graphs: numbers associated with edges
- multigraphs: multiple edges between same vertices allowed
- mixed graphs: both edges and arcs allowed
- hypergraphs: edges can involve more than 2 vertices
- infinite graphs: may have infinitely many vertices/edges

Graph Terminology

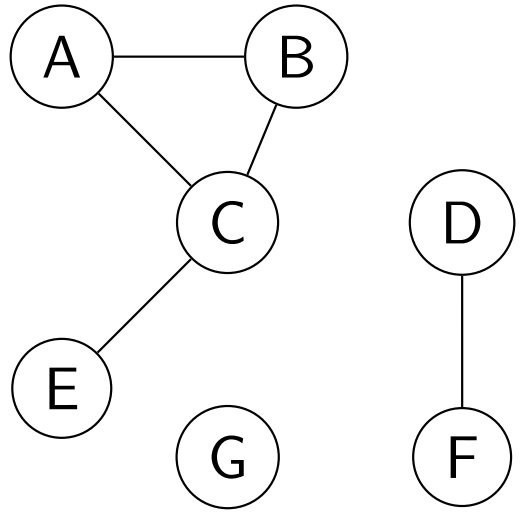
Definition (Graph Terminology)

Let $\langle V, E \rangle$ be a graph.

- u and v are the endpoints of the edge $\{u, v\} \in E$
- u and v are incident to the edge $\{u, v\} \in E$
- u and v are adjacent if $\{u, v\} \in E$
- the vertices adjacent with $v \in V$ are its neighbours $\text{neigh}(v)$: $\text{neigh}(v) = \{w \in V \mid \{v, w\} \in E\}$
- the number of neighbours of $v \in V$ is its degree $\deg(v)$: $\deg(v) = |\text{neigh}(v)|$

German: Endknoten, inzident, adjazent/benachbart, Nachbarn, Grad

Graph Terminology – Examples



endpoints, incident, adjacent, neighbours, degree

Directed Graph Terminology

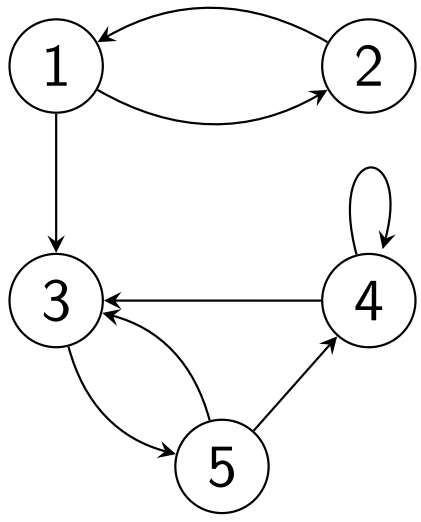
Definition (Directed Graph Terminology)

Let $\langle N, A \rangle$ be a directed graph.

- u is the tail and v is the head of the arc $\langle u, v \rangle \in A$; we say $\langle u, v \rangle$ is an arc from u to v
- u and v are incident to the arc $\langle u, v \rangle \in A$
- u is a predecessor of v and v is a successor of u if $\langle u, v \rangle \in A$
- the predecessors and successor of v are written as $\mathrm{pred}(v) = \{u \in N \mid \langle u, v \rangle \in A\}$ and $\mathrm{succ}(v) = \{w \in N \mid \langle v, w \rangle \in A\}$
- the number of predecessors/successors of $v \in N$ is its indegree/outdegree: $\mathrm{indeg}(v) = |\mathrm{pred}(v)|$, $\mathrm{outdeg}(v) = |\mathrm{succ}(v)|$

German: Fuss, Kopf, inzident, Vorgänger, Nachfolger, Eingangs-/Ausgangsgrad

Directed Graph Terminology – Examples



head, tail, predecessors, successors, indegree, outdegree

Induced Graphs and Degree Lemma

Induced Graph of a Directed Graph

Definition (undirected graph induced by a directed graph)

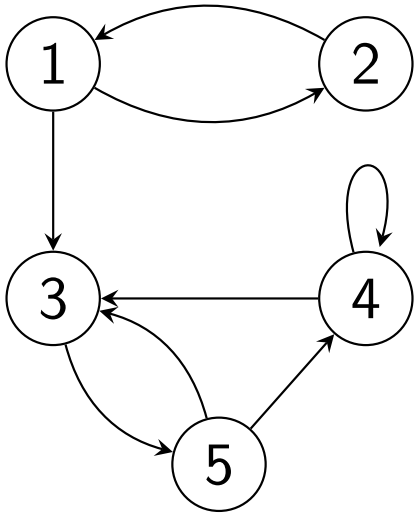
Let $G = \langle N, A \rangle$ be a directed graph.
The (undirected) graph induced by G is the graph $\langle N, E \rangle$ with $E = \{\{u, v\} \mid \langle u, v \rangle \in A, u \neq v\}$.

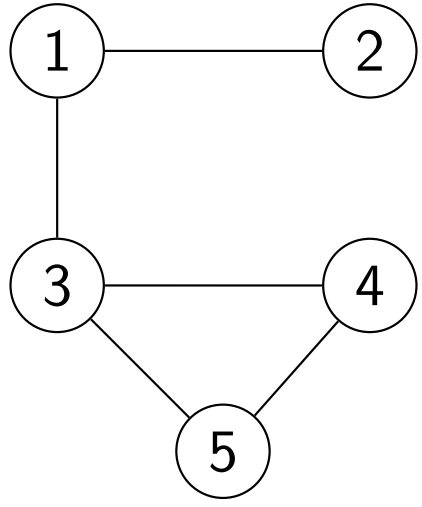
German: induziert

Questions:

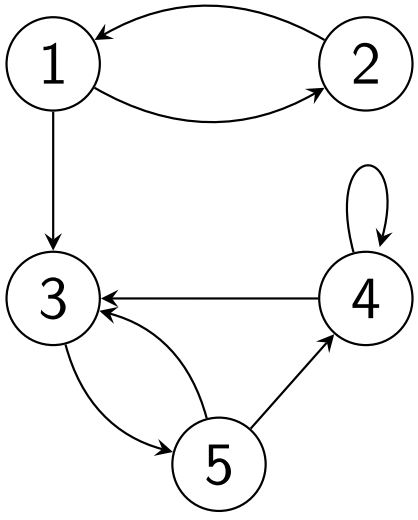
- Why require $u \neq v$?
- If $|N| = n$ and $|A| = m$, how many vertices and edges does the induced graph have?
- How does the answer change if G has no self-loops?

Induced Graph of a Directed Graph – Example

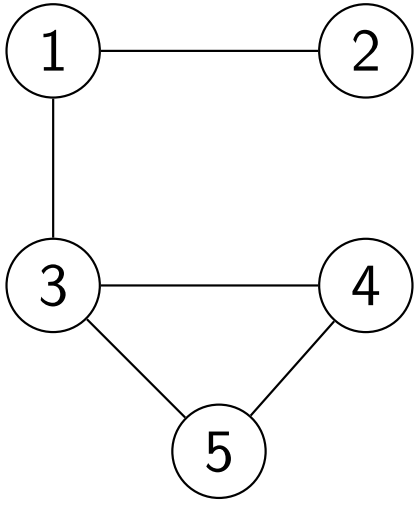




Induced Graph of a Directed Graph – Example



- $N = \{1, 2, 3, 4, 5\}$
- $A = \{\langle 1, 2\rangle, \langle 1, 3\rangle, \langle 2, 1\rangle, \langle 3, 5\rangle, \langle 4, 3\rangle, \langle 4, 4\rangle, \langle 5, 3\rangle, \langle 5, 4\rangle\}$



- $V = \{1, 2, 3, 4, 5\}$
- $E = \{\{1, 2\}, \{1, 3\}, \{3, 4\}, \{3, 5\}, \{4, 5\}\}$

Degree Lemma

Lemma (degree lemma for directed graphs)

Let $\langle N, A \rangle$ be a directed graph.
Then $\sum_{v \in N} \text{indeg}(v) = \sum_{v \in N} \text{outdeg}(v) = |A|$.

Intuitively: every arc contributes 1 to the indegree of one node and 1 to the outdegree of one node.

Degree Lemma

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Intuitively: every arc contributes 1 to the indegree of one node and 1 to the outdegree of one node.

Lemma (degree lemma for undirected graphs)

Let $\langle V, E \rangle$ be a graph.
Then $\sum_{v \in V} \deg(v) = 2|E|$.

Intuitively: every edge contributes 1 to the degree of two vertices.

Degree Lemma

Lemma (degree lemma for directed graphs)

Let $\langle N, A \rangle$ be a directed graph.
Then $\sum_{v \in N} \text{indeg}(v) = \sum_{v \in N} \text{outdeg}(v) = |A|$.

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Lemma (degree lemma for undirected graphs)

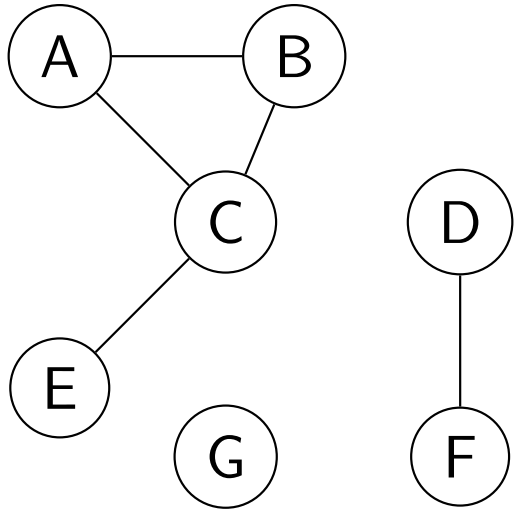
Let $\langle V, E \rangle$ be a graph.
Then $\sum_{v \in V} \deg(v) = 2|E|$.

Intuitively: every edge contributes 1 to the degree of two vertices.

Corollary

Every graph has an even number of vertices with odd degree.

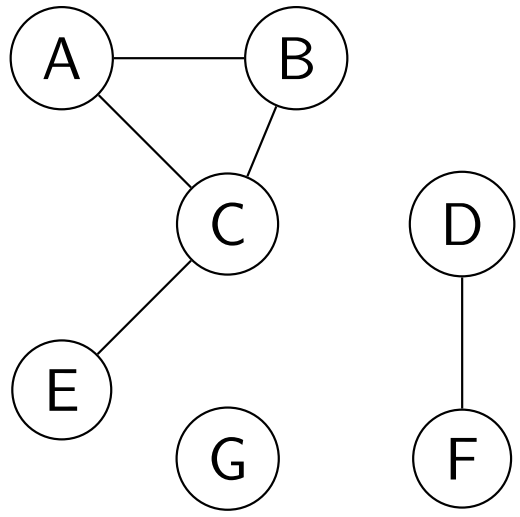
Degree Lemma – Example



$$\sum_{v \in V} \deg(v)$$

$$= \deg(A) + \deg(B) + \deg(C) + \deg(D) + \deg(E) + \deg(F) + \deg(G)$$

Degree Lemma – Example

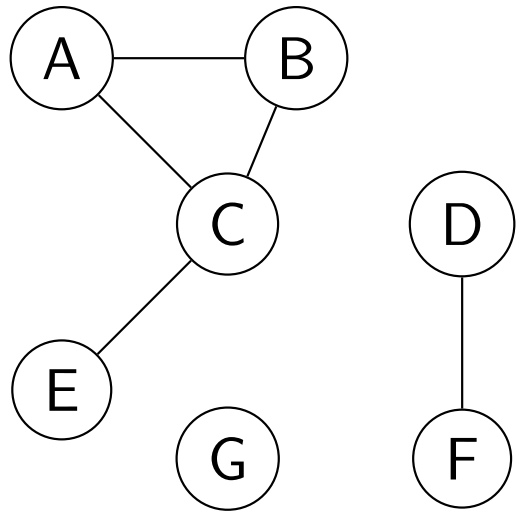


$$\sum_{v \in V} \deg(v)$$

$$= \deg(A) + \deg(B) + \deg(C) + \deg(D) + \deg(E) + \deg(F) + \deg(G)$$

$$= 2 + 2 + 3 + 1 + 1 + 1 + 0$$

Degree Lemma – Example



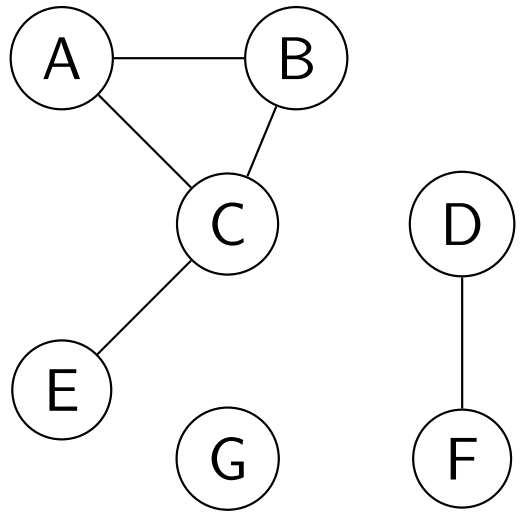
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$$= 10 = 2 \cdot 5 = 2|E|$$

Degree Lemma – Example



$$\sum_{v \in V} \deg(v)$$

$$= \deg(A) + \deg(B) + \deg(C) + \deg(D) + \deg(E) + \deg(F) + \deg(G)$$

$$= 2 + 2 + 3 + 1 + 1 + 1 + 0$$

$$= 10 = 2 \cdot 5 = 2|E|$$

4 vertices with odd degree

Degree Lemma – Proof (1)

Proof of degree lemma for directed graphs.

$$
\begin{aligned}
\sum_{v \in N} \mathrm{indeg}(v) &= \sum_{v \in N} |\mathrm{pred}(v)| \\
&= \sum_{v \in N} |\{u \mid u \in N, \langle u, v \rangle \in A\}| \\
&= \sum_{v \in N} |\{\langle u, v \rangle \mid u \in N, \langle u, v \rangle \in A\}| \\
&= \left| \bigcup_{v \in N} \{\langle u, v \rangle \mid u \in N, \langle u, v \rangle \in A\} \right| \\
&= |\{\langle u, v \rangle \mid u \in N, v \in N, \langle u, v \rangle \in A\}| \\
&= |A|.
\end{aligned}
$$

$\sum_{v \in N} \mathrm{outdeg}(v) = |A|$ is analogous. □

Degree Lemma – Proof (2)

We omit the proof for undirected graphs,
which can be conducted similarly.

One possible proof strategy that reuses the result we proved:

- Define directed graph $\langle V, A \rangle$ from the graph $\langle V, E \rangle$ by orienting each edge into an arc arbitrarily.
- Observe $\deg(v) = \mathrm{indeg}(v) + \mathrm{outdeg}(v)$, where deg refers to the graph and indeg/outdeg to the directed graph.
- Use the degree lemma for directed graphs:
  $\sum_{v \in V} \deg(v) = \sum_{v \in V} (\mathrm{indeg}(v) + \mathrm{outdeg}(v)) =$
  $\sum_{v \in V} \mathrm{indeg}(v) + \sum_{v \in V} \mathrm{outdeg}(v) = |A| + |A| = 2|A| = 2|E|$