Planning and Optimization
F4. Landmarks: Minimum Hitting Set Heuristic

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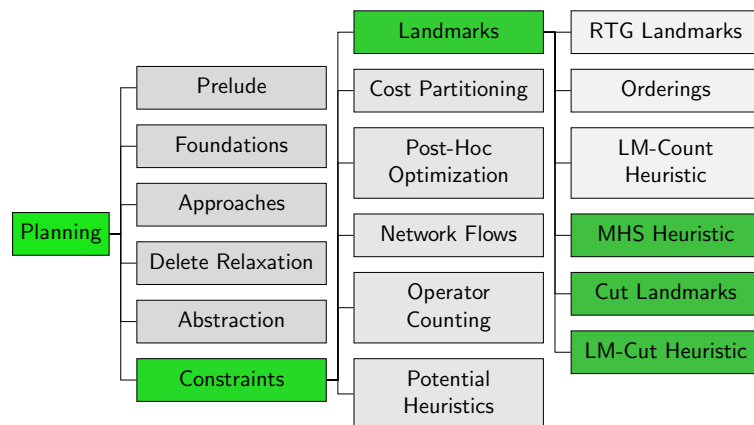

Content of the Course

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Planning ─┬─ Prelude
          ├─ Foundations
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          ├─ Abstraction
          └─ Constraints ─┬─ Landmarks ─────┬─ RTG Landmarks
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                          ├─ Operator          │  └─ LM-Cut Heuristic
                          │  Counting
                          └─ Potential
                             Heuristics
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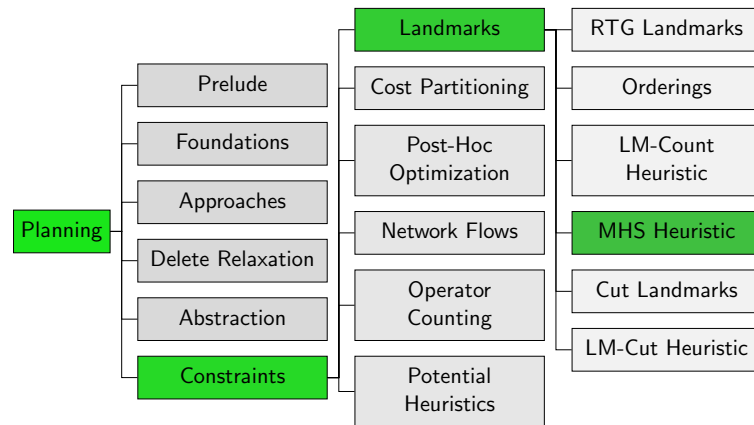
The remaining landmark topics focus on
disjunctive action landmarks.

F4.1 Minimum Hitting Set Heuristic

Content of the Course

Exploiting Disjunctive Action Landmarks

- The cost $cost(L)$ of a disjunctive action landmark L is an admissible heuristic, but it is usually not very informative.
- Landmark heuristics typically aim to combine multiple disjunctive action landmarks.

How can we exploit a given set $\mathcal{L}$ of disjunctive action landmarks?

- Sum of costs $\sum_{L \in \mathcal{L}} cost(L)$?
  ⤳ not admissible!
- Maximize costs $\max_{L \in \mathcal{L}} cost(L)$?
  ⤳ usually very weak heuristic
- better: Hitting sets

Hitting Sets

Definition (Hitting Set)

Let X be a set, $\mathcal{F} = \{F_1, \ldots, F_n\} \subseteq 2^X$ be a family of subsets of X and $c : X \to \mathbb{R}_0^+$ be a cost function for X.

A hitting set is a subset $H \subseteq X$ that "hits" all subsets in $\mathcal{F}$, i.e., $H \cap F \neq \emptyset$ for all $F \in \mathcal{F}$. The cost of H is $\sum_{x \in H} c(x)$.

A minimum hitting set (MHS) is a hitting set with minimal cost.

MHS is a "classical" NP-complete problem (Karp, 1972)

Example: Hitting Sets

Example

$X = \{o_1, o_2, o_3, o_4\}$

$\mathcal{F} = \{\{o_4\}, \{o_1, o_2\}, \{o_1, o_3\}, \{o_2, o_3\}\}$

$c(o_1) = 3, \quad c(o_2) = 4, \quad c(o_3) = 5, \quad c(o_4) = 0$

Specify a minimum hitting set.

Solution: $\{o_1, o_2, o_4\}$ with cost $3 + 4 + 0 = 7$

Hitting Sets for Disjunctive Action Landmarks

Idea: disjunctive action landmarks are interpreted as
instance of minimum hitting set

Definition (Hitting Set Heuristic)
Let $\mathcal{L}$ be a set of disjunctive action landmarks. The hitting set
heuristic $h^{\text{MHS}}(\mathcal{L})$ is defined as the cost of a minimum hitting set
for $\mathcal{L}$ with $c(o) = cost(o)$.

Proposition (Hitting Set Heuristic is Admissible)
Let $\mathcal{L}$ be a set of disjunctive action landmarks for state s.
Then $h^{MHS}(\mathcal{L})$ is an admissible estimate for s.

Hitting Set Heuristic: Discussion

▶ The hitting set heuristic is the best possible heuristic
that only uses the given information...
▶ ...but is NP-hard to compute.
▶ ⇝ Use approximations that can be efficiently computed.
⇒ LP-relaxation, cost partitioning (both discussed later)

F4.2 Summary

Summary

▶ Hitting sets yield the most accurate heuristic for a given set of
disjunctive action landmarks.
▶ The computation of a minimal hitting set is NP-hard.