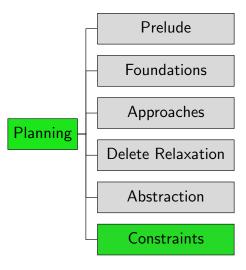
# Planning and Optimization F1. Constraints: Introduction

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#### Content of the Course



### Constraint-based Heuristics

### Coming Up with Heuristics in a Principled Way

#### General Procedure for Obtaining a Heuristic

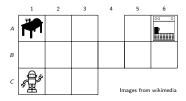
Solve a simplified version of the problem.

#### Major ideas for heuristics in the planning literature:

- delete relaxation
- abstraction
- critical paths
- landmarks
- network flows
- potential heuristic

Landmarks, network flows and potential heuristics are based on constraints that can be specified for a planning task.

### Constraints: Example



FDR planning task  $\langle V, I, O, \gamma \rangle$  with

- $lackbox{ }V=\{\textit{robot-at},\textit{dishes-at}\}\ \text{with}$ 
  - $dom(robot-at) = \{A1, ..., C3, B4, A5, ..., B6\}$
  - $\bullet dom(\textit{dishes-at}) = \{Table, Robot, Dishwasher\}$
- $I = \{ robot\text{-}at \mapsto C1, dishes\text{-}at \mapsto Table \}$
- operators
  - move-x-y to move from cell x to adjacent cell y
  - pickup dishes, and
  - load dishes into the dishwasher.
- $\gamma = (robot-at = B6) \land (dishes-at = Dishwasher)$

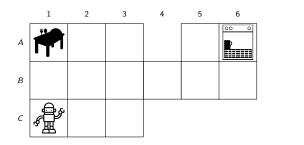
#### Constraints

Some heuristics exploit constraints that describe something that holds in every solution of the task.

For instance, every solution is such that

a variable takes a certain value in at least one visited state. (a fact landmark constraint)

Which values do robot-at and dishes-at take in every solution?

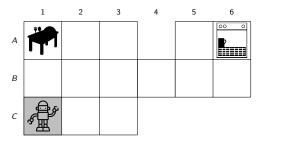








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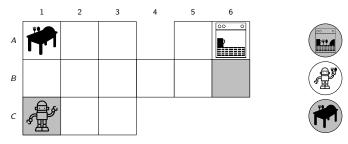






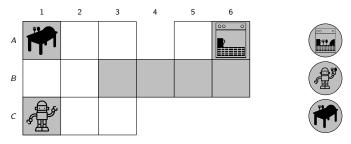
■ robot-at = C1, dishes-at = Table (initial state)

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- robot-at = B6, dishes-at = Dishwasher (goal state)

#### Which values do robot-at and dishes-at take in every solution?



- robot-at = C1, dishes-at = Table (initial state)
- robot-at = B6, dishes-at = Dishwasher (goal state)
- robot-at = A1, robot-at = B3, robot-at = B4, robot-at = B5, robot-at = A6, dishes-at = Robot

#### Constraints

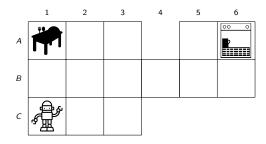
Some heuristics exploit constraints that describe something that holds in every solution of the task.

For instance, every solution is such that

- a variable takes some value in at least one visited state.
   (a fact landmark constraint)
- an action must be applied.(an action landmark constraint)

### Action Landmarks: Example

#### Which actions must be applied in every solution?



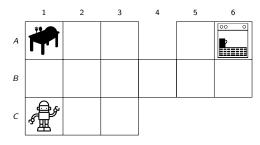


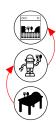




### Action Landmarks: Example

#### Which actions must be applied in every solution?

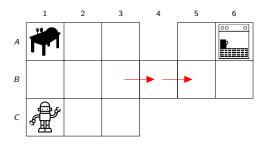


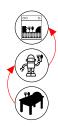


- pickup
- load

### Action Landmarks: Example

#### Which actions must be applied in every solution?





- pickup
- load
- move-B3-B4
- move-B4-B5

#### Constraints

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#### Constraints

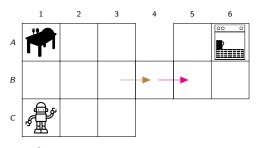
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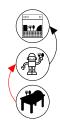
For instance, every solution is such that

- a variable takes some value in at least one visited state.
   (a fact landmark constraint)
- at least one action from a set of actions must be applied.
   (a disjunctive action landmark constraint)

### Disjunctive Action Landmarks: Example

Which set of actions is such that at least one must be applied?

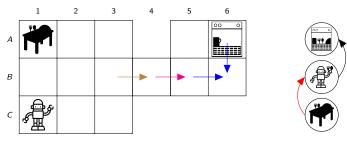




- {pickup}
- {load}
- {move-B3-B4}
- {move-B4-B5}

### Disjunctive Action Landmarks: Example

Which set of actions is such that at least one must be applied?



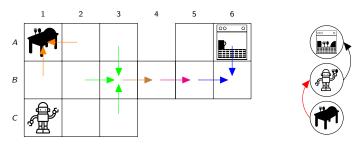
{pickup}

■ {move-A6-B6, move-B5-B6}

- {load}
- {move-B3-B4}
- {move-B4-B5}

### Disjunctive Action Landmarks: Example

#### Which set of actions is such that at least one must be applied?



- {pickup}
- {load}
- {move-B3-B4}
- {move-B4-B5}

- {move-A6-B6, move-B5-B6}
- {move-A3-B3, move-B2-B3, move-C3-B3}
- {move-B1-A1, move-A2-A1}
- **...**

#### Constraints

Some heuristics exploit constraints that describe something that holds in every solution of the task.

For instance, every solution is such that

- a variable takes some value in at least one visited state.
   (a fact landmark constraint)
- at least one action from a set of actions must be applied.
   (a disjunctive action landmark constraint)
- fact consumption and production is "balanced".
   (a network flow constraint)

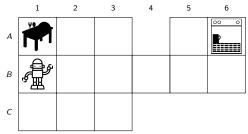
### Network Flow: Example

Consider the fact robot-at = B1. How often are actions used that enter this cell?

	1	2	3	4	5	6
Α						00 0 P
В						
С						

### Network Flow: Example

Consider the fact robot-at = B1. How often are actions used that enter this cell?



Answer: as often as actions that leave this cell

If Count $_o$  denotes how often operator o is applied, we have:

$$\begin{split} \mathsf{Count}_{\mathsf{move-A1-B1}} + \mathsf{Count}_{\mathsf{move-B2-B1}} + \mathsf{Count}_{\mathsf{move-C1-B1}} = \\ \mathsf{Count}_{\mathsf{move-B1-A1}} + \mathsf{Count}_{\mathsf{move-B1-B2}} + \mathsf{Count}_{\mathsf{move-B1-C1}} \end{split}$$

## Multiple Heuristics

### Combining Admissible Heuristics Admissibly

#### Major ideas to combine heuristics admissibly:

- maximize
- canoncial heuristic (for abstractions)
- minimum hitting set (for landmarks)
- cost partitioning
- operator counting

Often computed as solution to a (integer) linear program.

### Combining Heuristics Admissibly: Example

#### Example

Consider an FDR planning task  $\langle V, I, \{o_1, o_2, o_3, o_4\}, \gamma \rangle$  with  $V = \{v_1, v_2, v_3\}$  with  $dom(v_1) = \{A, B\}$  and  $dom(v_2) = dom(v_3) = \{A, B, C\}, I = \{v_1 \mapsto A, v_2 \mapsto A, v_3 \mapsto A\},$   $o_1 = \langle v_1 = A, v_1 := B, 1 \rangle$ 

$$o_1 = \langle v_1 - A, v_1 := B, 1 \rangle$$
 $o_2 = \langle v_2 = A \land v_3 = A, v_2 := B \land v_3 := B, 1 \rangle$ 
 $o_3 = \langle v_2 = B, v_2 := C, 1 \rangle$ 
 $o_4 = \langle v_3 = B, v_3 := C, 1 \rangle$ 

and 
$$\gamma = (v_1 = B) \land (v_2 = C) \land (v_3 = C)$$
.

Let C be the pattern collection that contains all atomic projections. What is the canonical heuristic function  $h^{C}$ ?

### Combining Heuristics Admissibly: Example

#### Example

Consider an FDR planning task  $\langle V, I, \{o_1, o_2, o_3, o_4\}, \gamma \rangle$  with  $V = \{v_1, v_2, v_3\}$  with  $dom(v_1) = \{A, B\}$  and  $dom(v_2) = dom(v_3) = \{A, B, C\}, I = \{v_1 \mapsto A, v_2 \mapsto A, v_3 \mapsto A\},$ 

$$o_1 = \langle v_1 = A, v_1 := B, 1 \rangle$$
  
 $o_2 = \langle v_2 = A \land v_3 = A, v_2 := B \land v_3 := B, 1 \rangle$   
 $o_3 = \langle v_2 = B, v_2 := C, 1 \rangle$   
 $o_4 = \langle v_3 = B, v_3 := C, 1 \rangle$ 

and 
$$\gamma = (v_1 = B) \land (v_2 = C) \land (v_3 = C)$$
.

Let C be the pattern collection that contains all atomic projections. What is the canonical heuristic function  $h^C$ ?

Answer: Let  $h_i := h^{v_i}$ . Then  $h^C = \max\{h_1 + h_2, h_1 + h_3\}$ .

### Reminder: Orthogonality and Additivity

Why can we add  $h_1$  and  $h_2$  ( $h_1$  and  $h_3$ ) admissibly?

#### Theorem (Additivity for Orthogonal Abstractions)

Let  $h^{\alpha_1}, \ldots, h^{\alpha_n}$  be abstraction heuristics of the same transition system such that  $\alpha_i$  and  $\alpha_j$  are orthogonal for all  $i \neq j$ .

Then  $\sum_{i=1}^{n} h^{\alpha_i}$  is a safe, goal-aware, admissible and consistent heuristic for  $\Pi$ .

The proof exploits that every concrete transition induces state-changing transition in at most one abstraction.

### Combining Heuristics (In)admissibly: Example

Let 
$$h = h_1 + h_2 + h_3$$
.  
 $o_2, o_3, o_4$ 
 $o_2, o_3, o_4$ 
 $h_1$ 
 $O_1, O_4$ 
 $O_1, O_2$ 
 $O_1, O_2$ 
 $O_1, O_2$ 
 $O_1, O_2$ 
 $O_1, O_2$ 
 $O_2, O_3, O_4$ 
 $O_1, O_2$ 
 $O_1, O_2$ 
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 $O_1, O_3$ 
 $O_1, O_2$ 
 $O_2, O_3$ 
 $O_1,$ 

 $\langle o_2, o_3, o_4 \rangle$  is a plan for  $s = \langle B, A, A \rangle$  but h(s) = 4.

### Combining Heuristics (In)admissibly: Example

Let 
$$h = h_1 + h_2 + h_3$$
.

 $o_2, o_3, o_4$ 
 $o_1, o_2$ 
 $o_1, o_2$ 
 $o_1, o_2$ 
 $o_1, o_2$ 
 $o_1, o_2$ 
 $o_2, o_3, o_4$ 
 $o_1, o_4$ 
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 $o_2, o_3, o_4$ 
 $o_1, o_2$ 

 $\langle o_2, o_3, o_4 \rangle$  is a plan for  $s = \langle B, A, A \rangle$  but h(s) = 4. Heuristics  $h_2$  and  $h_3$  both account for the single application of  $o_2$ .

### Prevent Inadmissibility

The reason that  $h_2$  and  $h_3$  are not additive is because the cost of  $o_2$  is considered in both.

Is there anything we can do about this?

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The reason that  $h_2$  and  $h_3$  are not additive is because the cost of  $o_2$  is considered in both.

Is there anything we can do about this?

Solution: We can ignore the cost of  $o_2$  in one heuristic by setting its cost to 0 (e.g.,  $cost_3(o_2) = 0$ ).

### Combining Heuristics Admissibly: Example

Let 
$$h' = h_1 + h_2 + h'_3$$
, where  $h'_3 = h^{v_3}$  assuming  $cost_3(o_2) = 0$ .

 $o_2, o_3, o_4$ 
 $o_1, o_2$ 
 $o_1, o_2$ 
 $o_1, o_2$ 
 $o_2, o_3, o_4$ 
 $o_1, o_2$ 
 $o_1, o_2$ 
 $o_1, o_2$ 
 $o_1, o_2$ 
 $o_2, o_3, o_4$ 
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 $o_1, o_2$ 
 $o_1, o_2$ 
 $o_1, o_2$ 
 $o_1, o_2$ 
 $o_1, o_2$ 
 $o_2, o_3, o_4$ 
 $o_1, o_4$ 

 $\langle o_2, o_3, o_4 \rangle$  is an optimal plan for  $s = \langle B, A, A \rangle$  and h'(s) = 3 an admissible estimate.

### Cost partitioning

Using the cost of every operator only in one heuristic is called a zero-one cost partitioning.

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More generally, heuristics are additive if all operator costs are distributed in a way that the sum of the individual costs is no larger than the cost of the operator.

This can also be expressed as a constraint, the cost partitioning constraint:

$$\sum_{i=1}^n cost_i(o) \leq cost(o) \text{ for all } o \in O$$

(more details later)

# Summary

### Summary

- Landmarks and network flows are constraints that describe something that holds in every solution of the task.
- Heuristics can be combined admissibly if the cost partitioning constraint is satisfied.