

Planning and Optimization

E12. Merge-and-Shrink: Merge Strategies & Outlook

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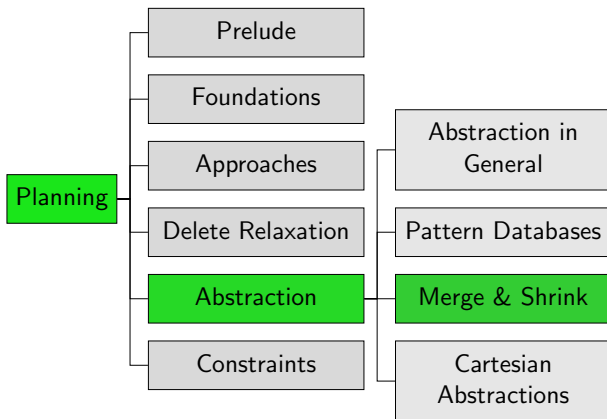
E12.1 Merge Strategies

E12.2 Outlook: Label Reduction and Pruning

E12.3 Summary

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Content of the Course



E12.1 Merge Strategies

Reminder: Generic Algorithm Template

Generic Merge & Shrink Algorithm for planning task Π

$F := F(\Pi)$

while $|F| > 1$:

select $type \in \{\text{merge}, \text{shrink}\}$

if $type = \text{merge}$:

select $\mathcal{T}_1, \mathcal{T}_2 \in F$

$F := (F \setminus \{\mathcal{T}_1, \mathcal{T}_2\}) \cup \{\mathcal{T}_1 \otimes \mathcal{T}_2\}$

if $type = \text{shrink}$:

select $\mathcal{T} \in F$

choose an abstraction mapping β on $\mathcal{T}$

$F := (F \setminus \{\mathcal{T}\}) \cup \{\mathcal{T}^\beta\}$

return the remaining factor $\mathcal{T}^\alpha$ in F

Remaining Question:

- Which abstractions to select for merging? $\rightsquigarrow$ **merge strategy**

Linear vs. Non-linear Merge Strategies

Linear Merge Strategy

In each iteration after the first, choose the abstraction computed in the previous iteration as $\mathcal{T}_1$.

Rationale: only maintains one “complex” abstraction at a time

- ▶ Fully defined by an ordering of atomic projections/variables.
- ▶ Each merge-and-shrink heuristic computed with a non-linear merge strategy can also be computed with a linear merge strategy.
- ▶ However, linear merging can require a super-polynomial blow-up of the final representation size.
- ▶ Recent research turned from linear to non-linear strategies, also because better label reduction techniques (later in this chapter) enabled a more efficient computation.

Classes of Merge Strategies

We can distinguish two major types of merge strategies:

- ▶ **precomputed merge strategies** fix a unique merge order up-front.
One-time effort but cannot react to other transformations applied to the factors.
- ▶ **stateless merge strategies** only consider the current FTS and decide what factors to merge.
Typically computing a score for each pair of factors and naturally non-linear; easy to implement but cannot capture dependencies between more than two factors.

Hybrid strategies combine ideas from precomputed and stateless strategies.

Example Linear Precomputed Merge Strategy

Idea: Use similar causal graph criteria as for growing patterns.

Example: Strategy of h_{HHH}

h_{HHH} : Ordering of atomic projections

- ▶ Start with a goal variable.
- ▶ Add variables that appear in preconditions of operators affecting previous variables.
- ▶ If that is not possible, add a goal variable.

Rationale: increases h quickly

Example Non-linear Precomputed Merge Strategy

Idea: Build clusters of variables with strong interactions and first merge variables within each cluster.

Example: MIASM (“maximum intermediate abstraction size minimizing merging strategy”)

MIASM strategy

- ▶ Measure interaction by ratio of unnecessary states in the merged system (= states not traversed by any abstract plan).
- ▶ Best-first search to identify interesting variable sets.
- ▶ Disjoint variable sets chosen by a greedy algorithm for maximum weighted set packing.

Rationale: increase power of pruning (later in this chapter)

Example Non-linear Stateless Merge Strategy

Idea: Preferably merge transition systems that must synchronize on labels that occur close to a goal state.

Example: DFP (named after Dräger, Finkbeiner and Podelski)

DFP strategy

- ▶ $labelrank(\ell, \mathcal{T}) = \min\{h^*(t) \mid \langle s, \ell, t \rangle \text{ transition in } \mathcal{T}\}$
- ▶ $score(\mathcal{T}, \mathcal{T}') = \min\{\max\{labelrank(\ell, \mathcal{T}), labelrank(\ell, \mathcal{T}')\} \mid \ell \text{ label in } \mathcal{T} \text{ and } \mathcal{T}'\}$
- ▶ Select two transition systems with minimum score.

Rationale: abstraction fine-grained in the goal region, which is likely to be searched by A^* .

Example Hybrid Merge Strategy

Idea: first combine the variables within each strongly connected component of the causal graph.

Example: SCC framework

SCC strategy

- ▶ Compute strongly connected components of causal graph
- ▶ Secondary strategies for order in which
 - ▶ the SCCs are considered (e.g. topologic order),
 - ▶ the factors within an SCC are merged, and
 - ▶ the resulting product systems are merged.

Rationale: reflect strong interactions of variables well

State of the art: SCC+DFP or a stateless MIASM variant

E12.2 Outlook: Label Reduction and Pruning

Further Transformations

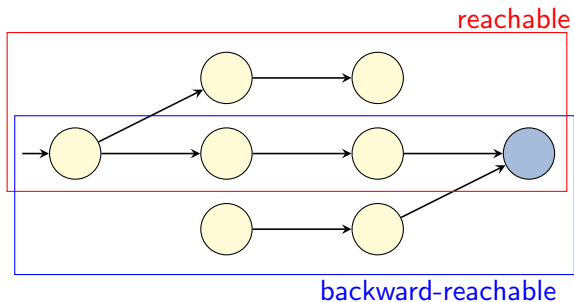
State-of-the-art Merge & Shrink uses two further transformations:

- ▶ Label reduction
- ▶ Pruning

Label Reduction

- ▶ Do no longer distinguish certain **labels**, similar to abstraction that does not distinguish certain states.
- ▶ A **label reduction** $\langle \lambda, c' \rangle$ for a FTS F with label set L is given by a function $\lambda : L \rightarrow L'$, where L' is an arbitrary set of labels, and a label cost function c' on L' such that for all $\ell \in L$, $c'(\lambda(\ell)) \leq c(\ell)$.
The label-reduced TSs have L' and c' for the labels and cost, and in each transition the original label ℓ is replaced with $\lambda(\ell)$.
- ▶ Label reduction is a **conservative** transformation.
- ▶ There are also clear criteria when label reduction is **exact**.
- ▶ **Reduces the time and memory requirement** for merge and shrink steps and **enables coarser bisimulation abstractions**.

Alive States



- ▶ state s is **reachable** if we can reach it from the initial state
- ▶ state s is **backward-reachable** if we can reach the goal from s
- ▶ state s is **alive** if it is reachable and backward-reachable
→ only alive states can be traversed by a solution
- ▶ a state s is **dead** if it is not alive.

Pruning States (1)

- ▶ If in a factor, state s is dead/not backward-reachable then all states that “cover” s in a synchronized product are dead/not backward-reachable in the synchronized product.
- ▶ Removing such states and all adjacent transitions in a factor does not remove any solutions from the synchronized product.
- ▶ This pruning leads to states in the original state space for which the merge-and-shrink abstraction does not define an abstract state.
→ use heuristic estimate ∞

Pruning States (2)

- ▶ Keeping **exactly all backward-reachable states** we still obtain safe, consistent, goal-aware and admissible (with conservative transformations) or perfect heuristics (with exact transformations).
- ▶ Pruning unreachable, backward-reachable states can render the heuristic unsafe because pruned states lead to infinite estimates.
- ▶ However, all reachable states in the original state space will have admissible estimates, so we can use the heuristic like an admissible one in a forward state-space search such as A^* (but not in other contexts like such as orbit search).

We usually prune all dead states to keep the factors small.

E12.3 Summary

Summary

- ▶ There is a wide range of merge strategies. We only covered some important ones.
- ▶ **Label reduction** is crucial for the performance of the merge-and-shrink algorithm, especially when using bisimilarity for shrinking.
- ▶ **Pruning** is used to keep the size of the factors small. It depends on the intended application how aggressive the pruning can be.

E12.4 Literature

Literature (1)

References on merge-and-shrink abstractions:



Klaus Dräger, Bernd Finkbeiner and Andreas Podelski.

Directed Model Checking with Distance-Preserving Abstractions.

Proc. SPIN 2006, pp. 19–34, 2006.

Introduces merge-and-shrink abstractions (for model checking) and **DFP** merging strategy.



Malte Helmert, Patrik Haslum and Jörg Hoffmann.

Flexible Abstraction Heuristics for Optimal Sequential Planning.

Proc. ICAPS 2007, pp. 176–183, 2007.

Introduces merge-and-shrink abstractions **for planning**.

Literature (2)



Raz Nissim, Jörg Hoffmann and Malte Helmert.
Computing Perfect Heuristics in Polynomial Time:
On Bisimulation and Merge-and-Shrink Abstractions
in Optimal Planning.

Proc. IJCAI 2011, pp. 1983–1990, 2011.

Introduces **bisimulation-based shrinking**.



Malte Helmert, Patrik Haslum, Jörg Hoffmann
and Raz Nissim.

Merge-and-Shrink Abstraction: A Method
for Generating Lower Bounds in Factored State Spaces.

Journal of the ACM 61 (3), pp. 16:1–63, 2014.

Detailed **journal version** of the previous two publications.

Literature (3)



Silvan Sievers, Martin Wehrle and Malte Helmert.

Generalized Label Reduction for Merge-and-Shrink Heuristics.

Proc. AAAI 2014, pp. 2358–2366, 2014.

Introduces modern version of **label reduction**.

(There was a more complicated version before.)



Gaojian Fan, Martin Müller and Robert Holte.

Non-linear merging strategies for merge-and-shrink
based on variable interactions.

Proc. SoCS 2014, pp. 53–61, 2014.

Introduces UMC and **MIASM merging strategies**

Literature (4)



Malte Helmert, Gabriele Röger and Silvan Sievers.

On the Expressive Power of Non-Linear Merge-and-Shrink Representations.

Proc. ICAPS 2015, pp. 106–114, 2015.

Shows that **linear merging can require a super-polynomial blow-up** in representation size.



Silvan Sievers and Malte Helmert.

Merge-and-Shrink: A Compositional Theory of Transformations of Factored Transition Systems.

JAIR 71, pp. 781–883, 2021.

Detailed theoretical analysis of task transformations as **sequence of transformations**.

Literature (5)



Silvan Sievers, Florian Pommerening , Thomas Keller and Malte Helmert.

Cost-Partitioned Merge-and-Shrink Heuristics for Optimal Classical Planning.

Proc. IJCAI 2020, pp. 4152–4160, 2020.

Extends **saturated cost partitioning** to merge-and-shrink.