### Planning and Optimization

E5. Abstractions: Additive Abstractions

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# Planning and Optimization November 10, 2025 — E5. Abstractions: Additive Abstractions

E5.1 Additivity

E5.2 Outlook

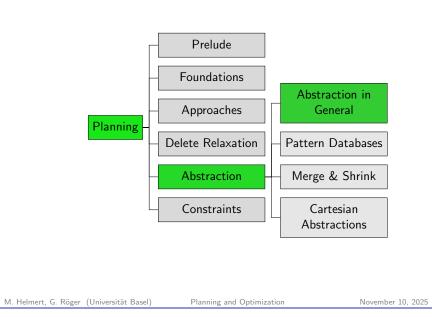
E5.3 Summary

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### Content of the Course



E5. Abstractions: Additive Abstractions

Additivity

E5.1 Additivity

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### Orthogonality of Abstractions

### Definition (Orthogonal)

Let  $\alpha_1$  and  $\alpha_2$  be abstractions of transition system  $\mathcal{T}$ .

We say that  $\alpha_1$  and  $\alpha_2$  are orthogonal if for all transitions  $s \xrightarrow{\ell} t$  of  $\mathcal{T}$ , we have  $\alpha_1(s) = \alpha_1(t)$  or  $\alpha_2(s) = \alpha_2(t)$ .

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### Affecting Transition Labels

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### Definition (Affecting Transition Labels)

Let  $\mathcal T$  be a transition system, and let  $\ell$  be one of its labels. We say that  $\ell$  affects  $\mathcal T$  if  $\mathcal T$  has a transition  $s \xrightarrow{\ell} t$  with  $s \neq t$ .

### Theorem (Affecting Labels vs. Orthogonality)

Let  $\alpha_1$  and  $\alpha_2$  be abstractions of transition system  $\mathcal{T}$ . If no label of  $\mathcal{T}$  affects both  $\mathcal{T}^{\alpha_1}$  and  $\mathcal{T}^{\alpha_2}$ , then  $\alpha_1$  and  $\alpha_2$  are orthogonal.

(Easy proof omitted.)

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Additivity

### Orthogonal Abstractions: Example

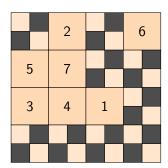
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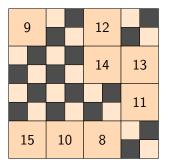
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Are the abstractions orthogonal?

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### Orthogonal Abstractions: Example





Are the abstractions orthogonal?

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### Orthogonality and Additivity

### Theorem (Additivity for Orthogonal Abstractions)

Let  $h^{\alpha_1}, \ldots, h^{\alpha_n}$  be abstraction heuristics of the same transition system such that  $\alpha_i$  and  $\alpha_i$  are orthogonal for all  $i \neq j$ .

Then  $\sum_{i=1}^{n} h^{\alpha_i}$  is a safe, goal-aware, admissible and consistent heuristic for  $\Pi$ .

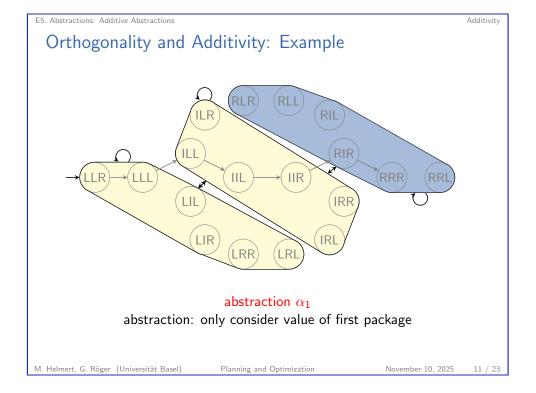
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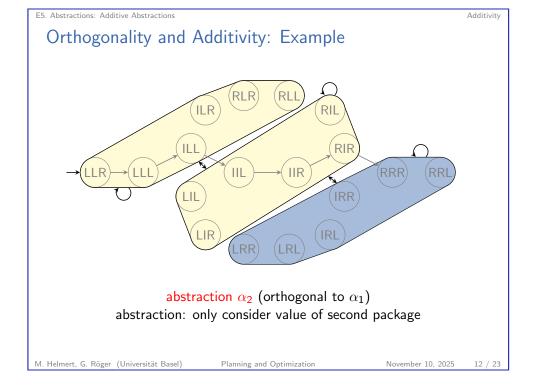
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Additivity

### Orthogonality and Additivity: Proof (1)

### Proof.

We prove goal-awareness and consistency; the other properties follow from these two.

Let  $\mathcal{T} = \langle S, L, c, \mathcal{T}, s_0, S_{\star} \rangle$  be the concrete transition system.

Let  $h = \sum_{i=1}^n h^{\alpha_i}$ .

Goal-awareness: For goal states  $s \in S_{\star}$ ,

 $h(s) = \sum_{i=1}^{n} h^{\alpha_i}(s) = \sum_{i=1}^{n} 0 = 0$  because all individual abstraction heuristics are goal-aware.

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### Orthogonality and Additivity: Proof (2)

### Proof (continued).

Consistency: Let  $s \stackrel{o}{\to} t \in T$ . We must prove  $h(s) \leq c(o) + h(t)$ .

Because the abstractions are orthogonal,  $\alpha_i(s) \neq \alpha_i(t)$ 

for at most one  $i \in \{1, \ldots, n\}$ .

Case 1:  $\alpha_i(s) = \alpha_i(t)$  for all  $i \in \{1, ..., n\}$ .

Then 
$$h(s) = \sum_{i=1}^{n} h^{\alpha_i}(s)$$
  
 $= \sum_{i=1}^{n} h^*_{\mathcal{T}^{\alpha_i}}(\alpha_i(s))$   
 $= \sum_{i=1}^{n} h^*_{\mathcal{T}^{\alpha_i}}(\alpha_i(t))$   
 $= \sum_{i=1}^{n} h^{\alpha_i}(t)$   
 $= h(t) \le c(o) + h(t)$ .

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### Orthogonality and Additivity: Proof (3)

Proof (continued).

Case 2:  $\alpha_i(s) \neq \alpha_i(t)$  for exactly one  $i \in \{1, ..., n\}$ .

Let  $k \in \{1, ..., n\}$  such that  $\alpha_k(s) \neq \alpha_k(t)$ .

Then 
$$h(s) = \sum_{i=1}^{n} h^{\alpha_i}(s)$$
  
 $= \sum_{i \in \{1,...,n\} \setminus \{k\}} h^*_{\mathcal{T}^{\alpha_i}}(\alpha_i(s)) + h^{\alpha_k}(s)$   
 $\leq \sum_{i \in \{1,...,n\} \setminus \{k\}} h^*_{\mathcal{T}^{\alpha_i}}(\alpha_i(t)) + c(o) + h^{\alpha_k}(t)$   
 $= c(o) + \sum_{i=1}^{n} h^{\alpha_i}(t)$   
 $= c(o) + h(t),$ 

where the inequality holds because  $\alpha_i(s) = \alpha_i(t)$  for all  $i \neq k$  and  $h^{\alpha_k}$  is consistent.

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Outlook

E5.2 Outlook

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### Using Abstraction Heuristics in Practice

In practice, there are conflicting goals for abstractions:

- we want to obtain an informative heuristic, but
- want to keep its representation small.

Abstractions have small representations if

- there are few abstract states and
- $\blacktriangleright$  there is a succinct encoding for  $\alpha$ .

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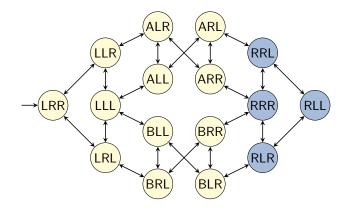
# Counterexample: One-State Abstraction Outlook Counterexample: One-State Abstraction One-state abstraction: $\alpha(s) := \text{const.}$ + very few abstract states and succinct encoding for $\alpha$ - completely uninformative heuristic

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## Counterexample: Identity Abstraction

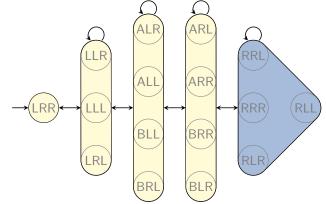


Identity abstraction:  $\alpha(s) := s$ .

- $+\,$  perfect heuristic and succinct encoding for  $\alpha$
- too many abstract states

Outlook

# Counterexample: Perfect Abstraction



Perfect abstraction:  $\alpha(s) := h^*(s)$ .

- + perfect heuristic and usually few abstract states
- usually no succinct encoding for  $\alpha$

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### Automatically Deriving Good Abstraction Heuristics

Abstraction Heuristics for Planning: Main Research Problem Automatically derive effective abstraction heuristics for planning tasks.

we will study three state-of-the-art approaches in the following chapters

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### Summary

- ► Abstraction heuristics from orthogonal abstractions can be added without losing admissibility or consistency.
- ▶ One sufficient condition for orthogonality is that all abstractions are affected by disjoint sets of labels.
- ▶ Practically useful abstractions are those which give informative heuristics, yet have a small representation.
- Coming up with good abstractions automatically is the main research challenge when applying abstraction heuristics in planning.

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