Planning and Optimization

B7. Computational Complexity of Planning: Results

Malte Helmert and Gabriele Röger

Universität Basel

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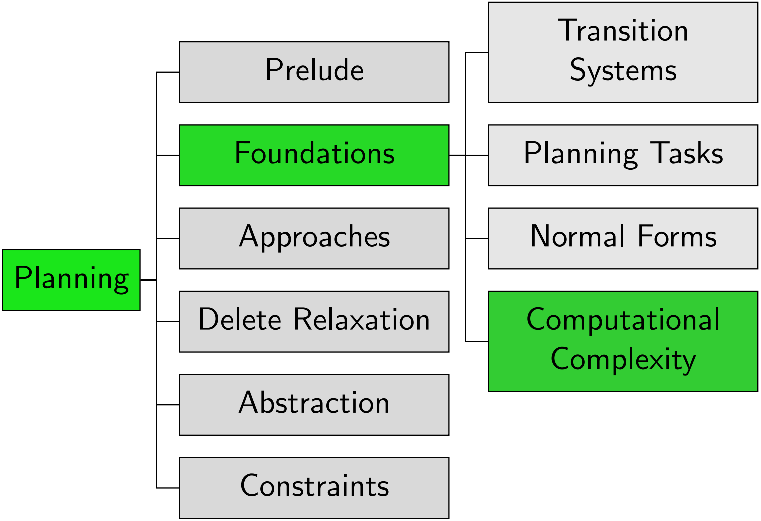
Content of the Course



(Bounded-Cost) Plan Existence

Decision Problems for Planning

Definition (Plan Existence)

Plan existence (PLANEX) is the following decision problem:

GIVEN: planning task Π

QUESTION: Is there a plan for Π?

$\rightsquigarrow$ decision problem analogue of satisficing planning

Definition (Bounded-Cost Plan Existence)

Bounded-cost plan existence (BCPLANEX) is the following decision problem:

GIVEN: planning task Π, cost bound $K \in \mathbb{N}_0$

QUESTION: Is there a plan for Π with cost at most K?

$\rightsquigarrow$ decision problem analogue of optimal planning

Plan Existence vs. Bounded-Cost Plan Existence

Theorem (Reduction from PLANEX to BCPLANEX)

$\textsc{PlanEx} \leq_p \textsc{BCPlanEx}$

Proof.

Consider a planning task Π with state variables V.

Let c_{max} be the maximal cost of all operators of Π.

Compute the number of states of Π as $N = 2^{|V|}$.

Π is solvable iff there is solution with cost at most $c_{max} \cdot (N - 1)$ because a solution need not visit any state twice.

$\rightsquigarrow$ map instance Π of $\textsc{PlanEx}$ to instance $\langle \Pi, c_{max} \cdot (N - 1) \rangle$ of $\textsc{BCPlanEx}$

$\rightsquigarrow$ polynomial reduction □

PSPACE-Completeness of Planning

Membership in PSPACE

Theorem

$\mathrm{BCP{\scriptsize LAN}E{\scriptsize X}} \in$ PSPACE

Proof.

Show $\mathrm{BCP{\scriptsize LAN}E{\scriptsize X}} \in$ NPSPACE and use Savitch's theorem.

Nondeterministic algorithm:

def plan($\langle V, I, O, \gamma \rangle$, K):
  $s := I$
  $k := K$
  loop forever:
    if $s \models \gamma$: **accept**
    guess $o \in O$
    if o is not applicable in s: **fail**
    if $cost(o) > k$: **fail**
    $s := s[\![o]\!]$
    $k := k - cost(o)$

□

PSPACE-Hardness

Idea: generic reduction

- For an arbitrary fixed DTM M with space bound polynomial p and input w, generate propositional planning task which is solvable iff M accepts w in space $p(|w|)$.
- Without loss of generality, we assume $p(n) \geq n$ for all n.

Reduction: State Variables

Let $M = \langle \Sigma, \square, Q, q_0, q_Y, \delta \rangle$ be the fixed DTM, and let p be its space-bound polynomial.

Given input $w_1 \dots w_n$, define relevant tape positions $X := \{-p(n), \dots, p(n)\}$

State Variables

- state_q for all $q \in Q$
- head_i for all $i \in X \cup \{-p(n) - 1, p(n) + 1\}$
- $\text{content}_{i,a}$ for all $i \in X$, $a \in \Sigma_\square$

$\rightsquigarrow$ allows encoding a Turing machine configuration

Reduction: Initial State

Let $M = \langle \Sigma, \square, Q, q_0, q_Y, \delta \rangle$ be the fixed DTM, and let p be its space-bound polynomial.

Given input $w_1 \dots w_n$, define relevant tape positions $X := \{-p(n), \dots, p(n)\}$

Initial State

Initially true:

- state_{q_0}
- head_1
- $\mathrm{content}_{i,w_i}$ for all $i \in \{1, \dots, n\}$
- $\mathrm{content}_{i,\square}$ for all $i \in X \setminus \{1, \dots, n\}$

Initially false:

- all others

Reduction: Operators

Let $M = \langle \Sigma, \Box, Q, q_0, q_Y, \delta \rangle$ be the fixed DTM,
and let p be its space-bound polynomial.

Given input $w_1 \dots w_n$, define relevant tape positions
$X := \{-p(n), \dots, p(n)\}$

Operators

One operator for each transition rule $\delta(q, a) = \langle q', a', d \rangle$
and each cell position $i \in X$:

- precondition: $\mathrm{state}_q \land \mathrm{head}_i \land \mathrm{content}_{i,a}$
- effect: $\neg\mathrm{state}_q \land \neg\mathrm{head}_i \land \neg\mathrm{content}_{i,a}$
  $\land\, \mathrm{state}_{q'} \land \mathrm{head}_{i+d} \land \mathrm{content}_{i,a'}$

Note that add-after-delete semantics are important here!

Reduction: Goal

Let $M = \langle \Sigma, \square, Q, q_0, q_Y, \delta \rangle$ be the fixed DTM, and let p be its space-bound polynomial.

Given input $w_1 \dots w_n$, define relevant tape positions $X := \{-p(n), \dots, p(n)\}$

Goal

state_{q_Y}

PSPACE-Completeness of STRIPS Plan Existence

Theorem (PSPACE-Completeness; Bylander, 1994)

PLANEX *and* BCPLANEX *are* PSPACE*-complete.*
This is true even if only STRIPS tasks are allowed.

Proof.

Membership for BCPLANEX was already shown.

Hardness for PLANEX follows because we just presented a polynomial reduction from an arbitrary problem in PSPACE to PLANEX. (Note that the reduction only generates STRIPS tasks, after trivial cleanup to make them conflict-free.)

Membership for PLANEX and hardness for BCPLANEX follow from the polynomial reduction from PLANEX to BCPLANEX. □

More Complexity Results

More Complexity Results

In addition to the basic complexity result presented in this chapter, there are many special cases, generalizations, variations and related problems studied in the literature:

- different planning formalisms
  - e.g., nondeterministic effects, partial observability, schematic operators, numerical state variables
- syntactic restrictions of planning tasks
  - e.g., without preconditions, without conjunctive effects, STRIPS without delete effects
- semantic restrictions of planning task
  - e.g., restricting variable dependencies ("causal graphs")
- particular planning domains
  - e.g., Blocksworld, Logistics, FreeCell

Complexity Results for Different Planning Formalisms

Some results for different planning formalisms:

- nondeterministic effects:
  - fully observable: EXP-complete (Littman, 1997)
  - unobservable: EXPSPACE-complete (Haslum & Jonsson, 1999)
  - partially observable: 2-EXP-complete (Rintanen, 2004)
- schematic operators:
  - usually adds one exponential level to PlanEx complexity
  - e.g., classical case EXPSPACE-complete (Erol et al., 1995)
- numerical state variables:
  - undecidable in most variations (Helmert, 2002)
  - decidable in restricted setting with at most two numeric state variables (Helal and Lakemeyer, 2025)

Summary

Summary

- Classical planning is PSPACE-complete.
- This is true both for satisficing and optimal planning (rather, the corresponding decision problems).
- The hardness proof is a polynomial reduction that translates an arbitrary polynomial-space DTM into a STRIPS task:
  - DTM configurations are encoded by state variables.
  - Operators simulate transitions between DTM configurations.
  - The DTM accepts an input iff there is a plan for the corresponding STRIPS task.
- This implies that there is no polynomial algorithm for classical planning unless P = PSPACE.
- It also means that planning is not polynomially reducible to any problem in NP unless NP = PSPACE.