

# Planning and Optimization

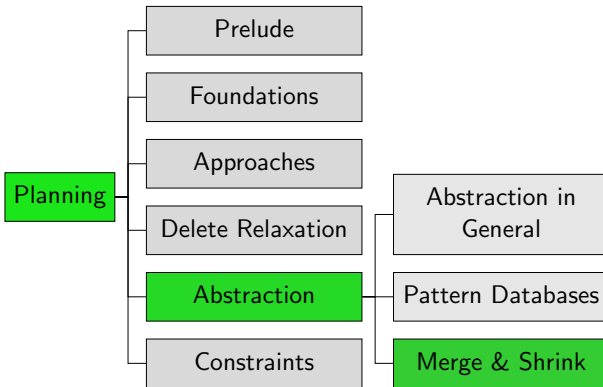
## E9. Merge-and-Shrink: Factored Transition Systems

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# Content of the Course

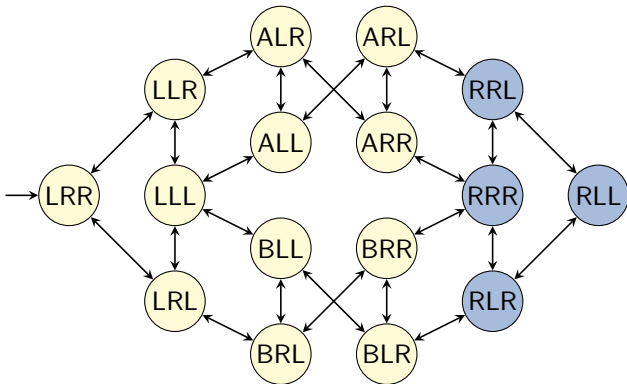


# Motivation

# Beyond Pattern Databases

- Despite their popularity, pattern databases have some **fundamental limitations** ( $\rightsquigarrow$  example on next slides).
- Today and next time, we study a class of abstractions called **merge-and-shrink abstractions**.
- Merge-and-shrink abstractions can be seen as a **proper generalization** of pattern databases.
  - They can do everything that pattern databases can do (modulo polynomial extra effort).
  - They can do some things that pattern databases cannot.

## Back to the Running Example

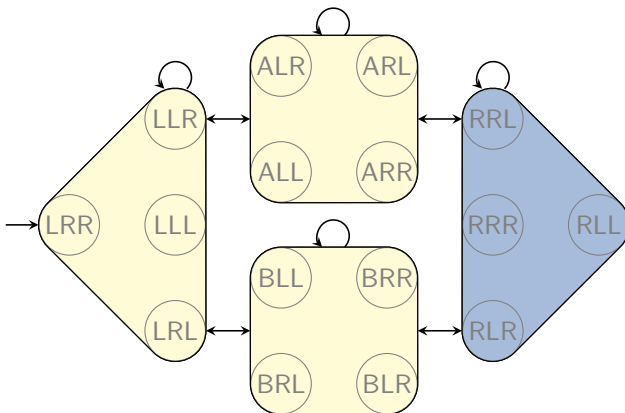


Logistics problem with one package, two trucks, two locations:

- state variable **package**:  $\{L, R, A, B\}$
- state variable **truck A**:  $\{L, R\}$
- state variable **truck B**:  $\{L, R\}$

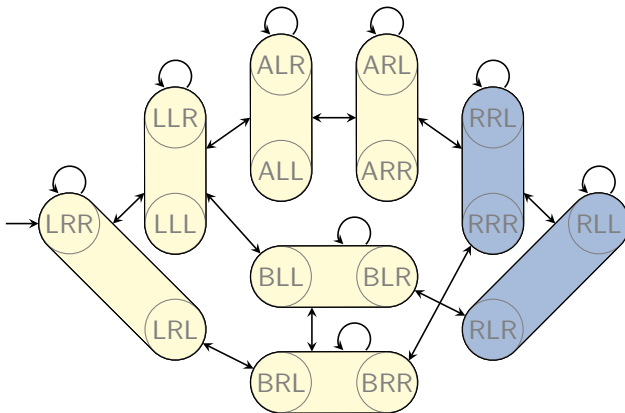
# Example: Projection (1)

$\mathcal{T}^{\pi}\{\text{package}\}$ :



# Example: Projection (2)

$\mathcal{T}^\pi\{\text{package, truck A}\}:$



# Limitations of Projections

How accurate is the PDB heuristic?

- consider **generalization of the example**:  
 $N$  trucks, 1 package
- consider **any** pattern that is a proper subset of variable set  $V$
- $h(s_0) \leq 2 \rightsquigarrow$  **no better** than atomic projection to **package**

These values cannot be improved by maximizing over several patterns or using additive patterns.

**Merge-and-shrink abstractions** can represent heuristics with  $h(s_0) \geq 3$  for tasks of this kind of any size.

Time and space requirements are **linear in  $N$** .

(In fact, with time/space  $O(N^2)$  we can construct a merge-and-shrink abstraction that gives the **perfect heuristic  $h^*$**  for such tasks, but we do not show this here.)



# Main Idea

# Merge-and-Shrink Abstractions: Main Idea

## Main Idea of Merge-and-shrink Abstractions

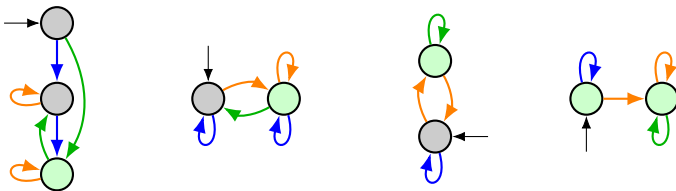
(due to Dräger, Finkbeiner & Podelski, 2006):

Instead of **perfectly** reflecting **a few** state variables, reflect **all** state variables, but in a **potentially lossy** way.

- Represent planning task as **factored transition system** (FTS): a set of (small) abstract transition systems (**factors**) that jointly represent the full transition system of the task.
- Iteratively **transform** FTS by:
  - **merging**: combining two factors into one
  - **shrinking**: reducing the size of a single factor by abstraction
- When only a single factor is left, its goal distances are the merge-and-shrink heuristic values.

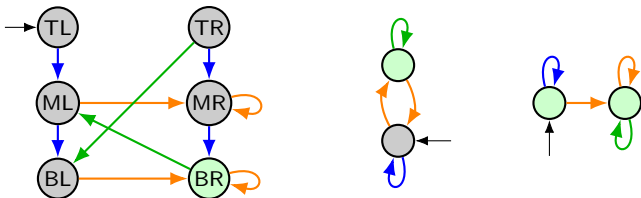
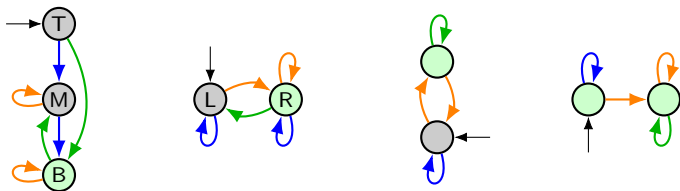
# Merge-and-Shrink Abstractions: Idea

Start from atomic factors (projections to single state variables)



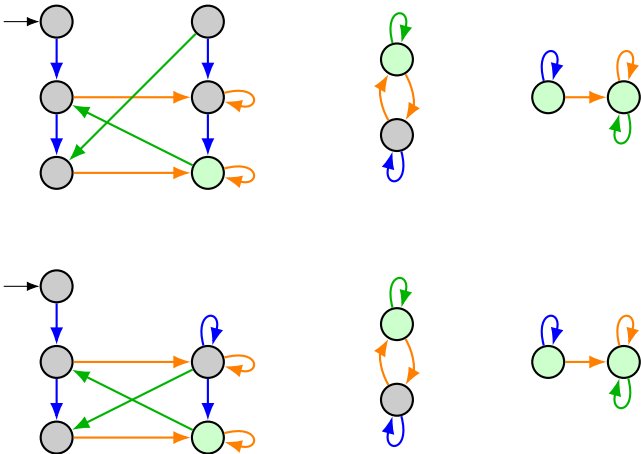
# Merge-and-Shrink Abstractions: Idea

Merge: replace two factors with their product

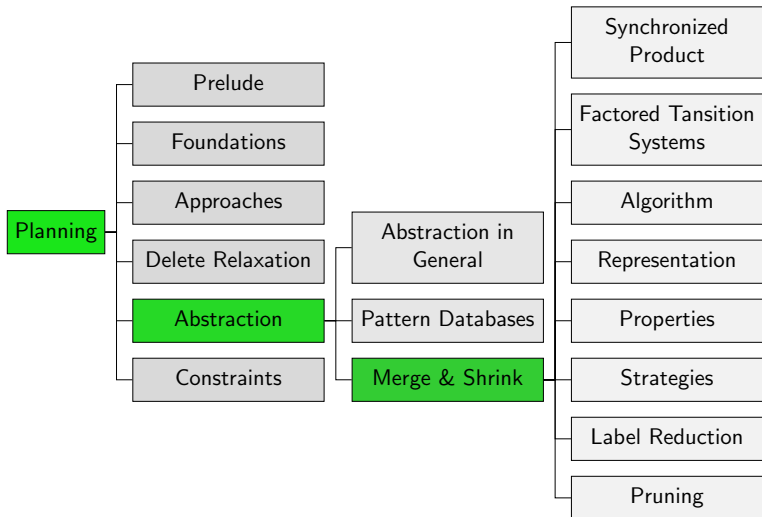


# Merge-and-Shrink Abstractions: Idea

Shrink: replace a factor by an abstraction of it



# Content of the Course



# Atomic Projections

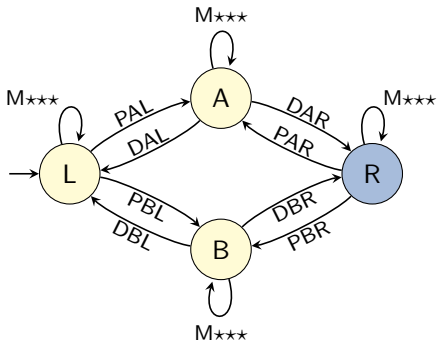
## Running Example: Explanations

- **Atomic projections** (projections to a single state variable) play an important role for merge-and-shrink abstractions.
- Unlike previous chapters, **transition labels** are critically important for merge-and-shrink.
- Hence we now look at the transition systems for atomic projections of our example task, including transition labels.
- We abbreviate labels (operator names) as in these examples:
  - **MALR**: **m**ove truck **A** from **l**eft to **r**ight
  - **DAR**: **d**rop package from truck **A** at **r**ight location
  - **PBL**: **p**ick up package with truck **B** at **l**eft location
- We abbreviate parallel arcs with **commas** and **wildcards** (**\***) as in these examples:
  - **PAL, DAL**: two parallel arcs labeled **PAL** and **DAL**
  - **MA\***: two parallel arcs labeled **MALR** and **MARL**



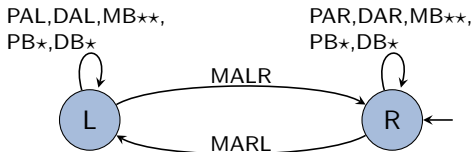
# Running Example: Atomic Projection for Package

$\mathcal{T}^{\pi}\{\text{package}\}$ :



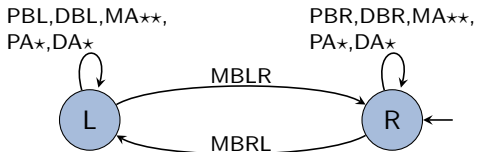
# Running Example: Atomic Projection for Truck A

$\mathcal{T}^\pi\{\text{truck A}\}$ :



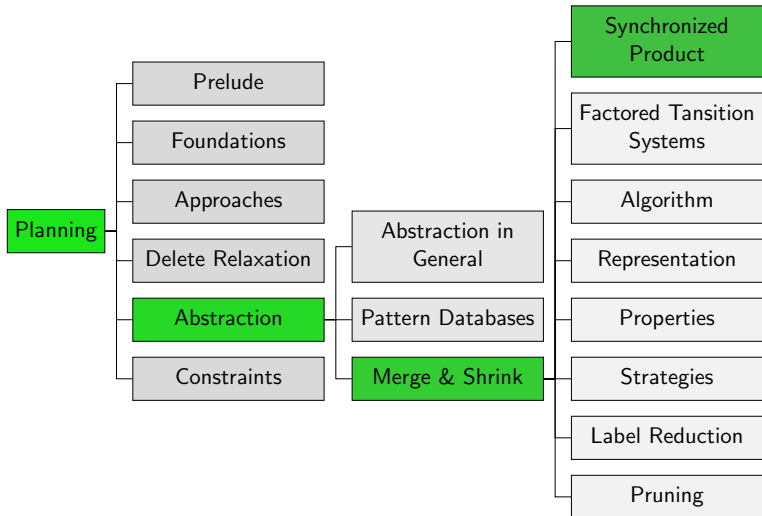
# Running Example: Atomic Projection for Truck B

$\mathcal{T}^\pi\{\text{truck B}\}$ :



# Synchronized Product

# Content of the Course



## Synchronized Product: Idea

- Given two abstract transition systems with the same labels, we can compute a **product transition system**.
- The product transition system **captures all information** of both transition systems.
- A sequence of labels is a solution for the product iff it is a solution for both factors.

# Synchronized Product of Transition Systems

## Definition (Synchronized Product of Transition Systems)

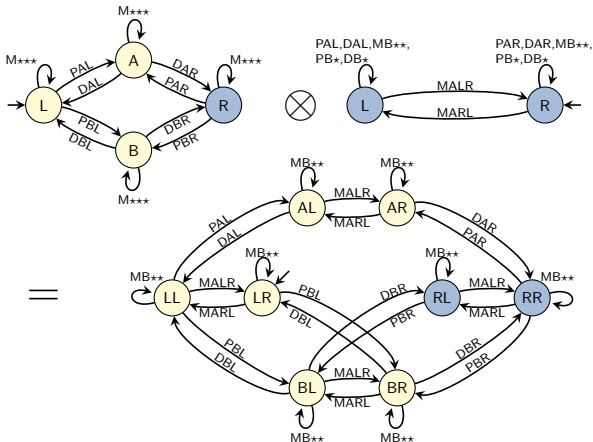
For  $i \in \{1, 2\}$ , let  $\mathcal{T}_i = \langle S_i, L, c, T_i, s_{0i}, S_{*i} \rangle$  be transition systems with the same labels and cost function.

The **synchronized product** of  $\mathcal{T}_1$  and  $\mathcal{T}_2$ , in symbols  $\mathcal{T}_1 \otimes \mathcal{T}_2$ , is the transition system  $\mathcal{T}_\otimes = \langle S_\otimes, L, c, T_\otimes, s_{0\otimes}, S_{*\otimes} \rangle$  with

- $S_\otimes = S_1 \times S_2$
- $T_\otimes = \{ \langle s_1, s_2 \rangle \xrightarrow{\ell} \langle t_1, t_2 \rangle \mid s_1 \xrightarrow{\ell} t_1 \in T_1 \text{ and } s_2 \xrightarrow{\ell} t_2 \in T_2 \}$
- $s_{0\otimes} = \langle s_{01}, s_{02} \rangle$
- $S_{*\otimes} = S_{*1} \times S_{*2}$

# Example: Synchronized Product

$$\mathcal{T}^\pi\{\text{package}\} \otimes \mathcal{T}^\pi\{\text{truck A}\}:$$

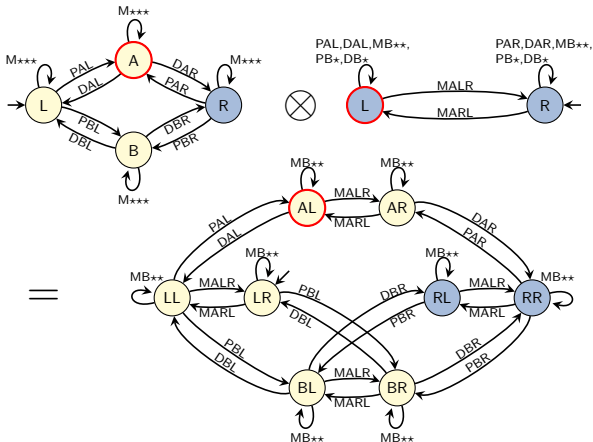




# Example: Synchronized Product

$$\mathcal{T}^\pi\{\text{package}\} \otimes \mathcal{T}^\pi\{\text{truck A}\}:$$

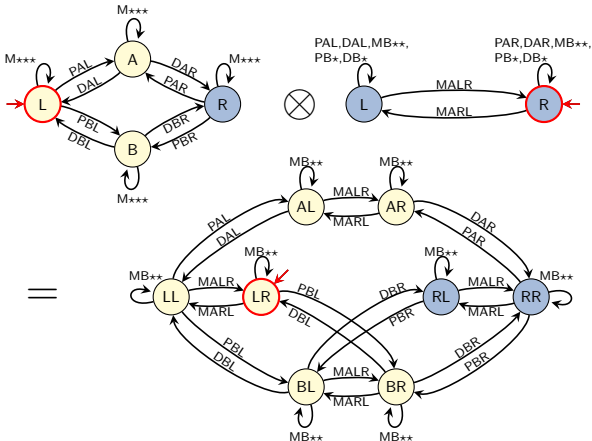
$$S_\otimes = S_1 \times S_2$$



# Example: Synchronized Product

$$\mathcal{T}^\pi\{\text{package}\} \otimes \mathcal{T}^\pi\{\text{truck A}\}:$$

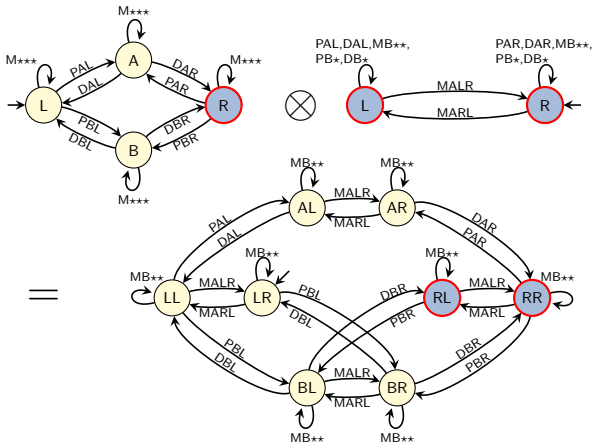
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# Example: Synchronized Product

$$\mathcal{T}^\pi\{\text{package}\} \otimes \mathcal{T}^\pi\{\text{truck A}\}:$$

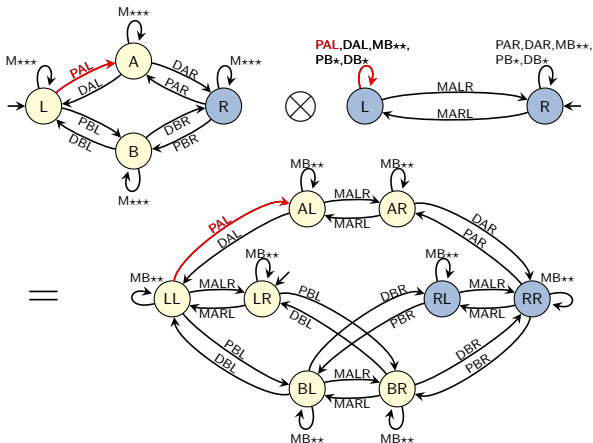
$$S_{*\otimes} = S_{*1} \times S_{*2}$$



# Example: Synchronized Product

$\mathcal{T}^\pi\{\text{package}\} \otimes \mathcal{T}^\pi\{\text{truck A}\}$ :

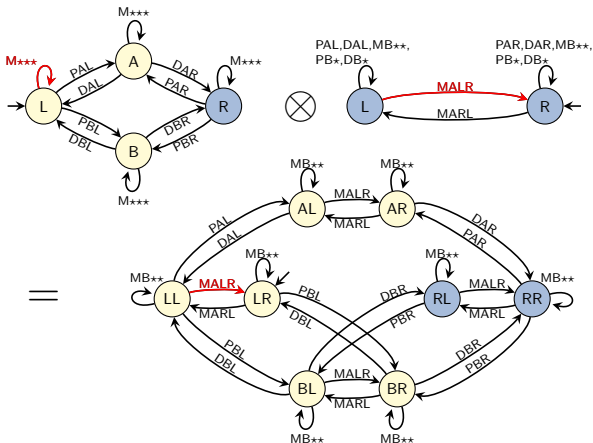
$$T_\otimes = \{ \langle s_1, s_2 \rangle \xrightarrow{\ell} \langle t_1, t_2 \rangle \mid s_1 \xrightarrow{\ell} t_1 \in T_1 \text{ and } s_2 \xrightarrow{\ell} t_2 \in T_2 \}$$



# Example: Synchronized Product

$\mathcal{T}^\pi\{\text{package}\} \otimes \mathcal{T}^\pi\{\text{truck A}\}$ :

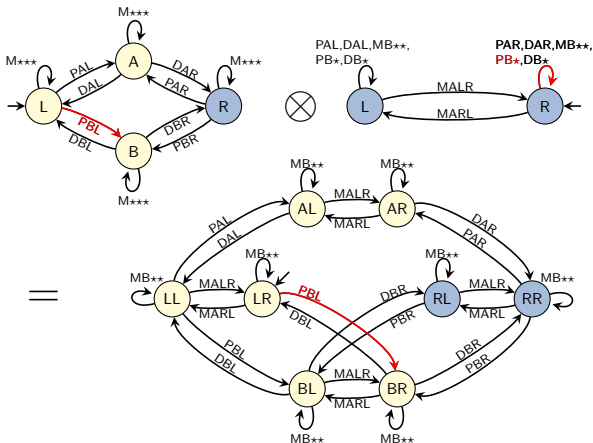
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# Example: Synchronized Product

$\mathcal{T}^\pi\{\text{package}\} \otimes \mathcal{T}^\pi\{\text{truck A}\}$ :

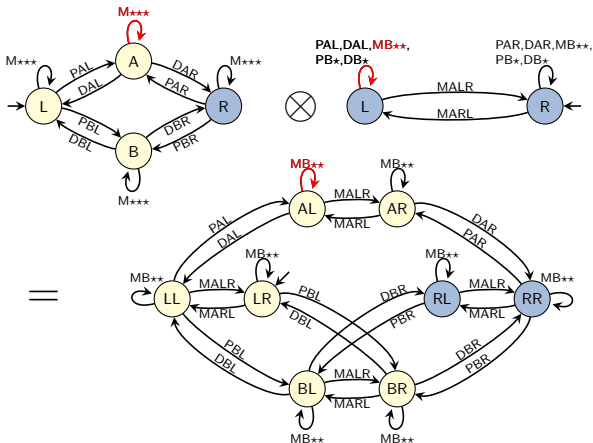
$$T_\otimes = \{\langle s_1, s_2 \rangle \xrightarrow{\ell} \langle t_1, t_2 \rangle \mid s_1 \xrightarrow{\ell} t_1 \in T_1 \text{ and } s_2 \xrightarrow{\ell} t_2 \in T_2\}$$



# Example: Synchronized Product

$\mathcal{T}^\pi\{\text{package}\} \otimes \mathcal{T}^\pi\{\text{truck A}\}$ :

$$T_\otimes = \{\langle s_1, s_2 \rangle \xrightarrow{\ell} \langle t_1, t_2 \rangle \mid s_1 \xrightarrow{\ell} t_1 \in T_1 \text{ and } s_2 \xrightarrow{\ell} t_2 \in T_2\}$$



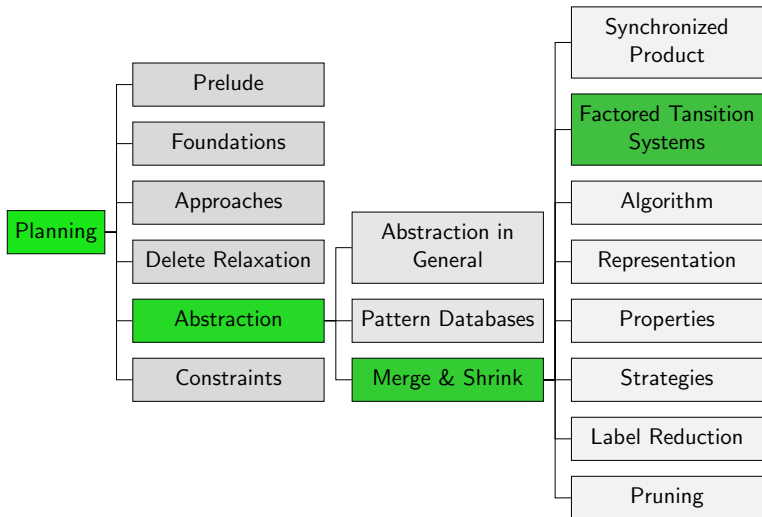
# Associativity and Commutativity

- Up to isomorphism (“names of states”), products are associative and commutative:
  - $(\mathcal{T} \otimes \mathcal{T}') \otimes \mathcal{T}'' \sim \mathcal{T} \otimes (\mathcal{T}' \otimes \mathcal{T}'')$
  - $\mathcal{T} \otimes \mathcal{T}' \sim \mathcal{T}' \otimes \mathcal{T}$
- We do not care about names of states and thus treat products as associative and commutative.
- We can then define the product of a **set**  $F = \{\mathcal{T}_1, \dots, \mathcal{T}_n\}$  of transition systems:  $\bigotimes F := \mathcal{T}_1 \otimes \dots \otimes \mathcal{T}_n$



# Factored Transition Systems

# Content of the Course



# Factored Transition System

## Definition (Factored Transition System)

A finite set  $F = \{\mathcal{T}_1, \dots, \mathcal{T}_n\}$  of transition systems with the same labels and cost function is called a **factored transition system (FTS)**.

$F$  **represents** the transition system  $\otimes F$ .

A planning task gives rise to an FTS via its atomic projections:

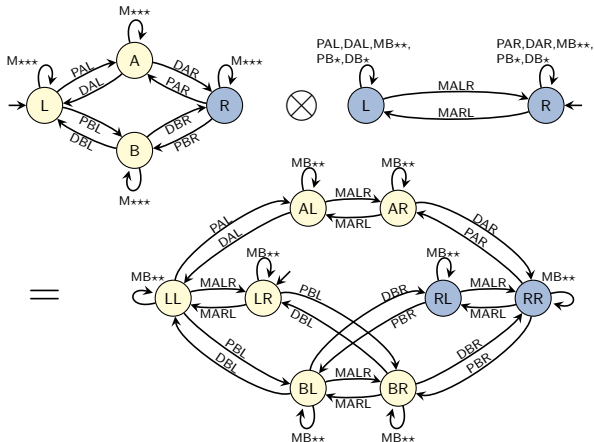
## Definition (Factored Transition System Induced by Planning Task)

Let  $\Pi$  be a planning task with state variables  $V$ .

The **factored transition system induced by  $\Pi$**  is the FTS  $F(\Pi) = \{\mathcal{T}^{\pi\{v\}} \mid v \in V\}$ .

# Back to the Example Product

$$\mathcal{T}^\pi\{\text{package}\} \otimes \mathcal{T}^\pi\{\text{truck A}\}:$$



We have  $\mathcal{T}^\pi\{\text{package}\} \otimes \mathcal{T}^\pi\{\text{truck A}\} \sim \mathcal{T}^\pi\{\text{package, truck A}\}$ . Coincidence?

# Products of Projections

## Theorem (Products of Projections)

Let  $\Pi$  be a **SAS<sup>+</sup>** planning task with variable set  $V$ ,  
and let  $V_1$  and  $V_2$  be **disjoint subsets** of  $V$ .

Then  $\mathcal{T}^{\pi_{V_1}} \otimes \mathcal{T}^{\pi_{V_2}} \sim \mathcal{T}^{\pi_{V_1 \cup V_2}}$ .

(Proof omitted.)

↪ products allow us to build finer projections from coarser ones

# Recovering $\mathcal{T}(\Pi)$ from the Factored Transition System

- By repeated application of the theorem, we can recover **all pattern database heuristics** of a  $SAS^+$  planning task as products of atomic factors.
- Moreover, by computing the product of **all** atomic projections, we can recover the **identity abstraction**  $id = \pi_V$ .

This implies:

Corollary (Recovering  $\mathcal{T}(\Pi)$  from the Factored Transition System)

Let  $\Pi$  be a  $SAS^+$  planning task. Then  $\bigotimes F(\Pi) \sim \mathcal{T}(\Pi)$ .

This is an important result because it shows that  $F(\Pi)$  **represents all important information** about  $\Pi$ .

# Summary

# Summary

- A **factored transition system** is a set of transition systems that represents a larger transition system by focusing on its individual components (**factors**).
- For planning tasks, these factors are the **atomic projections** (projections to single state variables).
- The **synchronized product**  $\mathcal{T} \otimes \mathcal{T}'$  of two transition systems with the same labels captures their “joint behaviour”.
- For SAS<sup>+</sup> tasks, all **projections** can be obtained as products of atomic projections.
- In particular, the product of all factors of a SAS<sup>+</sup> task results in the **full** transition system of the task.