

# Planning and Optimization

## C5. SAT Planning: Parallel Encoding

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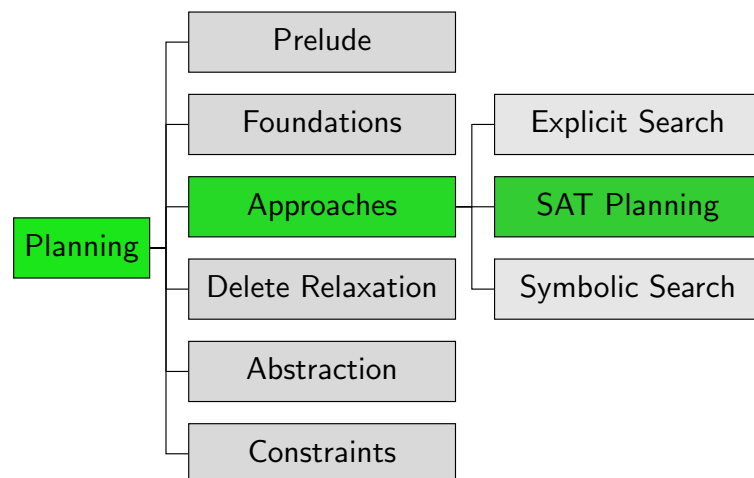
October 14, 2024 — C5. SAT Planning: Parallel Encoding

C5.1 Introduction

C5.2 Adapting the SAT Encoding

C5.3 Summary

## Content of the Course



## C5.1 Introduction

## Efficiency of SAT Planning

- ▶ All other things being equal, the most important aspect for efficient SAT solving is the **number of propositional variables** in the input formula.
- ▶ For sufficiently difficult inputs, runtime scales **exponentially** in the number of variables.
- ↪ Can we make SAT planning more efficient by **using fewer variables**?

## Number of Variables

### Reminder:

- ▶ given propositional planning task  $\Pi = \langle V, I, O, \gamma \rangle$
- ▶ given **horizon**  $T \in \mathbb{N}_0$

### Variables of the SAT Formula

- ▶ propositional variables  $v^i$  for all  $v \in V, 0 \leq i \leq T$   
encode state after  $i$  steps of the plan
- ▶ propositional variables  $o^i$  for all  $o \in O, 1 \leq i \leq T$   
encode operator(s) applied in  $i$ -th step of the plan

↪  $|V| \cdot (T + 1) + |O| \cdot T$  variables

↪ SAT solving runtime usually **exponential in  $T$**

## Parallel Plans and Interference

Can we get away with shorter horizons?

### Idea:

- ▶ allow **parallel plans** in the SAT encoding:  
**multiple operators** can be applied in the same step  
if they do not **interfere**

### Definition (Interference)

Let  $O' = \{o_1, \dots, o_n\}$  be a set of operators applicable in state  $s$ .

We say that  $O'$  is **interference-free** in  $s$  if

- ▶ for all permutations  $\pi$  of  $O'$ ,  $s[\pi]$  is defined, and
- ▶ for all permutations  $\pi, \pi'$  of  $O'$ ,  $s[\pi] = s[\pi']$ .

We say that  $O'$  **interfere** in  $s$  if they are not interference-free in  $s$ .

## Parallel Plan Extraction

- ▶ If we can rule out interference, we can allow multiple operators at the same time in the SAT encoding.
- ▶ A parallel plan (with multiple  $o^i$  used for the same  $i$ ) extracted from the SAT formula can then be converted into a “regular” plan by ordering the operators within each time step arbitrarily.

## Challenges for Parallel SAT Encodings

Two challenges remain:

- ▶ our current SAT encoding **does not allow concurrent operators**
- ▶ how do we ensure that our plans are **interference-free**?

## C5.2 Adapting the SAT Encoding

## Reminder: Sequential SAT Encoding (1)

### Sequential SAT Formula (1)

initial state clauses:

- ▶  $v^0$  for all  $v \in V$  with  $I(v) = \mathbf{T}$
- ▶  $\neg v^0$  for all  $v \in V$  with  $I(v) = \mathbf{F}$

goal clauses:

- ▶  $\gamma^T$

operator selection clauses:

- ▶  $o_1^i \vee \dots \vee o_n^i$  for all  $1 \leq i \leq T$

operator exclusion clauses:

- ▶  $\neg o_j^i \vee \neg o_k^i$  for all  $1 \leq i \leq T, 1 \leq j < k \leq n$

↔ operator exclusion clauses must be adapted

## Sequential SAT Encoding (2)

### Sequential SAT Formula (2)

precondition clauses:

- ▶  $o^i \rightarrow pre(o)^{i-1}$  for all  $1 \leq i \leq T, o \in O$

positive and negative effect clauses:

- ▶  $(o^i \wedge \alpha^{i-1}) \rightarrow v^i$  for all  $1 \leq i \leq T, o \in O, v \in V$
- ▶  $(o^i \wedge \delta^{i-1} \wedge \neg \alpha^{i-1}) \rightarrow \neg v^i$  for all  $1 \leq i \leq T, o \in O, v \in V$

positive and negative frame clauses:

- ▶  $(o^i \wedge v^{i-1} \wedge \neg v^i) \rightarrow \delta^{i-1}$  for all  $1 \leq i \leq T, o \in O, v \in V$
- ▶  $(o^i \wedge \neg v^{i-1} \wedge v^i) \rightarrow \alpha^{i-1}$  for all  $1 \leq i \leq T, o \in O, v \in V$

where  $\alpha = effcond(v, eff(o)), \delta = effcond(\neg v, eff(o))$ .

↔ frame clauses must be adapted

## Adapting the Operator Exclusion Clauses: Idea

Reminder: operator exclusion clauses  $\neg o_j^i \vee \neg o_k^i$   
for all  $1 \leq i \leq T, 1 \leq j < k \leq n$

- ▶ **Ideally:** replace with clauses that express “for all states  $s$ , the operators selected at time  $i$  are **interference-free** in  $s$ ”
- ▶ **but:** testing if a given set of operators interferes in any state is itself an NP-complete problem
- ↪ use something less heavy: a **sufficient condition** for interference-freeness that can be expressed at the level of **pairs** of operators

## Conflicting Operators

- ▶ Intuitively, two operators **conflict** if
  - ▶ one can disable the precondition of the other,
  - ▶ one can override an effect of the other, or
  - ▶ one can enable or disable an effect condition of the other.
- ▶ If no two operators in a set  $O'$  conflict, then  $O'$  is interference-free in all states.
- ▶ This is still difficult to test, so we restrict attention to the **STRIPS** case in the following.

### Definition (Conflicting STRIPS Operator)

Operators  $o$  and  $o'$  of a STRIPS task  $\Pi$  **conflict** if

- ▶  $o$  deletes a precondition of  $o'$  or vice versa, or
- ▶  $o$  deletes an add effect of  $o'$  or vice versa.

## Adapting the Operator Exclusion Clauses: Solution

Reminder: operator exclusion clauses  $\neg o_j^i \vee \neg o_k^i$   
for all  $1 \leq i \leq T, 1 \leq j < k \leq n$

### Solution:

Parallel SAT Formula: Operator Exclusion Clauses  
operator exclusion clauses:

- ▶  $\neg o_j^i \vee \neg o_k^i$  for all  $1 \leq i \leq T, 1 \leq j < k \leq n$   
such that  $o_j$  and  $o_k$  conflict

## Adapting the Frame Clauses: Idea

Reminder: frame clauses

$$\begin{aligned} (o^i \wedge v^{i-1} \wedge \neg v^i) &\rightarrow \delta^{i-1} && \text{for all } 1 \leq i \leq T, o \in O, v \in V \\ (o^i \wedge \neg v^{i-1} \wedge v^i) &\rightarrow \alpha^{i-1} && \text{for all } 1 \leq i \leq T, o \in O, v \in V \end{aligned}$$

### What is the problem?

- ▶ These clauses express that if  $o$  is applied at time  $i$  and the value of  $v$  changes, then  $o$  **caused the change**.
- ▶ This is no longer true if we want to be able to apply two operators concurrently.
- ↪ Instead, say “If the value of  $v$  changes, then **some operator** must have caused the change.”

## Adapting the Frame Clauses: Solution

Reminder: frame clauses

$$(o^i \wedge v^{i-1} \wedge \neg v^i) \rightarrow \delta^{i-1} \quad \text{for all } 1 \leq i \leq T, o \in O, v \in V$$

$$(o^i \wedge \neg v^{i-1} \wedge v^i) \rightarrow \alpha^{i-1} \quad \text{for all } 1 \leq i \leq T, o \in O, v \in V$$

**Solution:**

Parallel SAT Formula: Frame Clauses

positive and negative frame clauses:

$$\blacktriangleright (v^{i-1} \wedge \neg v^i) \rightarrow ((o_1^i \wedge \delta_{o_1}^{i-1}) \vee \dots \vee (o_n^i \wedge \delta_{o_n}^{i-1}))$$

for all  $1 \leq i \leq T, v \in V$

$$\blacktriangleright (\neg v^{i-1} \wedge v^i) \rightarrow ((o_1^i \wedge \alpha_{o_1}^{i-1}) \vee \dots \vee (o_n^i \wedge \alpha_{o_n}^{i-1}))$$

for all  $1 \leq i \leq T, v \in V$

where  $\alpha_o = \text{effcond}(v, \text{eff}(o))$ ,  $\delta_o = \text{effcond}(\neg v, \text{eff}(o))$ ,  
 $O = \{o_1, \dots, o_n\}$ .

For STRIPS, these are in clause form.

## C5.3 Summary

## Summary

- ▶ As a rule of thumb, SAT solvers generally perform better on formulas with fewer variables.
- ▶ **Parallel encodings** reduce the number of variables by shortening the horizon needed to solve a planning task.
- ▶ Parallel encodings replace the constraint that operators are not applied concurrently by the constraint that **conflicting** operators are not applied concurrently.
- ▶ To make parallelism possible, the **frame clauses** also need to be adapted.