

# Discrete Mathematics in Computer Science

## B4. Operations on Relations

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# Relations: Recap

- ▶ A **relation over sets**  $S_1, \dots, S_n$  is a set  $R \subseteq S_1 \times \dots \times S_n$ .
- ▶ A **binary** relation is a relation over two sets.
- ▶ A **homogeneous** relation  $R$  over set  $S$  is a binary relation  $R \subseteq S \times S$ .

# Set Operations

- ▶ Relations are **sets** of tuples, so we can build their union, intersection, complement, . . . .
- ▶ Let  $R$  be a relation over  $S_1, \dots, S_n$  and  $R'$  a relation over  $S'_1, \dots, S'_n$ . Then  $R \cup R'$  is a relation over  $S_1 \cup S'_1, \dots, S_n \cup S'_n$ .  
With the standard relations  $<, =$  and  $\leq$  for  $\mathbb{N}_0$ ,  
relation  $\leq$  corresponds to the union of relations  $<$  and  $=$ .
- ▶ Let  $R$  and  $R'$  be relations over  $n$  sets.  
Then  $R \cap R'$  is a relation.  
**Over which sets?**  
With the standard relations  $\leq, =$  and  $\geq$  for  $\mathbb{N}_0$ ,  
relation  $=$  corresponds to the intersection of  $\leq$  and  $\geq$ .
- ▶ If  $R$  is a relation over  $S_1, \dots, S_n$   
then so is the **complementary relation**  $\bar{R} = (S_1 \times \dots \times S_n) \setminus R$ .  
With the standard relations for  $\mathbb{N}_0$ , relation  $=$  is the  
complementary relation of  $\neq$  and  $>$  the one of  $\leq$ .

# Inverse of a Relation

## Definition

Let  $R \subseteq A \times B$  be a binary relation over  $A$  and  $B$ .

The **inverse relation** of  $R$  is the relation  $R^{-1} \subseteq B \times A$  given by  $R^{-1} = \{(b, a) \mid (a, b) \in R\}$ .

- ▶ The inverse of the  $<$  relation over  $\mathbb{N}_0$  is the  $>$  relation.
- ▶ Relation  $R$  with  $xRy$  iff person  $x$  has a key for  $y$ .  
Inverse:  $Q$  with  $aQb$  iff lock  $a$  can be opened by person  $b$ .

German: inverse Relation oder Umkehrrelation

# Composition of Relations

## Definition (Composition of relations)

Let  $R_1$  be a relation over  $A$  and  $B$  and  $R_2$  a relation over  $B$  and  $C$ . The **composition of  $R_1$  and  $R_2$**  is the relation  $R_2 \circ R_1$  over  $A$  and  $C$  with:

$$R_2 \circ R_1 = \{(a, c) \mid \text{there is a } b \in B \text{ with} \\ (a, b) \in R_1 \text{ and } (b, c) \in R_2\}$$

How can we illustrate this graphically?

German: Komposition oder Rückwärtsverkettung

# Composition of Relations: Example

$$S_1 = \{1, 2, 3, 4\}$$

$$S_2 = \{A, B, C, D, E\}$$

$$S_3 = \{a, b, c, d\}$$

$$R_1 = \{(1, A), (1, B), (3, B), (4, D)\} \text{ over } S_1 \text{ and } S_2$$

$$R_2 = \{(B, a), (C, c), (D, a), (D, d)\} \text{ over } S_2 \text{ and } S_3$$

$$R_2 \circ R_1 =$$

# Composition is Associative

## Theorem (Associativity of composition)

Let  $S_1, \dots, S_4$  be sets and  $R_1, R_2, R_3$  relations with  $R_i \subseteq S_i \times S_{i+1}$ .  
Then

$$R_3 \circ (R_2 \circ R_1) = (R_3 \circ R_2) \circ R_1.$$

## Proof.

It holds that  $(x_1, x_4) \in R_3 \circ (R_2 \circ R_1)$  iff there is an  $x_3$  with  $(x_1, x_3) \in R_2 \circ R_1$  and  $(x_3, x_4) \in R_3$ .

As  $(x_1, x_3) \in R_2 \circ R_1$  iff there is an  $x_2$  with  $(x_1, x_2) \in R_1$  and  $(x_2, x_3) \in R_2$ , we have overall that  $(x_1, x_4) \in R_3 \circ (R_2 \circ R_1)$  iff there are  $x_2, x_3$  with  $(x_1, x_2) \in R_1$ ,  $(x_2, x_3) \in R_2$  and  $(x_3, x_4) \in R_3$ .

This is the case iff there is an  $x_2$  with  $(x_1, x_2) \in R_1$  and  $(x_2, x_4) \in R_3 \circ R_2$ , which holds iff  $(x_1, x_4) \in (R_3 \circ R_2) \circ R_1$ . □

## (Reflexive) Transitive Closure

### Definition ((Reflexive) transitive closure)

Let  $R$  be a relation over set  $S$ .

The **transitive closure**  $R^+$  of  $R$  is the **smallest relation over  $S$  that is transitive and has  $R$  as a subset.**

The **reflexive transitive closure**  $R^*$  of  $R$  is the **smallest relation over  $S$  that is reflexive, transitive and has  $R$  as a subset.**

The (reflexive) transitive closure always exists. **Why?**

**Example:** If  $aRb$  specifies that there is a direct flight from  $a$  to  $b$ , what do  $R^+$  and  $R^*$  express?

German: (reflexive) transitive Hülle

# Transitive Closure and $n$ -fold Composition

Define the  $n$ -fold composition of a relation  $R$  over  $S$  as

$$\begin{aligned} R_0 &= \{(x, x) \mid x \in S\} && \text{and} \\ R_i &= R \circ R_{i-1} && \text{for } i > 1. \end{aligned}$$

## Theorem

Let  $R$  be a relation over set  $S$ .

Then  $R^+ = \bigcup_{i=1}^{\infty} R_i$  and  $R^* = \bigcup_{i=0}^{\infty} R_i$ .

Without proof.

German:  $n$ -fache Komposition

# Other Operators

- ▶ There are many more operators, also for general relations.
- ▶ Highly relevant for **queries over relational databases**.
- ▶ For example, **join** operators combine relations based on common entries.
- ▶ Example for a **natural join**:

*Employee*

Name	Empld	DeptName
Harry	3415	Finance
Sally	2241	Sales
George	3401	Finance
Harriet	2202	Sales
Mary	1257	Human Resources

*Dept*

DeptName	Manager
Finance	George
Sales	Harriet
Production	Charles

*Employee ⋈ Dept*

Name	Empld	DeptName	Manager
Harry	3415	Finance	George
Sally	2241	Sales	Harriet
George	3401	Finance	George
Harriet	2202	Sales	Harriet

(Source: Wikipedia)

# Summary

- ▶ Relations: general, binary, homogeneous
- ▶ Properties: reflexivity, symmetry, transitivity (and related properties)
- ▶ Special relations: equivalence relations, order relations
- ▶ Operations: inverse, composition, transitive closure