

Solution of Exercise in A5

Gabi Röger

1 Basic definitions (from the slides)

The set \mathcal{B} of binary trees is inductively defined as follows:

- \square is a binary tree (a *leaf*)
- If L and R are binary trees, then $\langle L, \circ, R \rangle$ is a binary tree (with *inner node* \circ).

The number of *leaves* of a binary tree B , written $leaves(B)$, is defined as follows:

$$\begin{aligned}leaves(\square) &= 1 \\leaves(\langle L, \circ, R \rangle) &= leaves(L) + leaves(R)\end{aligned}$$

The *height* of a binary tree B , written $height(B)$, is defined as follows:

$$\begin{aligned}height(\square) &= 0 \\height(\langle L, \circ, R \rangle) &= \max\{height(L), height(R)\} + 1\end{aligned}$$

2 Example: Structural Induction

We will establish the following theorem:

Theorem 1. For all binary trees B : $leaves(B) \leq 2^{height(B)}$.

Proof. We prove the theorem by structural induction.

Induction basis: If $B = \square$, we have $leaves(\square) = 1 = 2^0 = 2^{height(\square)}$. ✓

Induction hypothesis: Suppose that for binary trees L and R it holds that $leaves(L) \leq 2^{height(L)}$ and $leaves(R) \leq 2^{height(R)}$.

Inductive step: Consider $B = \langle L, \circ, R \rangle$.

By the definition of *leaves*, it holds that $leaves(B) = leaves(L) + leaves(R)$. Using the induction hypothesis, we get

$$\begin{aligned}leaves(B) &= leaves(L) + leaves(R) \\ &\stackrel{\text{IH}}{\leq} 2^{\text{height}(L)} + 2^{\text{height}(R)} \\ &\stackrel{(*)}{\leq} 2^{\max\{\text{height}(L), \text{height}(R)\}} + 2^{\max\{\text{height}(L), \text{height}(R)\}} \\ &= 2 \cdot 2^{\max\{\text{height}(L), \text{height}(R)\}} \\ &\stackrel{(**)}{=} 2 \cdot 2^{\text{height}(B)-1} \\ &= 2^{\text{height}(B)}\end{aligned}$$

For inequality (*), we use that $\text{height}(L) \leq \max\{\text{height}(L), \text{height}(R)\}$ and $\text{height}(R) \leq \max\{\text{height}(L), \text{height}(R)\}$. For equality (**), we use the definition of *height*. \square