

# Planning and Optimization

## D7. Delete Relaxation: Analysis of $h^{\max}$ and $h^{\text{add}}$

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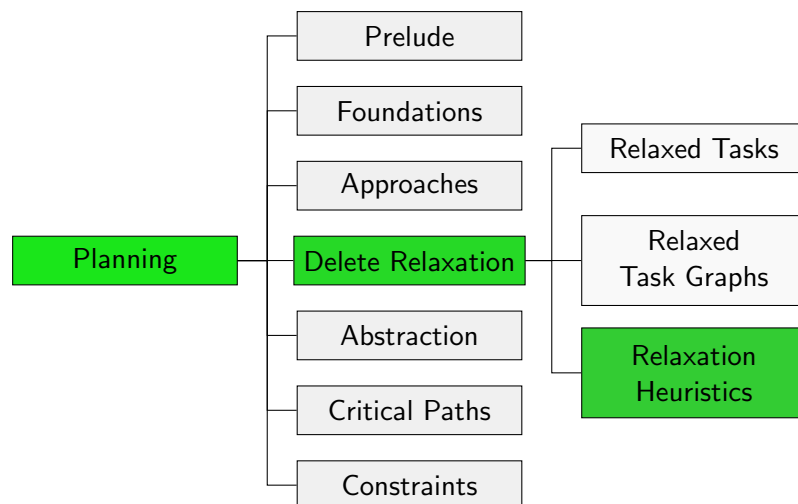
October 30, 2023 — D7. Delete Relaxation: Analysis of  $h^{\max}$  and  $h^{\text{add}}$

D7.1 Choice Functions

D7.2 Best Achievers

D7.3 Summary

## Content of this Course



D7. Delete Relaxation: Analysis of  $h^{\max}$  and  $h^{\text{add}}$

Choice Functions

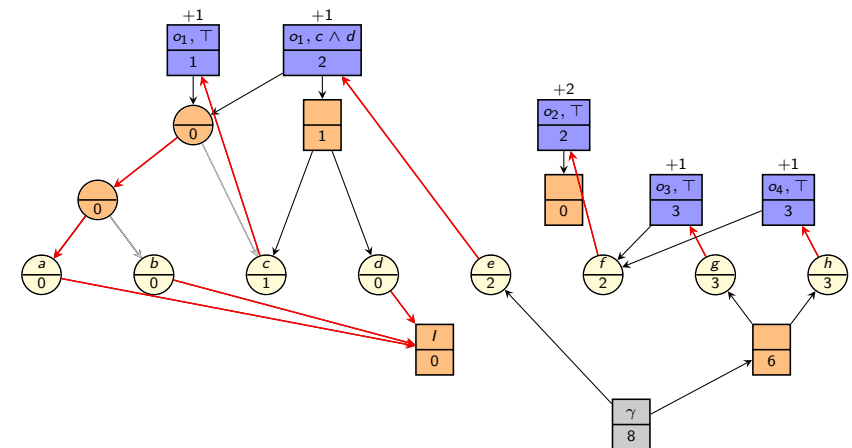
## D7.1 Choice Functions

## Motivation

- ▶ In this chapter, we analyze the behaviour of  $h^{\max}$  and  $h^{\text{add}}$  more deeply.
- ▶ Our goal is to understand their shortcomings.
  - ▶ In the next chapter we then used this understanding to devise an improved heuristic.
- ▶ As a preparation for our analysis, we need some further definitions that concern **choices** in AND/OR graphs.
- ▶ The key observation is that if we want to establish the value of a certain node  $n$ , we can to some extent **choose** how we want to achieve the OR nodes that are relevant to achieving  $n$ .

## Preview: Choice Function & Best Achievers

Preserve at most one outgoing arc of each OR node, but node values may not change.



## Choice Functions

### Definition (Choice Function)

Let  $G$  be an AND/OR graph with nodes  $N$  and OR nodes  $N_{\vee}$ .

A **choice function** for  $G$  is a function  $f : N' \rightarrow N$  defined on some set  $N' \subseteq N_{\vee}$  such that  $f(n) \in \text{succ}(n)$  for all  $n \in N'$ .

- ▶ In words, choice functions select (at most) **one** successor for each OR node of  $G$ .
- ▶ Intuitively,  $f(n)$  selects by which disjunct  $n$  is achieved.
- ▶ If  $f(n)$  is undefined for a given  $n$ , the intuition is that  $n$  is not achieved.

## Reduced Graphs

Once we have decided how to achieve an OR node, we can remove the other alternatives:

### Definition (Reduced Graph)

Let  $G$  be an AND/OR graph, and let  $f$  be a choice function for  $G$  defined on nodes  $N'$ .

The **reduced graph** for  $f$  is the subgraph of  $G$  where all outgoing arcs of OR nodes are removed except for the chosen arcs  $\langle n, f(n) \rangle$  with  $n \in N'$ .

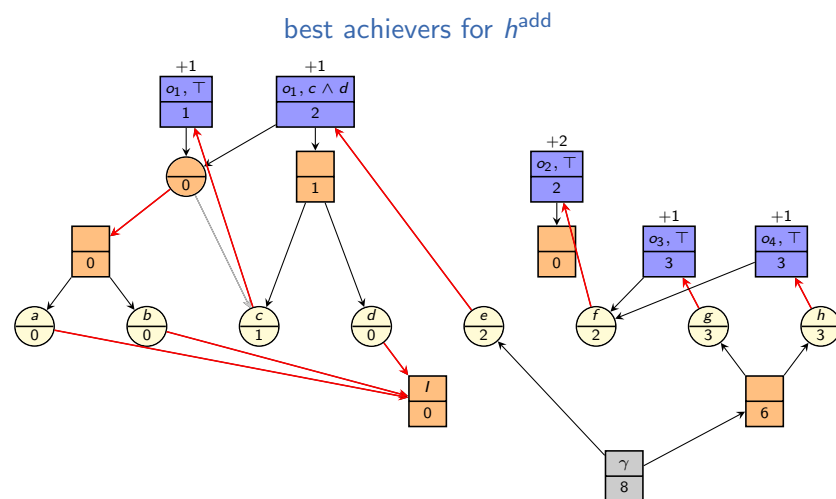
## D7.2 Best Achievers

## Choice Functions Induced by $h^{\max}$ and $h^{\text{add}}$

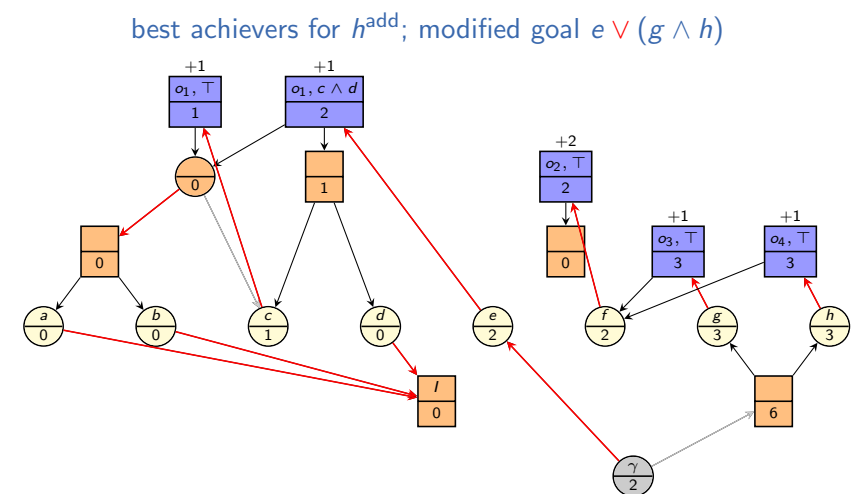
Which choices do  $h^{\max}$  and  $h^{\text{add}}$  make?

- ▶ At every OR node  $n$ , we set the cost of  $n$  to the **minimum** of the costs of the successors of  $n$ .
- ▶ The motivation for this is to achieve  $n$  via the successor that can be achieved **most cheaply** according to our cost estimates.
- ↪ This corresponds to defining a choice function  $f$  with  $f(n) \in \arg \min_{n' \in N'} n'.cost$  for all reached OR nodes  $n$ , where  $N' \subseteq succ(n)$  are all successors of  $n$  processed before  $n$ .
- ▶ The successors chosen by this cost function are called **best achievers** (according to  $h^{\max}$  or  $h^{\text{add}}$ ).
- ▶ Note that the best achiever function  $f$  is in general not well-defined because there can be multiple minimizers. We assume that ties are broken arbitrarily.

### Example: Best Achievers (1)



### Example: Best Achievers (2)



## Best Achiever Graphs

- ▶ **Observation:** The  $h^{\max}/h^{\text{add}}$  costs of nodes remain the same if we replace the RTG by the reduced graph for the respective best achiever function.
- ▶ The AND/OR graph that is obtained by removing all nodes with infinite cost from this reduced graph is called the **best achiever graph** for  $h^{\max}/h^{\text{add}}$ .
  - ▶ We write  $G^{\max}$  and  $G^{\text{add}}$  for the best achiever graphs.
- ▶  $G^{\max}$  ( $G^{\text{add}}$ ) is always **acyclic**: for all arcs  $\langle n, n' \rangle$  it contains,  $n$  is processed by  $h^{\max}$  (by  $h^{\text{add}}$ ) after  $n'$ .

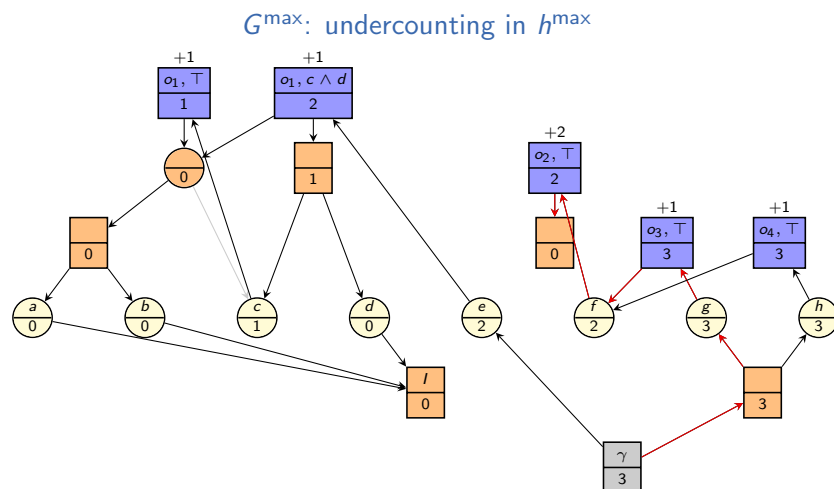
## Paths in Best Achiever Graphs

Let  $n$  be a node of the best achiever graph.  
 Let  $N_{\text{eff}}$  be the set of effect nodes of the best achiever graph.  
 The **cost** of an **effect node** is the cost of the associated operator.  
 The **cost** of a **path** in the best achiever graph is the sum of costs of all **effect nodes** on the path.

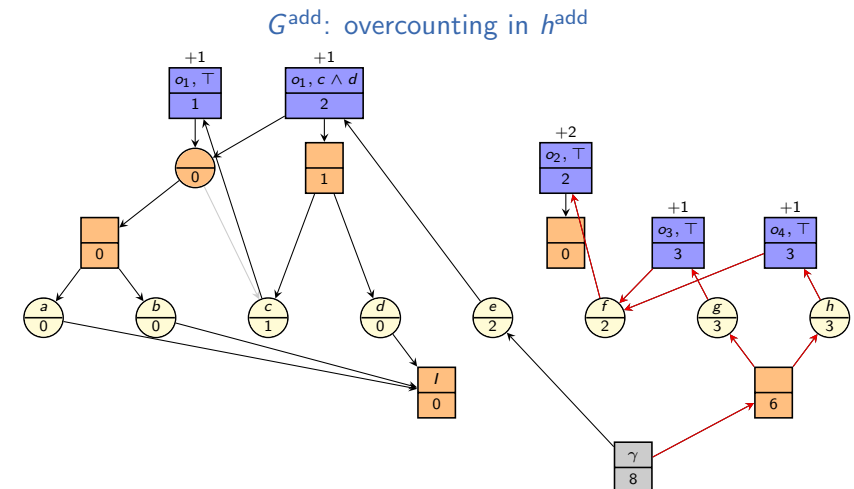
The following properties can be shown by induction:

- ▶  $h^{\max}(n)$  is the **maximum cost** of all paths originating from  $n$  in  $G^{\max}$ . A path achieving this maximum is called a **critical path**.
- ▶  $h^{\text{add}}(n)$  is the **sum**, over all effect nodes  $n'$ , of the cost of  $n'$  multiplied by the **number of paths** from  $n$  to  $n'$  in  $G^{\text{add}}$ .

In particular, these properties hold for the goal node  $n_\gamma$  if it is reachable.

Example: Undercounting in  $h^{\max}$ 

$\rightsquigarrow o_1$  and  $o_4$  not counted because they are off the critical path

Example: Overcounting in  $h^{\text{add}}$ 

$\rightsquigarrow o_2$  counted twice because there are two paths to  $n_{o_2}^\top$

## D7.3 Summary

## Summary

- ▶  $h^{\max}$  and  $h^{\text{add}}$  can be used to decide **how** to achieve OR nodes in a relaxed task graph  
~> **best achievers**
- ▶ **Best achiever graphs** help identify shortcomings of  $h^{\max}$  and  $h^{\text{add}}$  compared to the perfect delete relaxation heuristic  $h^+$ .
  - ▶  $h^{\max}$  **underestimates**  $h^+$  because it only considers the cost of a **critical path** for the relaxed planning task.
  - ▶  $h^{\text{add}}$  **overestimates**  $h^+$  because it double-counts operators occurring on **multiple paths** in the best achiever graph.