

# Planning and Optimization

## B5. Positive Normal Form and STRIPS

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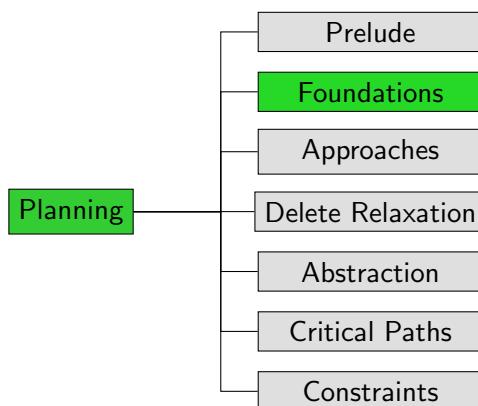
### B5.1 Motivation

### B5.2 Positive Normal Form

### B5.3 STRIPS

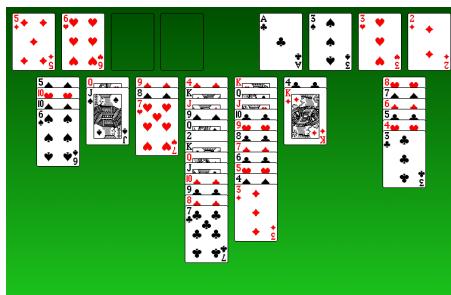
### B5.4 Summary

## Content of this Course



### B5.1 Motivation

## Example: Freecell



### Example (Good and Bad Effects)

If we move  $K\Diamond$  to a free tableau position,  
the **good effect** is that  $4\clubsuit$  is now accessible.  
The **bad effect** is that we lose one free tableau position.

## What is a Good or Bad Effect?

**Question:** Which operator effects are good, and which are bad?

Difficult to answer in general, because it depends on context:

- ▶ Locking our door is **good** if we want to keep burglars out.
- ▶ Locking our door is **bad** if we want to enter.

We now consider a reformulation of propositional planning tasks that makes the distinction between good and bad effects obvious.

## B5.2 Positive Normal Form

## Positive Formulas, Operators and Tasks

### Definition (Positive Formula)

A logical formula  $\varphi$  is **positive** if no negation symbols appear in  $\varphi$ .

**Note:** This includes the negation symbols implied by  $\rightarrow$  and  $\leftrightarrow$ .

### Definition (Positive Operator)

An operator  $o$  is **positive** if  $pre(o)$  and all effect conditions in  $eff(o)$  are positive.

### Definition (Positive Propositional Planning Task)

A propositional planning task  $\langle V, I, O, \gamma \rangle$  is **positive** if all operators in  $O$  and the goal  $\gamma$  are positive.

## Positive Normal Form

### Definition (Positive Normal Form)

A propositional planning task is in **positive normal form** if it is positive and all operator effects are flat.

## Positive Normal Form: Example

### Example (Transformation to Positive Normal Form)

$$\begin{aligned}
 V &= \{ \text{home}, \text{uni}, \text{lecture}, \text{bike}, \text{bike-locked} \} \\
 I &= \{ \text{home} \mapsto \mathbf{T}, \text{bike} \mapsto \mathbf{T}, \text{bike-locked} \mapsto \mathbf{T}, \\
 &\quad \text{uni} \mapsto \mathbf{F}, \text{lecture} \mapsto \mathbf{F} \} \\
 O &= \{ \langle \text{home} \wedge \text{bike} \wedge \neg \text{bike-locked}, \neg \text{home} \wedge \text{uni} \rangle, \\
 &\quad \langle \text{bike} \wedge \text{bike-locked}, \neg \text{bike-locked} \rangle, \\
 &\quad \langle \text{bike} \wedge \neg \text{bike-locked}, \text{bike-locked} \rangle, \\
 &\quad \langle \text{uni}, \text{lecture} \wedge ((\text{bike} \wedge \neg \text{bike-locked}) \triangleright \neg \text{bike}) \rangle \} \\
 \gamma &= \text{lecture} \wedge \text{bike}
 \end{aligned}$$

## Positive Normal Form: Example

### Example (Transformation to Positive Normal Form)

$$\begin{aligned}
 V &= \{ \text{home}, \text{uni}, \text{lecture}, \text{bike}, \text{bike-locked} \} \\
 I &= \{ \text{home} \mapsto \mathbf{T}, \text{bike} \mapsto \mathbf{T}, \text{bike-locked} \mapsto \mathbf{T}, \\
 &\quad \text{uni} \mapsto \mathbf{F}, \text{lecture} \mapsto \mathbf{F} \} \\
 O &= \{ \langle \text{home} \wedge \text{bike} \wedge \neg \text{bike-locked}, \neg \text{home} \wedge \text{uni} \rangle, \\
 &\quad \langle \text{bike} \wedge \text{bike-locked}, \neg \text{bike-locked} \rangle, \\
 &\quad \langle \text{bike} \wedge \neg \text{bike-locked}, \text{bike-locked} \rangle, \\
 &\quad \langle \text{uni}, \text{lecture} \wedge ((\text{bike} \wedge \neg \text{bike-locked}) \triangleright \neg \text{bike}) \rangle \} \\
 \gamma &= \text{lecture} \wedge \text{bike}
 \end{aligned}$$

Identify state variable  $v$  occurring negatively in conditions.

## Positive Normal Form: Example

### Example (Transformation to Positive Normal Form)

$$\begin{aligned}
 V &= \{ \text{home}, \text{uni}, \text{lecture}, \text{bike}, \text{bike-locked}, \text{bike-unlocked} \} \\
 I &= \{ \text{home} \mapsto \mathbf{T}, \text{bike} \mapsto \mathbf{T}, \text{bike-locked} \mapsto \mathbf{T}, \\
 &\quad \text{uni} \mapsto \mathbf{F}, \text{lecture} \mapsto \mathbf{F}, \text{bike-unlocked} \mapsto \mathbf{F} \} \\
 O &= \{ \langle \text{home} \wedge \text{bike} \wedge \neg \text{bike-locked}, \neg \text{home} \wedge \text{uni} \rangle, \\
 &\quad \langle \text{bike} \wedge \text{bike-locked}, \neg \text{bike-locked} \rangle, \\
 &\quad \langle \text{bike} \wedge \neg \text{bike-locked}, \text{bike-locked} \rangle, \\
 &\quad \langle \text{uni}, \text{lecture} \wedge ((\text{bike} \wedge \neg \text{bike-locked}) \triangleright \neg \text{bike}) \rangle \} \\
 \gamma &= \text{lecture} \wedge \text{bike}
 \end{aligned}$$

Introduce new variable  $\hat{v}$  with complementary initial value.



## Positive Normal Form: Example

### Example (Transformation to Positive Normal Form)

$$V = \{ \text{home}, \text{uni}, \text{lecture}, \text{bike}, \text{bike-locked}, \text{bike-unlocked} \}$$

$$I = \{ \text{home} \mapsto \mathbf{T}, \text{bike} \mapsto \mathbf{T}, \text{bike-locked} \mapsto \mathbf{T},$$

$$\quad \text{uni} \mapsto \mathbf{F}, \text{lecture} \mapsto \mathbf{F}, \text{bike-unlocked} \mapsto \mathbf{F} \}$$

$$O = \{ \langle \text{home} \wedge \text{bike} \wedge \text{bike-unlocked}, \neg \text{home} \wedge \text{uni} \rangle,$$

$$\quad \langle \text{bike} \wedge \text{bike-locked}, \neg \text{bike-locked} \wedge \text{bike-unlocked} \rangle,$$

$$\quad \langle \text{bike} \wedge \text{bike-unlocked}, \text{bike-locked} \wedge \neg \text{bike-unlocked} \rangle,$$

$$\quad \langle \text{uni}, \text{lecture} \wedge ((\text{bike} \wedge \text{bike-unlocked}) \triangleright \neg \text{bike}) \rangle \}$$

$$\gamma = \text{lecture} \wedge \text{bike}$$

## Positive Normal Form: Existence

### Theorem (Positive Normal Form)

For every propositional planning task  $\Pi$ , there is an equivalent propositional planning task  $\Pi'$  in positive normal form. Moreover,  $\Pi'$  can be computed from  $\Pi$  in polynomial time.

**Note:** Equivalence here means that the transition systems induced by  $\Pi$  and  $\Pi'$ , **restricted to the reachable states**, are isomorphic.

We prove the theorem by describing a suitable algorithm.  
(However, we do not prove its correctness or complexity.)

## Positive Normal Form: Algorithm

### Transformation of $\langle V, I, O, \gamma \rangle$ to Positive Normal Form

Replace all operators with equivalent conflict-free operators.

Convert all conditions to negation normal form (NNF).

**while** any condition contains a negative literal  $\neg v$ :

Let  $v$  be a variable which occurs negatively in a condition.

$V := V \cup \{\hat{v}\}$  for some new propositional state variable  $\hat{v}$

$$I(\hat{v}) := \begin{cases} \mathbf{F} & \text{if } I(v) = \mathbf{T} \\ \mathbf{T} & \text{if } I(v) = \mathbf{F} \end{cases}$$

Replace the effect  $v$  by  $(v \wedge \neg \hat{v})$  in all operators  $o \in O$ .

Replace the effect  $\neg v$  by  $(\neg v \wedge \hat{v})$  in all operators  $o \in O$ .

Replace  $\neg v$  by  $\hat{v}$  in all conditions.

Convert all operators  $o \in O$  to flat operators.

Here, **all conditions** refers to all operator preconditions, operator effect conditions and the goal.

## Why Positive Normal Form is Interesting

In the **absence of conditional effects**, positive normal form allows us to distinguish good and bad effects easily:

- ▶ Effects that make state variables true (**add effects**) are good.
- ▶ Effects that make state variables false (**delete effects**) are bad.

This is particularly useful for planning algorithms based on **delete relaxation**, which we will study in Part D.

(Why restriction “in the absence of conditional effects”?)

## B5.3 STRIPS

### STRIPS Operators: Remarks

- ▶ Every STRIPS operator is of the form

$$\langle v_1 \wedge \dots \wedge v_n, \ell_1 \wedge \dots \wedge \ell_m \rangle$$

where  $v_i$  are state variables and  $\ell_j$  are atomic effects.

- ▶ Often, STRIPS operators  $o$  are described via three **sets of state variables**:
  - ▶ the **preconditions** (state variables occurring in  $pre(o)$ )
  - ▶ the **add effects** (state variables occurring positively in  $eff(o)$ )
  - ▶ the **delete effects** (state variables occurring negatively in  $eff(o)$ )
- ▶ Definitions of STRIPS in the literature often do **not** require conflict-freeness. But it is easy to achieve and makes many things simpler.
- ▶ There exists a variant called **STRIPS with negation** where negative literals are also allowed in conditions.

## STRIPS Operators and Planning Tasks

### Definition (STRIPS Operator)

An operator  $o$  of a prop. planning task is a **STRIPS operator** if

- ▶  $pre(o)$  is a conjunction of state variables, and
- ▶  $eff(o)$  is a conflict-free conjunction of atomic effects.

### Definition (STRIPS Planning Task)

A propositional planning task  $\langle V, I, O, \gamma \rangle$  is a **STRIPS planning task** if all operators  $o \in O$  are STRIPS operators and  $\gamma$  is a conjunction of state variables.

Note: STRIPS operators are conflict-free and flat.  
STRIPS is a special case of positive normal form.

## Why STRIPS is Interesting

- ▶ STRIPS is **particularly simple**, yet expressive enough to capture general planning tasks.
- ▶ In particular, STRIPS planning is **no easier** than planning in general (as we will see in Chapter B6).
- ▶ Many algorithms in the planning literature are **only presented for STRIPS planning tasks** (generalization is often, but not always, obvious).

### STRIPS

STanford Research Institute Problem Solver  
(Fikes & Nilsson, 1971)

## Transformation to STRIPS

- ▶ Not every operator is equivalent to a STRIPS operator.
- ▶ However, each operator can be transformed into a **set** of STRIPS operators whose “combination” is equivalent to the original operator. (How?)
- ▶ However, this transformation may exponentially increase the number of operators. There are planning tasks for which such a blow-up is unavoidable.
- ▶ There are polynomial transformations of propositional planning tasks to STRIPS, but these do not lead to isomorphic transition systems (auxiliary states are needed). (They are, however, equivalent in a weaker sense.)

## B5.4 Summary

## Summary

- ▶ A **positive** task helps distinguish good and bad effects. The notion of positive tasks only exists for **propositional** tasks.
- ▶ A positive task with flat operators is in **positive normal form**.
- ▶ **STRIPS** is even more restrictive than positive normal form, forbidding complex preconditions and conditional effects.
- ▶ Both forms are expressive enough to capture general propositional planning tasks.
- ▶ Transformation to positive normal form is possible with polynomial size increase.
- ▶ Isomorphic transformations of propositional planning tasks to STRIPS can increase the number of operators exponentially; non-isomorphic polynomial transformations exist.