Discrete Mathematics in Computer Science

B7. Operations on Relations

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Relations: Recap

- A relation over sets $S_1, \dots, S_n$ is a set $R \subseteq S_1 \times \dots \times S_n$.
- A binary relation is a relation over two sets.
- A homogeneous relation R over set S is a binary relation $R \subseteq S \times S$.

Set Operations

- Relations are sets of tuples, so we can build their union, intersection, complement, ...
- Let R be a relation over $S_1, \dots, S_n$ and R' a relation over $S'_1, \dots, S'_n$. Then $R \cup R'$ is a relation over $S_1 \cup S'_1, \dots, S_n \cup S'_n$.
  With the standard relations $<, =$ and $\leq$ for $\mathbb{N}_0$, relation $\leq$ corresponds to the union of relations $<$ and $=$.
- Let R and R' be relations over n sets.
  Then $R \cap R'$ is a relation.
  Over which sets?
  With the standard relations $\leq, =$ and $\geq$ for $\mathbb{N}_0$, relation $=$ corresponds to the intersection of $\leq$ and $\geq$.
- If R is a relation over $S_1, \dots, S_n$ then so is the complementary relation $\bar{R} = (S_1 \times \dots \times S_n) \setminus R$.
  With the standard relations for $\mathbb{N}_0$, relation $=$ is the complementary relation of $\neq$ and $>$ the one of $\leq$.

Inverse of a Relation

Definition
Let $R \subseteq A \times B$ be a binary relation over A and B.
The inverse relation of R is the relation $R^{-1} \subseteq B \times A$ given by $R^{-1} = \{(b, a) \mid (a, b) \in R\}$.

- The inverse of the $<$ relation over $\mathbb{N}_0$ is the $>$ relation.
- Relation R with xRy iff person x has a key for y.
  Inverse: Q with aQb iff lock a can be openened by person b.

Composition of Relations

Definition (Composition of relations)
Let R_1 be a relation over A and B and R_2 a relation over B and C.
The composition of R_1 and R_2 is the relation $R_2 \circ R_1$ over A and C with:

$$R_2 \circ R_1 = \{(a, c) \mid \textit{there is a } b \in B \text{ with } (a, b) \in R_1 \text{ and } (b, c) \in R_2\}$$

How can we illustrate this graphically?

Composition of Relations: Example

$S_1 = \{1, 2, 3, 4\}$
$S_2 = \{A, B, C, D, E\}$
$S_3 = \{a, b, c, d\}$

$R_1 = \{(1, A), (1, B), (3, B), (4, D)\}$ over S_1 and S_2

$R_2 = \{(B, a), (C, c), (D, a), (D, d)\}$ over S_2 and S_3

$R_2 \circ R_1 =$

Composition is Associative

Theorem (Associativity of composition)

Let $S_1, \dots, S_4$ be sets and R_1, R_2, R_3 relations with $R_i \subseteq S_i \times S_{i+1}$. Then

$$R_3 \circ (R_2 \circ R_1) = (R_3 \circ R_2) \circ R_1.$$

Proof.

It holds that $(x_1, x_4) \in R_3 \circ (R_2 \circ R_1)$ iff there is an x_3 with $(x_1, x_3) \in R_2 \circ R_1$ and $(x_3, x_4) \in R_3$.

As $(x_1, x_3) \in R_2 \circ R_1$ iff there is an x_2 with $(x_1, x_2) \in R_1$ and $(x_2, x_3) \in R_2$, we have overall that $(x_1, x_4) \in R_3 \circ (R_2 \circ R_1)$ iff there are x_2, x_3 with $(x_1, x_2) \in R_1$, $(x_2, x_3) \in R_2$ and $(x_3, x_4) \in R_3$.

This is the case iff there is an x_2 with $(x_1, x_2) \in R_1$ and $(x_2, x_4) \in R_3 \circ R_2$, which holds iff $(x_1, x_4) \in (R_3 \circ R_2) \circ R_1$. □

(Reflexive) Transitive Closure

Definition ((Reflexive) transitive closure)

Let R be a relation over set S.

The transitive closure R^+ of R is the smallest relation over S that is transitive and has R as a subset.

The reflexive transitive closure R^* of R is the smallest relation over S that is reflexive, transitive and has R as a subset.

The (reflexive) transitive closure always exists. Why?

Example: If aRb specifies that there is a direct flight from a to b, what do R^+ and R^* express?

Transitive Closure and n-fold Composition

Define the n-fold composition of a relation R over S as

$$R^0 = \{(x, x) \mid x \in S\} \quad \text{and}$$
$$R^i = R \circ R^{i-1} \quad \text{for } i > 1.$$

Theorem

Let R be a relation over set S.
Then $R^+ = \bigcup_{i=1}^{\infty} R^i$ and $R^* = \bigcup_{i=0}^{\infty} R^i$.

Without proof.

Other Operators

- There are many more operators, also for general relations.
- Highly relevant for queries over relational databases.
- For example, join operators combine relations based on common entries.
- Example for a natural join:

Employee

Name	EmpId	DeptName
Harry	3415	Finance
Sally	2241	Sales
George	3401	Finance
Harriet	2202	Sales
Mary	1257	Human Resources

Dept

DeptName	Manager
Finance	George
Sales	Harriet
Production	Charles

Employee ⋈ Dept

Name	EmpId	DeptName	Manager
Harry	3415	Finance	George
Sally	2241	Sales	Harriet
George	3401	Finance	George
Harriet	2202	Sales	Harriet

(Source: Wikipedia)

Summary

- Relations: general, binary, homogeneous
- Properties: reflexivity, symmetry, transitivity
  (and related properties)
- Special relations: equivalence relations, order relations
- Operations: inverse, composition, transitive closure