

Discrete Mathematics in Computer Science

B1. Sets: Foundations

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Sets

Important Building Blocks of Discrete Mathematics

- sets
- relations
- functions

Sets

Definition

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- **unordered**: no notion of a “first” or “second” object,
e. g. $\{Alice, Bob, Charly\} = \{Charly, Bob, Alice\}$
- **distinct**: each object contained **at most once**,
e. g. $\{Alice, Bob, Charly\} = \{Alice, Charly, Bob, Alice\}$

Notation

- Specification of sets
 - **explicit**, listing all elements, e. g. $A = \{1, 2, 3\}$
 - **implicit** with **set-builder notation**, specifying a **property** characterizing all elements, e. g. $A = \{x \mid x \in \mathbb{N}_0 \text{ and } 1 \leq x \leq 3\}$,
 $B = \{n^2 \mid n \in \mathbb{N}_0\}$
 - **implicit**, as a **sequence with dots**, e. g. $\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$
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Question: Is it true that $1 \in \{\{1, 2\}, 3\}$?

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- **Rational numbers** $\mathbb{Q} = \{n/d \mid n \in \mathbb{Z}, d \in \mathbb{N}_1\}$

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- **Real numbers** $\mathbb{R} = (-\infty, \infty)$

Why do we use interval notation?

Why didn't we introduce it before?

Questions



Questions?

Russell's Paradox

Excursus: Barber Paradox

Barber Paradox

In a town there is only one barber, who is male.
The barber shaves all men in the town,
and only those, who do not shave themselves.



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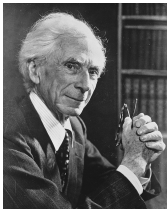
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We can exploit the self-reference to derive a contradiction.

Russell's Paradox

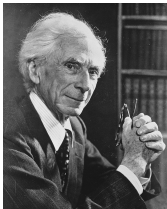


Bertrand Russell

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Is the collection of all sets that do not contain themselves as a member a set?

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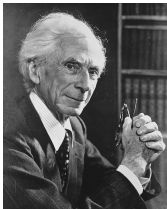
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Is $S = \{M \mid M \text{ is a set and } M \notin M\}$ a set?

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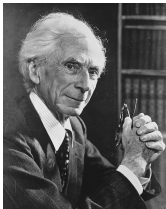
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Hence, there is no such set S .

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Hence, there is no such set S .

→ Not every property used in set-builder notation defines a set.

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Relations on Sets

Equality

Definition (Axiom of Extensionality)

Two sets A and B are **equal** (written $A = B$) if every element of A is an element of B and vice versa.

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We write $A \neq B$ to indicate that A and B are **not** equal.

Subsets and Supersets

- $A \subseteq B$: A is a **subset** of B ,
i. e., every element of A is an element of B
- $A \subset B$: A is a **strict subset** of B ,
i. e., $A \subseteq B$ and $A \neq B$.
- $A \supseteq B$: A is a **superset** of B if $B \subseteq A$.
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We write $A \not\subseteq B$ to indicate that A is **not** a subset of B .

Analogously: $\not\subset$, $\not\supseteq$, $\not\supset$

Power Set

Definition (Power Set)

The **power set** $\mathcal{P}(S)$ of a set S is the set of all subsets of S .
That is,

$$\mathcal{P}(S) = \{M \mid M \subseteq S\}.$$

Example: $\mathcal{P}(\{a, b\}) =$

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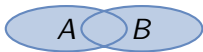
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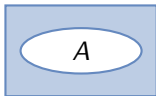
- **union** $A \cup B = \{x \mid x \in A \text{ or } x \in B\}$



- **set difference** $A \setminus B = \{x \mid x \in A \text{ and } x \notin B\}$



- **complement** $\bar{A} = B \setminus A$, where $A \subseteq B$ and B is the set of all considered objects (in a given context)



Properties of Set Operations: Commutativity

Theorem (Commutativity of \cup and \cap)

For all sets A and B it holds that

- $A \cup B = B \cup A$ and
- $A \cap B = B \cap A$.

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- $A \cap B = B \cap A$.

Question: Is the set difference also commutative,
i. e. is $A \setminus B = B \setminus A$ for all sets A and B ?

Properties of Set Operations: Associativity

Theorem (Associativity of \cup and \cap)

For all sets A, B and C it holds that

- $(A \cup B) \cup C = A \cup (B \cup C)$ and
- $(A \cap B) \cap C = A \cap (B \cap C)$.

Properties of Set Operations: Distributivity

Theorem (Union distributes over intersection and vice versa)

For all sets A , B and C it holds that

- $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$ and
- $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$.

Properties of Set Operations: De Morgan's Law



Augustus De Morgan

British mathematician (1806-1871)

Theorem (De Morgan's Law)

For all sets A and B it holds that

- $\overline{A \cup B} = \overline{A} \cap \overline{B}$ and
- $\overline{A \cap B} = \overline{A} \cup \overline{B}$.

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Finite Sets

Cardinality of Sets

The **cardinality** $|S|$ measures the size of set S .

A set is **finite** if it has a finite number of elements.

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Definition (Cardinality)

The **cardinality** of a finite set is the **number of elements** it contains.

- $|\emptyset| =$
- $|\{x \mid x \in \mathbb{N}_0 \text{ and } 2 \leq x < 5\}| =$
- $|\{3, 0, \{1, 3\}\}| =$
- $|\mathcal{P}(\{1, 2\})| =$

Cardinality of the Union of Sets

Theorem

For finite sets A and B it holds that $|A \cup B| = |A| + |B| - |A \cap B|$.

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Corollary

If finite sets A and B are *disjoint* then $|A \cup B| = |A| + |B|$.

Cardinality of the Power Set

Theorem

Let S be a finite set. Then $|\mathcal{P}(S)| = 2^{|S|}$.

Proof sketch.

We can construct a subset S' by iterating over all elements e of S and deciding whether e becomes a member of S' or not.

We make $|S|$ independent decisions, each between two options. Hence, there are $2^{|S|}$ possible outcomes.

Every subset of S can be constructed this way and different choices lead to different sets. Thus, $|\mathcal{P}(S)| = 2^{|S|}$. □

Alternative Proof by Induction

Proof.

By induction over $|S|$.

Basis ($|S| = 0$): Then $S = \emptyset$ and $|\mathcal{P}(S)| = |\{\emptyset\}| = 1 = 2^0$.

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Inductive Step ($n \rightarrow n + 1$):

Let S' be an arbitrary set with $|S'| = n + 1$ and let e be an arbitrary member of S' .

Let further $S = S' \setminus \{e\}$ and $X = \{S'' \cup \{e\} \mid S'' \in \mathcal{P}(S)\}$.

Then $\mathcal{P}(S') = \mathcal{P}(S) \cup X$. As $\mathcal{P}(S)$ and X are disjoint and $|X| = |\mathcal{P}(S)|$, it holds that $|\mathcal{P}(S')| = 2|\mathcal{P}(S)|$.

Since $|S| = n$, we can use the IH and get

$$|\mathcal{P}(S')| = 2 \cdot 2^{|S|} = 2 \cdot 2^n = 2^{n+1} = 2^{|S'|}.$$



Enumerating all Subsets

Determine a one-to-one mapping between numbers $0, \dots, 2^{|S|} - 1$ and all subsets of finite set S :

$$S = \{a, b, c\}$$

- Consider the binary representation of numbers $0, \dots, 2^{|S|} - 1$.
- Associate every bit with a different element of S .
- Every number is mapped to the set that contains exactly the elements associated with the 1-bits.

decimal	binary	set
	abc	
0	000	$\{\}$
1	001	$\{c\}$
2	010	$\{b\}$
3	011	$\{b, c\}$
4	100	$\{a\}$
5	101	$\{a, c\}$
6	110	$\{a, b\}$
7	111	$\{a, b, c\}$

Computer Representation as Bit String

Same representation as in enumeration of all subsets:

- **Required:** Fixed universe U of possible elements
- Represent sets as bitstrings of length $|U|$
- Associate every bit with one object from the universe
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Example:

- $U = \{o_0, \dots, o_9\}$
- Associate the i -th bit (0-indexed, from left to right) with o_i
- $\{o_2, o_4, o_5, o_9\}$ is represented as:
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How can the set operations be implemented?

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Summary

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- Sets are **unordered collections** of **distinct** objects.
- Important **set relations**: **equality** ($=$), **subset** (\subseteq), **superset** (\supseteq) and strict variants (\subset and \supset)
- The **power set** of a set S is the set of all subsets of S .
- Important **set operations** are **intersection**, **union**, **set difference** and **complement**.
 - Union and intersection are **commutative and associative**.
 - Union distributes over intersection and vice versa.
 - **De Morgan's law** for complement of union or intersection.
- The number of elements in a finite set is called its **cardinality**.
- Sets over a finite universe can be represented as bit strings.
→ also useful for enumerating all subsets