

# Planning and Optimization

## C4. SAT Planning: Core Idea and Sequential Encoding

Malte Helmert and Gabriele Röger

Universität Basel

October 12, 2022

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## C4.1 Introduction

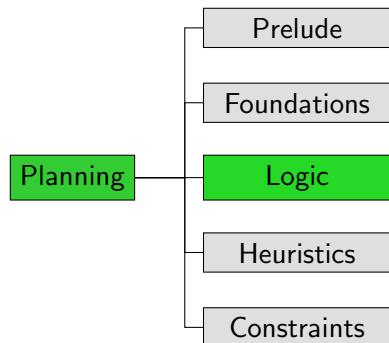
## C4.2 Formula Overview

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## Content of this Course



## C4.1 Introduction

## SAT Solvers

- ▶ **SAT solvers** (algorithms that find satisfying assignments to CNF formulas) are one of the major success stories in solving hard combinatorial problems.
- ▶ Can we leverage them for classical planning?
- ~~ **SAT planning** (a.k.a. planning as satisfiability)

background on SAT Solvers:

- ~~ Foundations of Artificial Intelligence Course, Ch. 31–32

## Solution: Iterative Deepening

- ▶ We can generate a propositional formula that tests if task  $\Pi$  has a plan with **horizon** (length bound)  $T$  in time  $O(\|\Pi\|^k \cdot T)$  (~~ pseudo-polynomial reduction).
- ▶ Use as building block of algorithm that probes increasing horizons (a bit like IDA\*).
- ▶ Can be efficient if there exist plans that are **not excessively long**.

## Complexity Mismatch

- ▶ The SAT problem is **NP-complete**, while PLANEX is **PSPACE-complete**.
- ~~ one-shot polynomial reduction from PLANEX to SAT not possible (unless NP = PSPACE)

## SAT Planning: Main Loop

basic SAT Planning algorithm:

```

SAT Planning
def satplan( $\Pi$ ):
    for  $T \in \{0, 1, 2, \dots\}$ :
         $\varphi := \text{build\_sat\_formula}(\Pi, T)$ 
         $I = \text{sat\_solver}(\varphi)$                                  $\triangleright$  returns a model or none
        if  $I$  is not none:
            return extract_plan( $\Pi, T, I$ )
    
```

Termination criterion for unsolvable tasks?

## C4.2 Formula Overview

### SAT Formula: CNF?

- ▶ SAT solvers require **conjunctive normal form** (CNF), i.e., formulas expressed as collection of **clauses**.
- ▶ We will make sure that our SAT formulas are in CNF when our input is a **STRIPS** task.
- ▶ We do allow fully general propositional tasks, but then the formula may need additional conversion to CNF.

### SAT Formula: Variables

- ▶ given propositional planning task  $\Pi = \langle V, I, O, \gamma \rangle$
- ▶ given **horizon**  $T \in \mathbb{N}_0$

#### Variables of the SAT Formula

- ▶ propositional variables  $v^i$  for all  $v \in V, 0 \leq i \leq T$   
encode **state after  $i$  steps** of the plan
- ▶ propositional variables  $o^i$  for all  $o \in O, 1 \leq i \leq T$   
encode **operator(s) applied in  $i$ -th step** of the plan

### Formulas with Time Steps

#### Definition (Time-Stamped Formulas)

Let  $\varphi$  be a propositional logic formula over the variables  $V$ .  
Let  $0 \leq i \leq T$ .

We write  $\varphi^i$  for the formula obtained from  $\varphi$  by replacing each  $v \in V$  with  $v^i$ .

**Example:**  $((a \wedge b) \vee \neg c)^3 = (a^3 \wedge b^3) \vee \neg c^3$

## SAT Formula: Motivation

We want to express a **formula** whose **models** are exactly the plans/traces with  $T$  steps.

For this, the formula must express four things:

- ▶ The variables  $v^0$  ( $v \in V$ ) define the initial state.
- ▶ The variables  $v^T$  ( $v \in V$ ) define a goal state.
- ▶ We select exactly one operator variable  $o^i$  ( $o \in O$ ) for each time step  $1 \leq i \leq T$ .
- ▶ If we select  $o^i$ , then variables  $v^{i-1}$  and  $v^i$  ( $v \in V$ ) describe a state transition from the  $(i-1)$ -th state of the plan to the  $i$ -th state of the plan (that uses operator  $o$ ).

The final formula is the **conjunction** of all these parts.

## SAT Formula: Initial State

### SAT Formula: Initial State

initial state clauses:

- ▶  $v^0$  for all  $v \in V$  with  $I(v) = \mathbf{T}$
- ▶  $\neg v^0$  for all  $v \in V$  with  $I(v) = \mathbf{F}$

## C4.3 Initial State, Goal, Operator Selection

## SAT Formula: Goal

### SAT Formula: Goal

goal clauses:

- ▶  $\gamma^T$

For STRIPS, this is a conjunction of unit clauses.  
For general goals, this may not be in clause form.

## SAT Formula: Operator Selection

Let  $\mathcal{O} = \{o_1, \dots, o_n\}$ .

### SAT Formula: Operator Selection

operator selection clauses:

- ▶  $o_1^i \vee \dots \vee o_n^i$  for all  $1 \leq i \leq T$

operator exclusion clauses:

- ▶  $\neg o_j^i \vee \neg o_k^i$  for all  $1 \leq i \leq T, 1 \leq j < k \leq n$

## SAT Formula: Transitions

We now get to the interesting/challenging bit: encoding the transitions.

Key observations: if we apply operator  $o$  at time  $i$ ,

- ▶ its **precondition** must be satisfied at time  $i - 1$ :  

$$o^i \rightarrow \text{pre}(o)^{i-1}$$
- ▶ variable  $v$  is true at time  $i$  iff its **regression** is true at  $i - 1$ :  

$$o^i \rightarrow (v^i \leftrightarrow \text{regr}(v, \text{eff}(o))^{i-1})$$

Question: Why  $\text{regr}(v, \text{eff}(o))$ , not  $\text{regr}(v, o)$ ?

## C4.4 Transitions

## SAT Formula: Transitions

## Simplifications and Abbreviations

- ▶ Let us pick the last formula apart to understand it better (and also get a CNF representation along the way).
- ▶ Let us call the formula  $\tau$  ("transition"):  

$$\tau = o^i \rightarrow (v^i \leftrightarrow \text{regr}(v, \text{eff}(o))^{i-1})$$
.
- ▶ First, some abbreviations:
  - ▶ Let  $e = \text{eff}(o)$ .
  - ▶ Let  $\rho = \text{regr}(v, e)$  ("regression").  
We have  $\rho = \text{effcond}(v, e) \vee (v \wedge \neg \text{effcond}(\neg v, e))$ .
  - ▶ Let  $\alpha = \text{effcond}(v, e)$  ("added").
  - ▶ Let  $\delta = \text{effcond}(\neg v, e)$  ("deleted").
- ~ $\tau = o^i \rightarrow (v^i \leftrightarrow \rho^{i-1})$  with  $\rho = \alpha \vee (v \wedge \neg \delta)$

## Picking it Apart (1)

Reminder:  $\tau = o^i \rightarrow (v^i \leftrightarrow \rho^{i-1})$  with  $\rho = \alpha \vee (v \wedge \neg\delta)$

$$\begin{aligned} \tau &= o^i \rightarrow (v^i \leftrightarrow \rho^{i-1}) \\ &\equiv o^i \rightarrow ((v^i \rightarrow \rho^{i-1}) \wedge (\rho^{i-1} \rightarrow v^i)) \\ &\equiv \underbrace{(o^i \rightarrow (v^i \rightarrow \rho^{i-1}))}_{\tau_1} \wedge \underbrace{(o^i \rightarrow (\rho^{i-1} \rightarrow v^i))}_{\tau_2} \end{aligned}$$

~~ consider this two **separate** constraints  $\tau_1$  and  $\tau_2$

## Interpreting the Constraints (1)

Can we give an **intuitive description** of  $\tau_{11}$  and  $\tau_{12}$ ?

~~ Yes!

►  $\tau_{11} = (o^i \wedge \neg\alpha^{i-1} \wedge \neg v^{i-1}) \rightarrow \neg v^i$

"When applying  $o$ , if  $v$  is false and  $o$  does not add it, it remains false."

► called **negative frame clause**

► in clause form:  $\neg o^i \vee \alpha^{i-1} \vee v^{i-1} \vee \neg v^i$

►  $\tau_{12} = (o^i \wedge \neg\alpha^{i-1} \wedge \delta^{i-1}) \rightarrow \neg v^i$

"When applying  $o$ , if  $o$  deletes  $v$  and does not add it, it is false afterwards." (Note the add-after-delete semantics.)

► called **negative effect clause**

► in clause form:  $\neg o^i \vee \alpha^{i-1} \vee \neg \delta^{i-1} \vee \neg v^i$

For STRIPS tasks, these are indeed clauses. (And in general?)

## Picking it Apart (2)

Reminder:  $\tau_1 = o^i \rightarrow (v^i \rightarrow \rho^{i-1})$  with  $\rho = \alpha \vee (v \wedge \neg\delta)$

$$\begin{aligned} \tau_1 &= o^i \rightarrow (v^i \rightarrow \rho^{i-1}) \\ &\equiv o^i \rightarrow (\neg\rho^{i-1} \rightarrow \neg v^i) \\ &\equiv (o^i \wedge \neg\rho^{i-1}) \rightarrow \neg v^i \\ &\equiv (o^i \wedge \neg(\alpha^{i-1} \vee (v^{i-1} \wedge \neg\delta^{i-1}))) \rightarrow \neg v^i \\ &\equiv (o^i \wedge (\neg\alpha^{i-1} \wedge (\neg v^{i-1} \vee \delta^{i-1}))) \rightarrow \neg v^i \\ &\equiv \underbrace{((o^i \wedge \neg\alpha^{i-1} \wedge \neg v^{i-1}) \rightarrow \neg v^i)}_{\tau_{11}} \wedge \underbrace{((o^i \wedge \neg\alpha^{i-1} \wedge \delta^{i-1}) \rightarrow \neg v^i)}_{\tau_{12}} \end{aligned}$$

~~ consider this two **separate** constraints  $\tau_{11}$  and  $\tau_{12}$

## Picking it Apart (3)

Almost done!

Reminder:  $\tau_2 = o^i \rightarrow (\rho^{i-1} \rightarrow v^i)$  with  $\rho = \alpha \vee (v \wedge \neg\delta)$

$$\begin{aligned} \tau_2 &= o^i \rightarrow (\rho^{i-1} \rightarrow v^i) \\ &\equiv (o^i \wedge \rho^{i-1}) \rightarrow v^i \\ &\equiv (o^i \wedge (\alpha^{i-1} \vee (v^{i-1} \wedge \neg\delta^{i-1}))) \rightarrow v^i \\ &\equiv \underbrace{((o^i \wedge \alpha^{i-1}) \rightarrow v^i)}_{\tau_{21}} \wedge \underbrace{((o^i \wedge v^{i-1} \wedge \neg\delta^{i-1}) \rightarrow v^i)}_{\tau_{22}} \end{aligned}$$

~~ consider this two **separate** constraints  $\tau_{21}$  and  $\tau_{22}$

## Interpreting the Constraints (2)

How about an **intuitive description** of  $\tau_{21}$  and  $\tau_{22}$ ?

- ▶  $\tau_{21} = (o^i \wedge \alpha^{i-1}) \rightarrow v^i$

"When applying  $o$ , if  $o$  adds  $v$ , it is true afterwards."

- ▶ called **positive effect clause**

- ▶ in clause form:  $\neg o^i \vee \neg \alpha^{i-1} \vee v^i$

- ▶  $\tau_{22} = (o^i \wedge v^{i-1} \wedge \neg \delta^{i-1}) \rightarrow v^i$

"When applying  $o$ , if  $v$  is true and  $o$  does not delete it, it remains true."

- ▶ called **positive frame clause**

- ▶ in clause form:  $\neg o^i \vee \neg v^{i-1} \vee \delta^{i-1} \vee v^i$

For STRIPS tasks, these are indeed clauses. (But not in general.)

## SAT Formula: Transitions

### SAT Formula: Transitions

precondition clauses:

- ▶  $\neg o^i \vee \text{pre}(o)^{i-1}$  for all  $1 \leq i \leq T$ ,  $o \in O$

positive and negative effect clauses:

- ▶  $\neg o^i \vee \neg \alpha^{i-1} \vee v^i$  for all  $1 \leq i \leq T$ ,  $o \in O$ ,  $v \in V$

- ▶  $\neg o^i \vee \alpha^{i-1} \vee \neg \delta^{i-1} \vee \neg v^i$  for all  $1 \leq i \leq T$ ,  $o \in O$ ,  $v \in V$

positive and negative frame clauses:

- ▶  $\neg o^i \vee \neg v^{i-1} \vee \delta^{i-1} \vee v^i$  for all  $1 \leq i \leq T$ ,  $o \in O$ ,  $v \in V$

- ▶  $\neg o^i \vee \alpha^{i-1} \vee v^{i-1} \vee \neg v^i$  for all  $1 \leq i \leq T$ ,  $o \in O$ ,  $v \in V$

where  $\alpha = \text{effcond}(v, \text{eff}(o))$ ,  $\delta = \text{effcond}(\neg v, \text{eff}(o))$ .

For STRIPS, all except the precondition clauses are in clause form.

The precondition clauses are easily convertible to CNF

(one clause  $\neg o^i \vee v^{i-1}$  for each precondition atom  $v$  of  $o$ ).

## C4.5 Summary

## Summary

- ▶ **SAT planning** (planning as satisfiability) expresses a sequence of bounded-horizon planning tasks as SAT formulas.
- ▶ Plans can be extracted from satisfying assignments; unsolvable tasks are challenging for the algorithm.
- ▶ For each **time step**, there are propositions encoding which state variables are true and which operators are applied.
- ▶ We describe a basic **sequential** encoding where one operator is applied at every time step.
- ▶ The encoding produces a **CNF** formula for **STRIPS** tasks.
- ▶ The encoding follows naturally (with some work) from using **regression** to link state variables in adjacent time steps.