Discrete Mathematics in Computer Science E6. Advanced Concepts in Predicate Logic and Outlook

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E6. Advanced Concepts in Predicate Logic and Outlook

Free and Bound Variables

1 / 25

E6.1 Free and Bound Variables



E6.1 Free and Bound Variables

E6.2 Reasoning in Predicate Logic

E6.3 Summary and Outlook

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Free and Bound Variables

2 / 25

Free and Bound Variables: Motivation

Question:

- Consider a signature with variable symbols {x₁, x₂, x₃,...}
 and an interpretation *I*.
- ▶ Which parts of the definition of α are relevant to decide whether $\mathcal{I}, \alpha \models (\forall x_4(\mathsf{R}(x_4, x_2) \lor (\mathsf{f}(x_3) = x_4)) \lor \exists x_3\mathsf{S}(x_3, x_2))$?
- α(x₁), α(x₅), α(x₆), α(x₇), ... are irrelevant since those variable symbols occur in no formula.
- α(x₄) also is irrelevant: the variable occurs in the formula, but all occurrences are bound by a surrounding quantifier.
- \blacktriangleright \rightarrow only assignments for free variables x_2 and x_3 relevant

German: gebundene und freie Variablen

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Free and Bound Variables

Variables of a Term

Definition (Variables of a Term)

Let t be a term. The set of variables that occur in t, written as var(t), is defined as follows:

- var(x) = {x}
 for variable symbols x
- Var(c) = ∅ for constant symbols c
- ► $var(f(t_1,...,t_k)) = var(t_1) \cup \cdots \cup var(t_k)$ for function terms

terminology: A term t with $var(t) = \emptyset$ is called ground term. German: Grundterm

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example: var(product(x, sum(k, y))) =
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5 / 25

7 / 25

Closed Formulas/Sentences

Note: Let φ be a formula and let α and β variable assignments with $\alpha(x) = \beta(x)$ for all free variables x of φ . Then $\mathcal{I}, \alpha \models \varphi$ iff $\mathcal{I}, \beta \models \varphi$.

In particular, α is completely irrelevant if $free(\varphi) = \emptyset$.

Definition (Closed Formulas/Sentences)

A formula φ without free variables (i. e., *free*(φ) = \emptyset) is called closed formula or sentence.

If φ is a sentence, then we often write $\mathcal{I} \models \varphi$ instead of $\mathcal{I}, \alpha \models \varphi$, since the definition of α does not influence whether φ is true under \mathcal{I} and α or not.

Formulas with at least one free variable are called open.

Closed formulas with no quantifiers are called ground formulas.

German: geschlossene Formel/Satz, offene Formel, Grundformel/variablenfreie Formel alte Helmert, Gabriele Röger (University of Discrete Mathematics in Computer Science

Free and Bound Variables of a Formula

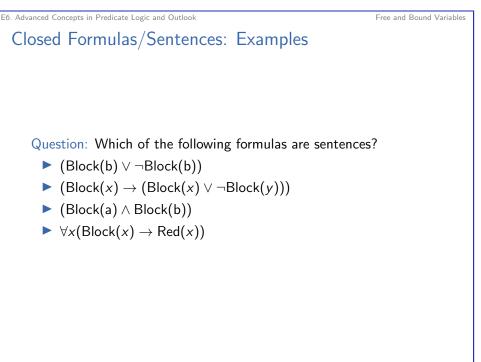
Definition (Free Variables)

Let φ be a predicate logic formula. The set of free variables of φ , written as *free*(φ), is defined as follows:

- $free(P(t_1,\ldots,t_k)) = var(t_1) \cup \cdots \cup var(t_k)$
- $\blacktriangleright free((t_1 = t_2)) = var(t_1) \cup var(t_2)$
- free($\neg \varphi$) = free(φ)
- $\blacktriangleright \ \textit{free}((\varphi \land \psi)) = \textit{free}((\varphi \lor \psi)) = \textit{free}(\varphi) \cup \textit{free}(\psi)$
- $\models free(\forall x \varphi) = free(\exists x \varphi) = free(\varphi) \setminus \{x\}$

Example: *free*(($\forall x_4(R(x_4, x_2) \lor (f(x_3) = x_4)) \lor \exists x_3S(x_3, x_2)))$ =

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6 / 25

Free and Bound Variables

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Reasoning in Predicate Logic

9 / 25

Sets of Formulas: Semantics

Definition (Satisfied/True Sets of Formulas)

Let S be a signature, Φ a set of formulas over S, \mathcal{I} an interpretation for S and α a variable assignment for Sand the universe of \mathcal{I} .

We say that \mathcal{I} and α satisfy the formulas Φ (also: Φ is true under \mathcal{I} and α), written as: $\mathcal{I}, \alpha \models \Phi$, if $\mathcal{I}, \alpha \models \varphi$ for all $\varphi \in \Phi$.

German: \mathcal{I} und α erfüllen Φ , Φ ist wahr unter \mathcal{I} und α

We may again write $\mathcal{I} \models \Phi$ if all formulas in Φ are sentences.

Reasoning in Predicate Logic

Terminology for Formulas

The terminology we introduced for propositional logic equally applies to predicate logic:

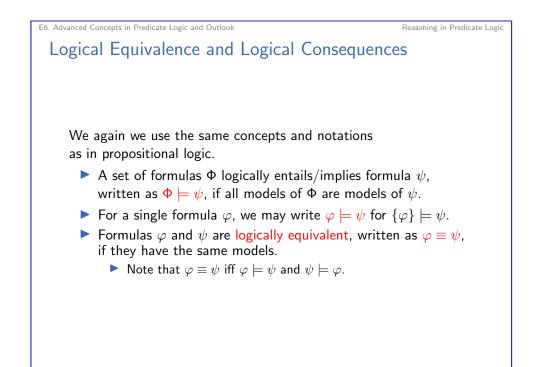
- Interpretation *I* and variable assignment α form a model of the formula φ if *I*, α ⊨ φ.
- Formula φ is satisfiable if $\mathcal{I}, \alpha \models \varphi$ for at least one \mathcal{I}, α .
- ▶ Formula φ is falsifiable if $\mathcal{I}, \alpha \not\models \varphi$. for at least one \mathcal{I}, α
- Formula φ is valid if $\mathcal{I}, \alpha \models \varphi$ for all \mathcal{I}, α .
- Formula φ is unsatisfiable if $\mathcal{I}, \alpha \not\models \varphi$ for all \mathcal{I}, α .

German: Modell, erfüllbar, falsifizierbar, gültig, unerfüllbar

All concepts can be used for the special case of sentences. In this case we usually omit α . Examples:

- lnterpretation \mathcal{I} is a model of a sentence φ if $\mathcal{I} \models \varphi$.
- Sentence φ is unsatisfiable if $\mathcal{I} \not\models \varphi$ for all \mathcal{I} .

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10 / 25

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Reasoning in Predicate Logic

Important Theorems about Logical Consequences

Theorem (Deduction Theorem) KB \cup { φ } $\models \psi$ *iff* KB \models ($\varphi \rightarrow \psi$)

German: Deduktionssatz

Theorem (Contraposition Theorem) KB $\cup \{\varphi\} \models \neg \psi \text{ iff } KB \cup \{\psi\} \models \neg \varphi$

German: Kontrapositionssatz

Theorem (Contradiction Theorem)

 $\mathsf{KB} \cup \{\varphi\} \text{ is unsatisfiable iff } \mathsf{KB} \models \neg \varphi$

German: Widerlegungssatz

These can be proved exactly the same way as in propositional logic.

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Reasoning in Predicate Logic

13 / 25

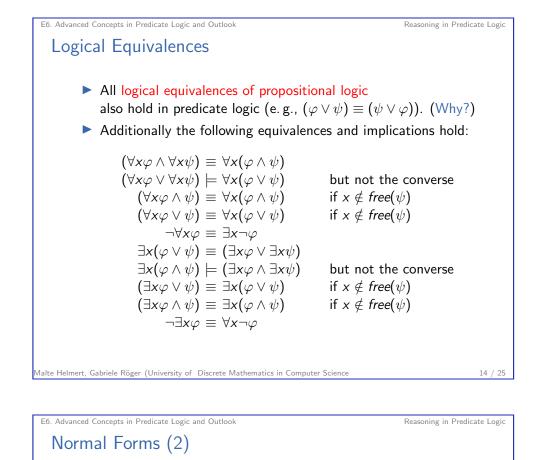
Normal Forms (1)

Analogously to DNF and CNF for propositional logic there are several normal forms for predicate logic, such as

- negation normal form (NNF): negation symbols (¬) are only allowed in front of atoms
- prenex normal form: quantifiers must form the outermost part of the formula
- Skolem normal form:

prenex normal form without existential quantifiers

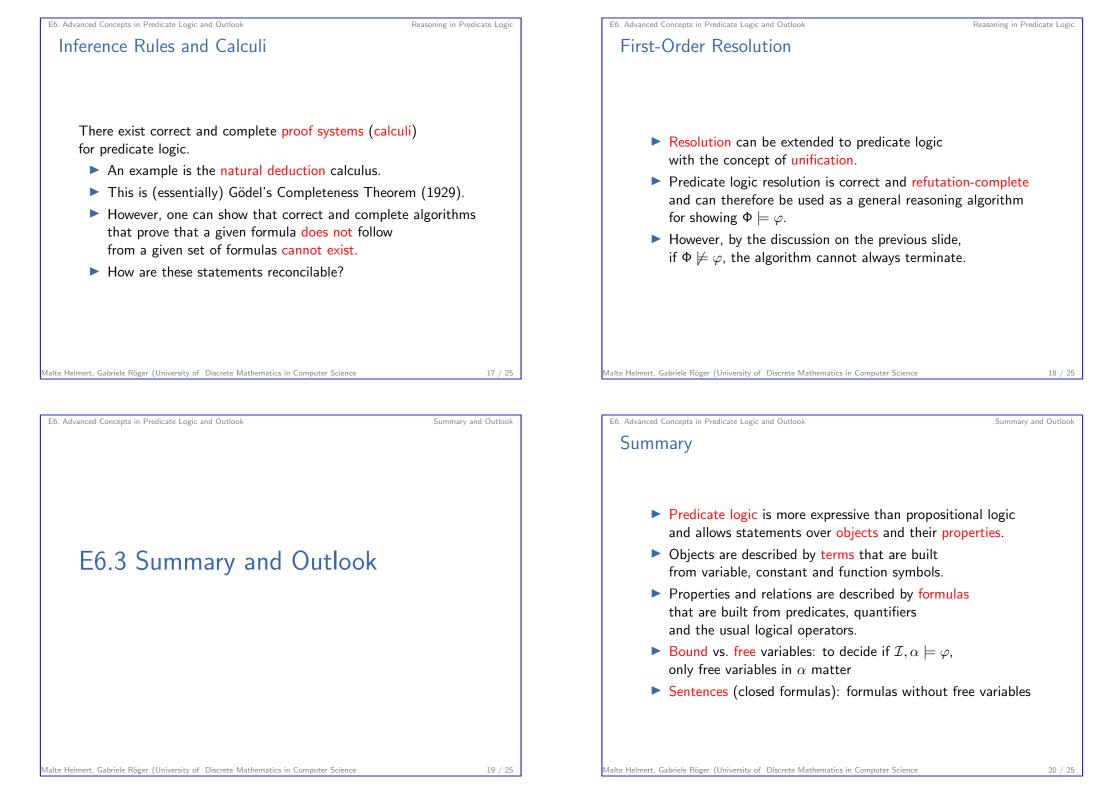
German: Negationsnormalform, Pränexnormalform, Skolemnormalform

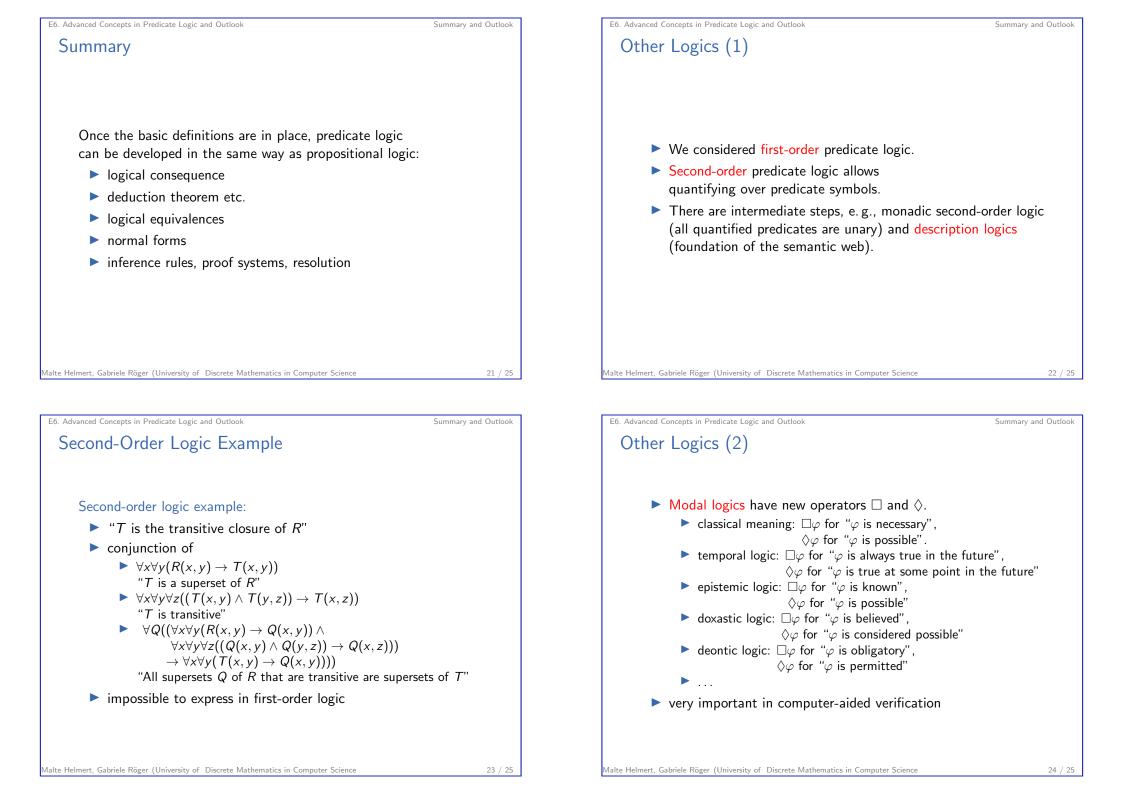


Efficient methods transform formula φ

- into an equivalent formula in negation normal form,
- ▶ into an equivalent formula in prenex normal form, or

into an equisatisfiable formula in Skolem normal form.
 German: erfüllbarkeitsäquivalent





Summary and Outlook

Other Logics (3)

- In fuzzy logic, formulas are not true or false but have values between 0 and 1.
- Intuitionist logic is "constructive" and excludes indirect proof methods such as the principle of the excluded third.
- Non-monotonic logics have rules with exceptions (e.g., default logic, cumulative logic).
- ▶ ...and there is a lot more

25 / 25