# Discrete Mathematics in Computer Science B7. Operations on Relations 

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- B7. Operations on Relations


## B7.1 Operations on Relations

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## Relations: Recap

- A relation over sets $S_{1}, \ldots, S_{n}$ is a set $R \subseteq S_{1} \times \cdots \times S_{n}$.
- A binary relation is a relation over two sets.
- A homogeneous relation $R$ over set $S$ is a binary relation $R \subseteq S \times S$.


## Set Operations

- Relations are sets of tuples, so we can build their union, intersection, complement, ....
- Let $R$ be a relation over $S_{1}, \ldots, S_{n}$ and $R^{\prime}$ a relation over $S_{1}^{\prime}, \ldots, S_{n}^{\prime}$. Then $R \cup R^{\prime}$ is a relation over $S_{1} \cup S_{1}^{\prime}, \ldots, S_{n} \cup S_{n}^{\prime}$.
With the standard relations $<,=$ and $\leq$ for $\mathbb{N}_{0}$, relation $\leq$ corresponds to the union of relations $<$ and $=$.
- Let $R$ and $R^{\prime}$ be relations over $n$ sets.

Then $R \cap R^{\prime}$ is a relation.
Over which sets?
With the standard relations $\leq,=$ and $\geq$ for $\mathbb{N}_{0}$, relation $=$ corresponds to the intersection of $\leq$ and $\geq$.

- If $R$ is a relation over $S_{1}, \ldots, S_{n}$ then so is the complementary relation $\bar{R}=\left(S_{1} \times \cdots \times S_{n}\right) \backslash R$. With the standard relations for $\mathbb{N}_{0}$, relation $=$ is the complementary relation of $\neq$ and $>$ the one of $\leq$.


## Inverse of a Relation

## Definition

Let $R \subseteq A \times B$ be a binary relation over $A$ and $B$.
The inverse relation of $R$ is the relation $R^{-1} \subseteq B \times A$ given by $R^{-1}=\{(b, a) \mid(a, b) \in R\}$.

- The inverse of the $<$ relation over $\mathbb{N}_{0}$ is the $>$ relation.
- Relation $R$ with $x R y$ iff person $x$ has a key for $y$. Inverse: $Q$ with $a Q b$ iff lock $a$ can be openened by person $b$.


## Composition of Relations

Definition (Composition of relations)
Let $R_{1}$ be a relation over $A$ and $B$ and $R_{2}$ be a relation over $B$ and $C$.

The composition of $R_{1}$ and $R_{2}$ is the relation $R_{2} \circ R_{1}$ with:

$$
\begin{aligned}
& R_{2} \circ R_{1}=\{(a, c) \mid \text { there is } a b \in B \text { with } \\
& \left.\qquad(a, b) \in R_{1} \text { and }(b, c) \in R_{2}\right\}
\end{aligned}
$$

How can we illustrate this graphically?

## Composition is Associative

## Theorem (Associativity of composition)

Let $S_{1}, \ldots, S_{4}$ be sets and $R_{1}, R_{2}, R_{3}$ relations with $R_{i} \subseteq S_{i} \times S_{i+1}$. Then

$$
R_{3} \circ\left(R_{2} \circ R_{1}\right)=\left(R_{3} \circ R_{2}\right) \circ R_{1} .
$$

## Proof.

It holds that $\left(x_{1}, x_{4}\right) \in R_{3} \circ\left(R_{2} \circ R_{1}\right)$ iff there is an $x_{3}$ with $\left(x_{1}, x_{3}\right) \in R_{2} \circ R_{1}$ and $\left(x_{3}, x_{4}\right) \in R_{3}$.

As $\left(x_{1}, x_{3}\right) \in R_{2} \circ R_{1}$ iff there is an $x_{2}$ with $\left(x_{1}, x_{2}\right) \in R_{1}$ and $\left(x_{2}, x_{3}\right) \in R_{2}$, we have overall that $\left(x_{1}, x_{4}\right) \in R_{3} \circ\left(R_{2} \circ R_{1}\right)$ iff there are $x_{2}, x_{3}$ with $\left(x_{1}, x_{2}\right) \in R_{1},\left(x_{2}, x_{3}\right) \in R_{2}$ and $\left(x_{3}, x_{4}\right) \in R_{3}$.

This is the case iff there is an $x_{2}$ with $\left(x_{1}, x_{2}\right) \in R_{1}$ and $\left(x_{2}, x_{4}\right) \in R_{3} \circ R_{2}$, which holds iff $\left(x_{1}, x_{4}\right) \in\left(R_{3} \circ R_{2}\right) \circ R_{1}$.

## Transitive Closure

## Definition (Transitive closure)

The transitive closure $R^{+}$of a relation $R$ over set $S$ is the smallest relation over $S$ that is transitive and has $R$ as a subset.

The transitive closure always exists. Why?
Example: If $a R b$ specifies that block $a$ lies on block $b$, what does $R^{+}$express?

## Transitive Closure II

Define the $i$-th power of a homogeneous relation $R$ as

$$
\begin{aligned}
& R^{1}=R \\
& R^{i}=R \circ R^{i-1} \\
& \text { if } i=1 \text { and } \\
& i>1
\end{aligned}
$$

Theorem
Let $R$ be a relation over set $S$. Then $R^{+}=\bigcup_{i=1}^{\infty} R^{i}$.
Without proof.

## Other Operators

- There are many more operators, also for general relations.
- Highly relevant for queries over relational databases.
- For example, join operators combine relations based on common entries.
- Example for a natural join:

| Employee |  |  |
| :--- | :--- | :--- |
| Name | Empld | DeptName |
| Harry | 3415 | Finance |
| Sally | 2241 | Sales |
| George | 3401 | Finance |
| Harriet | 2202 | Sales |
| Mary | 1257 | Human <br> Resources |


| Dept |  |
| :--- | :--- |
| DeptName | Manager |
| Finance | George |
| Sales | Harriet |
| Production | Charles |


| Employee $\pitchfork$ Dept |  |  |  |
| :--- | :--- | :--- | :--- |
| Name | Empld | DeptName | Manager |
| Harry | 3415 | Finance | George |
| Sally | 2241 | Sales | Harriet |
| George | 3401 | Finance | George |
| Harriet | 2202 | Sales | Harriet |

## Summary

- Relations: general, binary, homogeneous
- Properties: reflexivity, symmetry, transitivity (and related properties)
- Special relations: equivalence relations, order relations
- Operations: inverse, composition, transitive closure

