

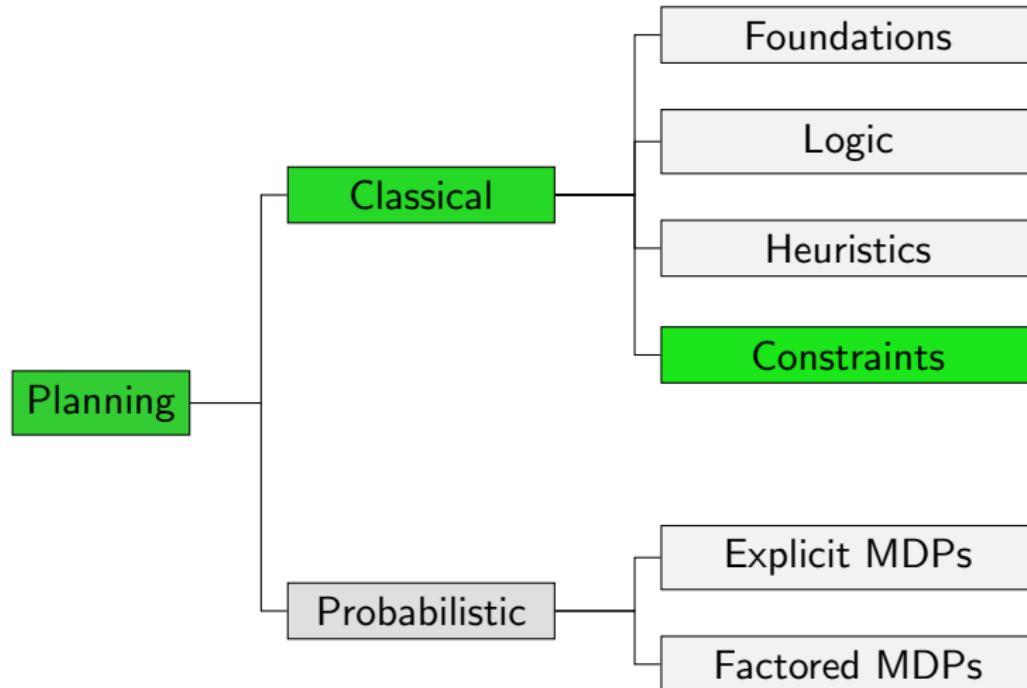
# Planning and Optimization

## E1. Constraints: Introduction

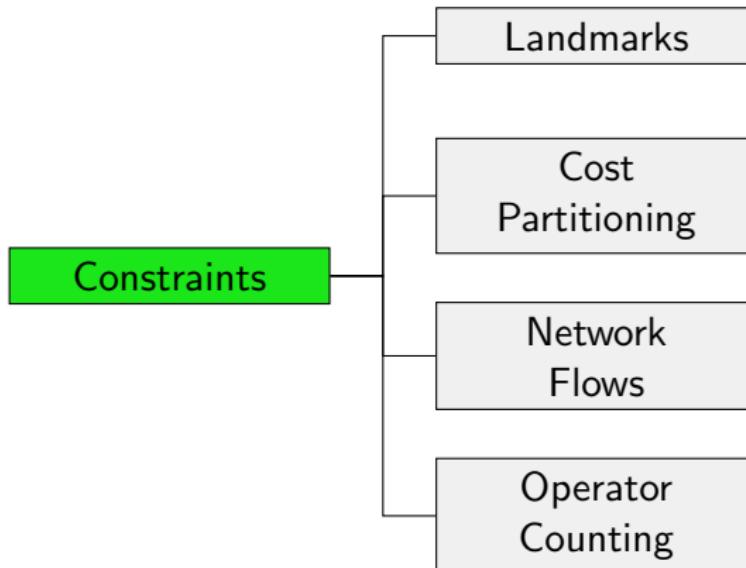
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Universität Basel

# Content of this Course



# Content of this Course: Constraints



# Constraint-based Heuristics

# Coming Up with Heuristics in a Principled Way

## General Procedure for Obtaining a Heuristic

Solve a simplified version of the problem.

Major ideas for heuristics in the planning literature:

- delete relaxation
- abstraction
- landmarks
- critical paths
- network flows
- potential heuristic

Landmarks, network flows and potential heuristics are based on **constraints** that can be specified for a planning task.

# Constraints: Example

	1	2	3	4	5	6
A						
B						
C						

Images from wikipedia

# Constraints: Example

## Example

Consider a FDR planning task  $\langle V, I, O, \gamma \rangle$  with

- $V = \{robot\text{-}at, dishes\text{-}at\}$  with
  - $\text{dom}(robot\text{-}at) = \{\text{A1}, \dots, \text{C3}, \text{B4}, \text{A5}, \dots, \text{B6}\}$
  - $\text{dom}(dishes\text{-}at) = \{\text{Table}, \text{Robot}, \text{Dishwasher}\}$
- $I = \{robot\text{-}at \mapsto \text{C1}, dishes\text{-}at \mapsto \text{Table}\}$
- operators
  - move- $x$ - $y$  to move from cell  $x$  to adjacent cell  $y$
  - pickup dishes, and
  - load dishes into the dishwasher.
- $\gamma = (robot\text{-}at = \text{B6}) \wedge (dishes\text{-}at = \text{Dishwasher})$

# Constraints

Some heuristics exploit **constraints** that describe something that holds in every solution of the task.

For instance, every solution is such that

- a variable takes some value in at least one visited state.  
(a **fact landmark** constraint)

# Fact Landmarks: Example

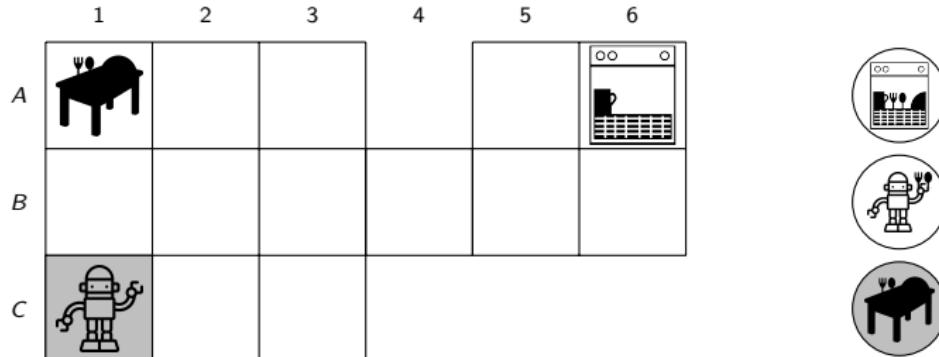
Which values do *robot-at* and *dishes-at* take in every solution?

	1	2	3	4	5	6
A						
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# Fact Landmarks: Example

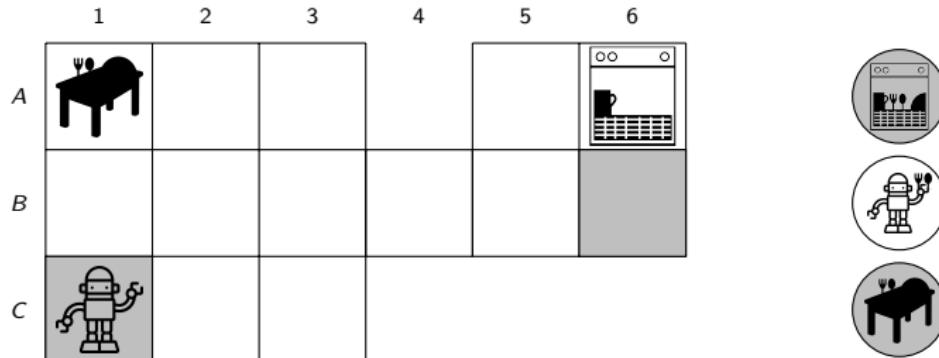
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- $\text{robot-at} = \text{C1}$ ,  $\text{dishes-at} = \text{Table}$  (initial state)

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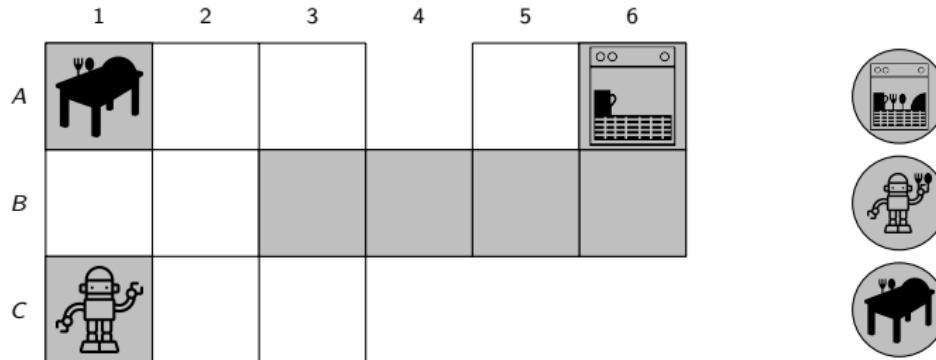
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- $\text{robot-at} = \text{C1}$ ,  $\text{dishes-at} = \text{Table}$  (initial state)
- $\text{robot-at} = \text{B6}$ ,  $\text{dishes-at} = \text{Dishwasher}$  (goal state)

## Fact Landmarks: Example

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- $\text{robot-at} = \text{A1}$ ,  $\text{robot-at} = \text{B3}$ ,  $\text{robot-at} = \text{B4}$ ,  
 $\text{robot-at} = \text{B5}$ ,  $\text{robot-at} = \text{A6}$ ,  $\text{dishes-at} = \text{Robot}$

# Constraints

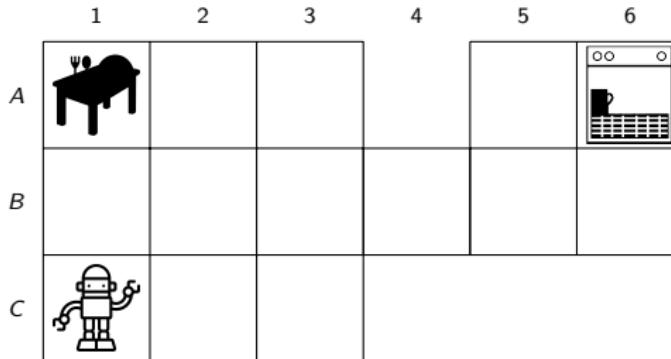
Some heuristics exploit **constraints** that describe something that holds in every solution of the task.

For instance, every solution is such that

- a variable takes some value in at least one visited state.  
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- an action must be applied.  
(an **action landmark** constraint)

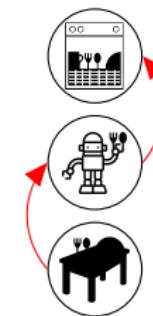
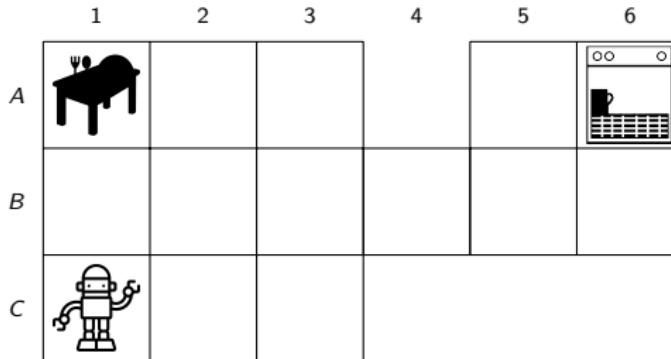
# Action Landmarks: Example

Which actions must be applied in every solution?



# Action Landmarks: Example

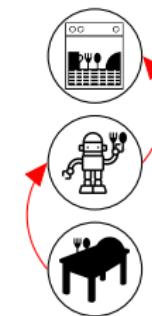
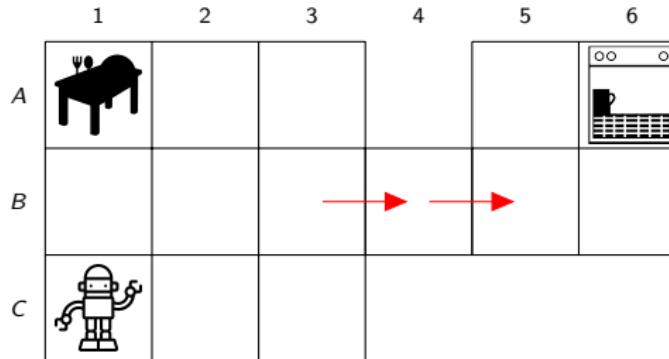
Which actions must be applied in every solution?



- pickup
- load

# Action Landmarks: Example

Which actions must be applied in every solution?



- pickup
- load
- move-B3-B4
- move-B4-B5

# Constraints

Some heuristics exploit **constraints** that describe something that holds in every solution of the task.

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# Constraints

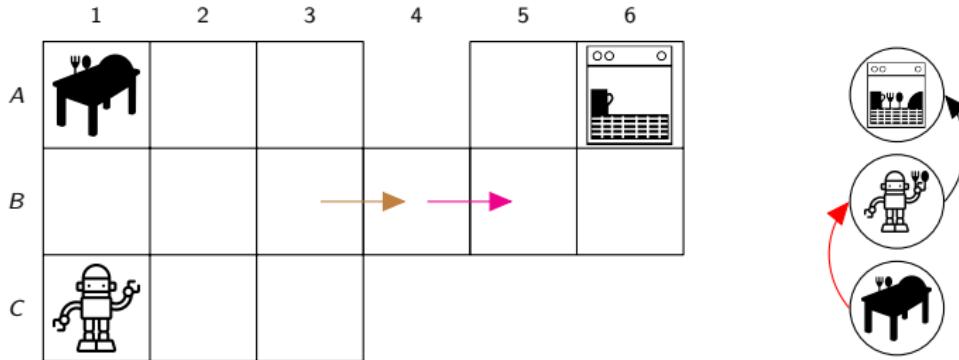
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For instance, every solution is such that

- a variable takes some **value** in at least one visited state.  
(a **fact landmark** constraint)
- at least one action from a set of actions must be applied.  
(a **disjunctive action landmark** constraint)

# Disjunctive Action Landmarks: Example

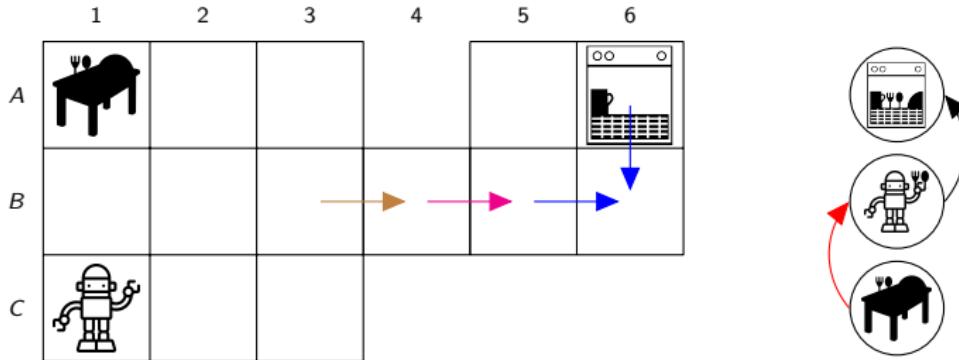
Which set of actions is such that at least one must be applied?



- {pickup}
- {load}
- {move-B3-B4}
- {move-B4-B5}

# Disjunctive Action Landmarks: Example

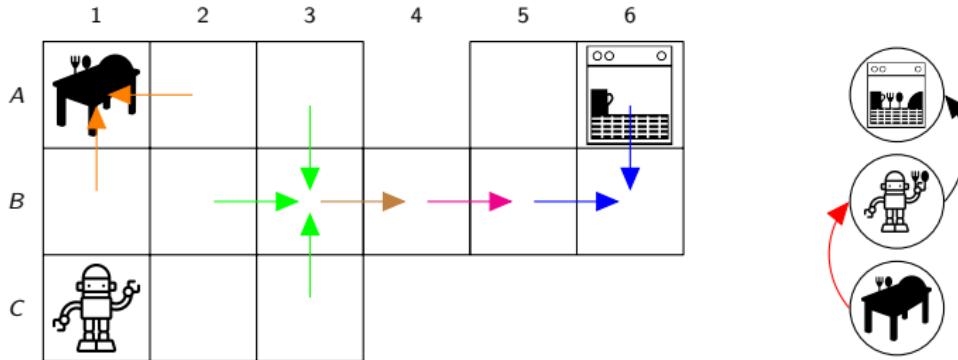
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- {pickup}
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- {move-A6-B6, move-B5-B6}

# Disjunctive Action Landmarks: Example

Which set of actions is such that at least one must be applied?



- {pickup}
- {load}
- {move-B3-B4}
- {move-B4-B5}
- {move-A6-B6, move-B5-B6}
- {move-A3-B3, move-B2-B3, move-C3-B3}
- {move-B1-A1, move-A2-A1}
- ...

# Constraints

Some heuristics exploit **constraints** that describe something that holds in every solution of the task.

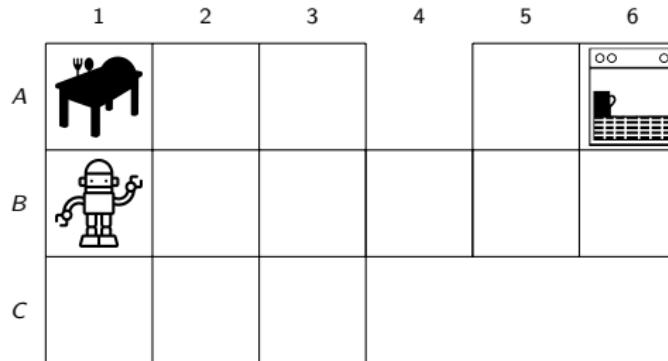
For instance, every solution is such that

- a variable takes some value in at least one visited state.  
(a **fact landmark** constraint)
- at least one action from a set of actions must be applied.  
(a **disjunctive action landmark** constraint)
- fact consumption and production is “balanced” .  
(a **network flow** constraint)

# Network Flow: Example

Consider the fact  $\text{robot-at} = B1$ .

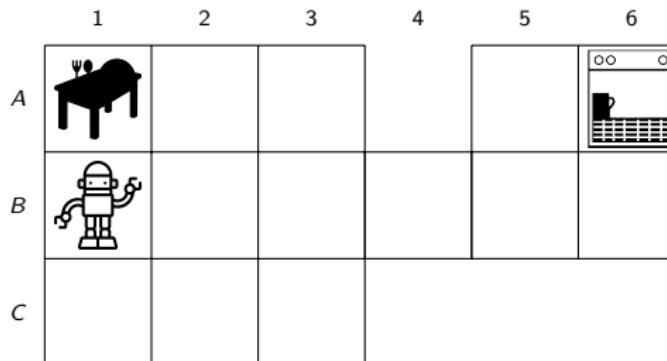
How often are actions used that enter this cell?



## Network Flow: Example

Consider the fact  $\text{robot-at} = B1$ .

How often are actions used that enter this cell?



Answer: as often as actions that leave this cell

If  $\text{Count}_o$  denotes how often operator  $o$  is applied, we have:

$$\begin{aligned}\text{Count}_{\text{move-A1-B1}} + \text{Count}_{\text{move-B2-B1}} + \text{Count}_{\text{move-C1-B1}} = \\ \text{Count}_{\text{move-B1-A1}} + \text{Count}_{\text{move-B1-B2}} + \text{Count}_{\text{move-B1-C1}}\end{aligned}$$

# Multiple Heuristics

# Combining Admissible Heuristics Admissibly

Major ideas to combine heuristics admissibly:

- maximize
- canoncial heuristic (for abstractions)
- **minimum hitting set** (for landmarks)
- **cost partitioning**
- **operator counting**

Often computed as solution to a **(integer) linear program**.

# Combining Heuristics Admissibly: Example

## Example

Consider an FDR planning task  $\langle V, I, \{o_1, o_2, o_3, o_4\}, \gamma \rangle$  with  $V = \{v_1, v_2, v_3\}$  with  $\text{dom}(v_1) = \{A, B\}$  and  $\text{dom}(v_2) = \text{dom}(v_3) = \{A, B, C\}$ ,  $I = \{v_1 \mapsto A, v_2 \mapsto A, v_3 \mapsto A\}$ ,

$$o_1 = \langle v_1 = A, v_1 := B, 1 \rangle$$

$$o_2 = \langle v_2 = A \wedge v_3 = A, v_2 := B \wedge v_3 := B, 1 \rangle$$

$$o_3 = \langle v_2 = B, v_2 := C, 1 \rangle$$

$$o_4 = \langle v_3 = B, v_3 := C, 1 \rangle$$

and  $\gamma = (v_1 = B) \wedge (v_2 = C) \wedge (v_3 = C)$ .

Let  $\mathcal{C}$  be the pattern collection that contains all atomic projections. What is the canonical heuristic function  $h^{\mathcal{C}}$ ?

# Combining Heuristics Admissibly: Example

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$$o_1 = \langle v_1 = A, v_1 := B, 1 \rangle$$

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and  $\gamma = (v_1 = B) \wedge (v_2 = C) \wedge (v_3 = C)$ .

Let  $\mathcal{C}$  be the pattern collection that contains all atomic projections. What is the canonical heuristic function  $h^{\mathcal{C}}$ ?

**Answer:** Let  $h_i := h^{v_i}$ . Then  $h^{\mathcal{C}} = \max \{h_1 + h_2, h_1 + h_3\}$ .

## Reminder: Orthogonality and Additivity

Why can we add  $h_1$  and  $h_2$  ( $h_1$  and  $h_3$ ) admissibly?

### Theorem (Additivity for Orthogonal Abstractions)

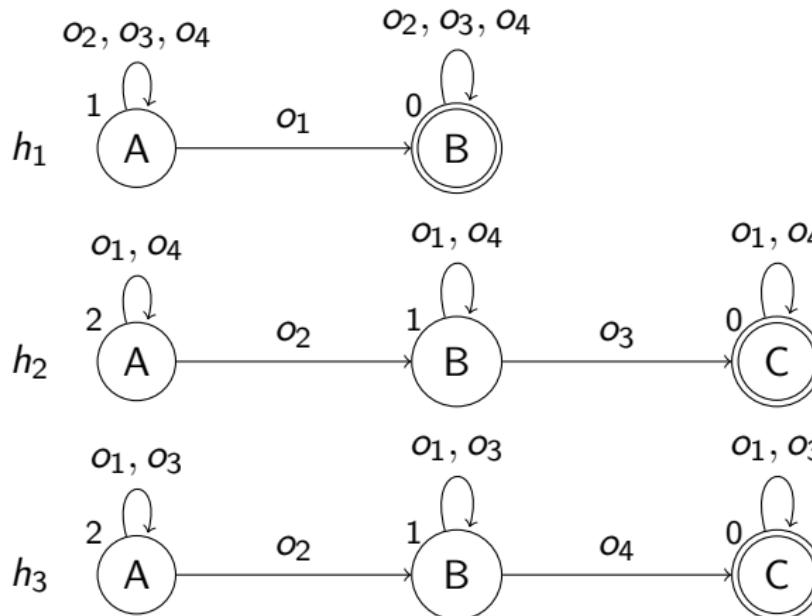
Let  $h^{\alpha_1}, \dots, h^{\alpha_n}$  be abstraction heuristics of the same transition system such that  $\alpha_i$  and  $\alpha_j$  are orthogonal for all  $i \neq j$ .

Then  $\sum_{i=1}^n h^{\alpha_i}$  is a safe, goal-aware, admissible and consistent heuristic for  $\Pi$ .

Consistency proof exploits that **every concrete transition** induces state-changing transition in **at most one abstraction**.

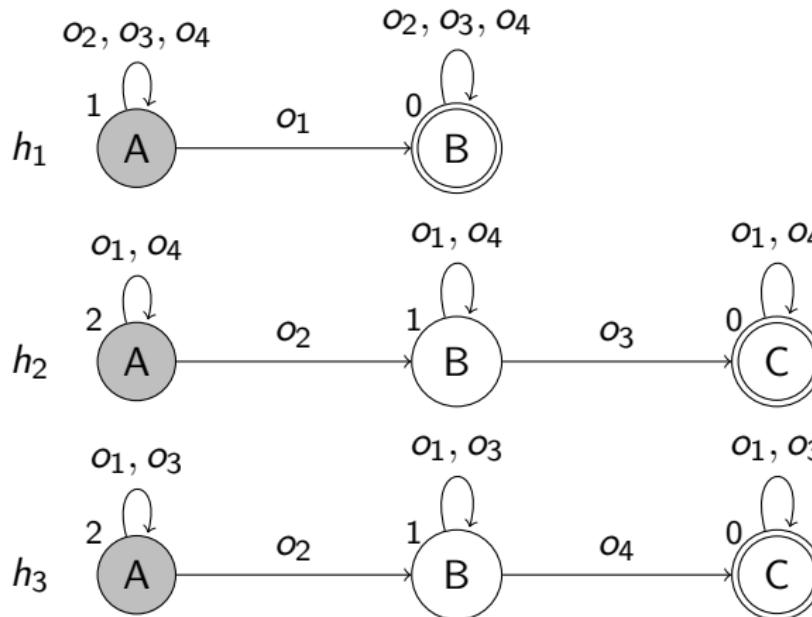
# Combining Heuristics Admissibly: Example

Let  $h = h_1 + h_2 + h_3$ . Where is consistency violated?



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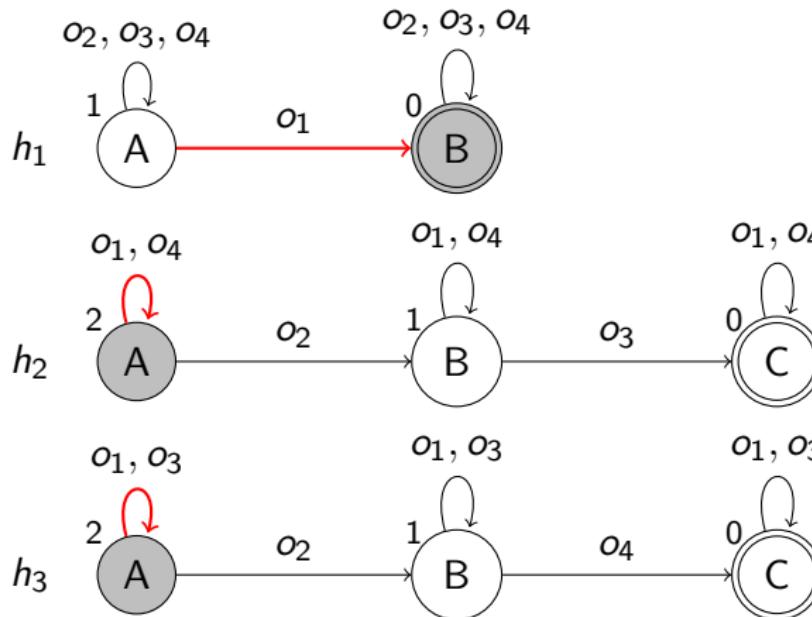
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Consider solution  $\langle o_1, o_2, o_3, o_4 \rangle$

# Combining Heuristics Admissibly: Example

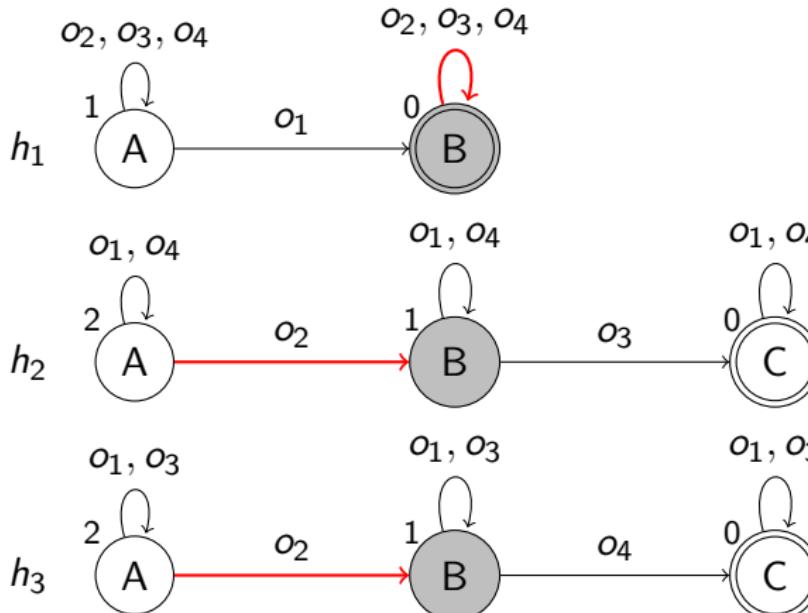
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Here:

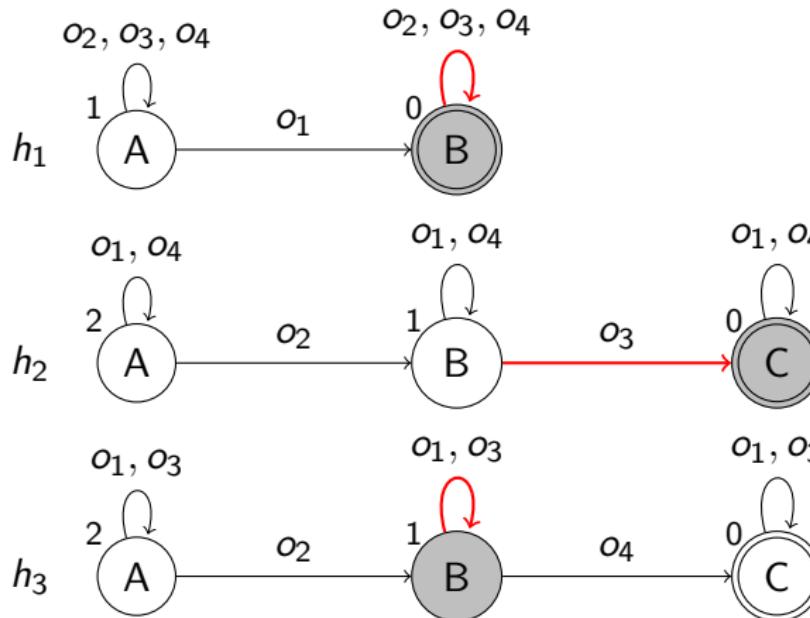
$$h(BAA) = 4$$
$$h(BBB) = 2$$

$h_2$  and  $h_3$   
not additive  
because of  $o_2$

Consider solution  $\langle o_1, o_2, o_3, o_4 \rangle$

# Combining Heuristics Admissibly: Example

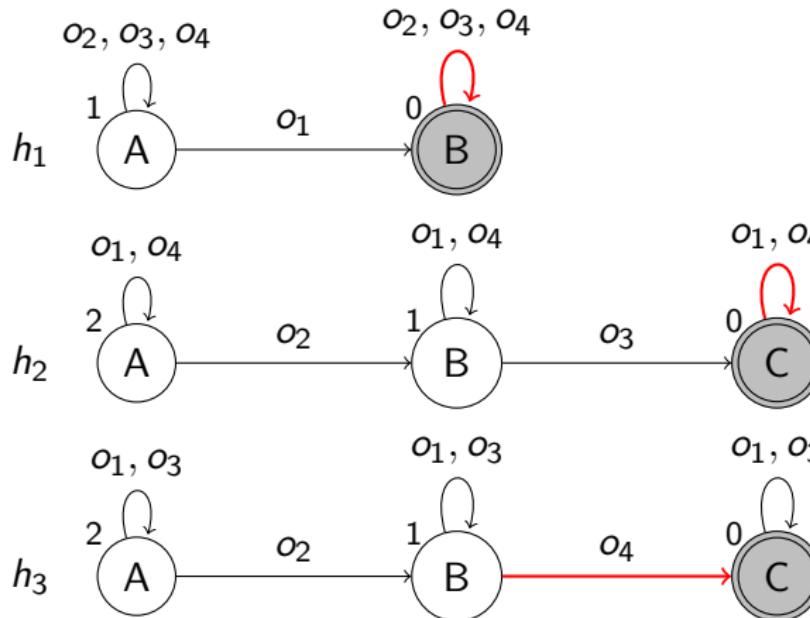
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Consider solution  $\langle o_1, o_2, \textcolor{red}{o_3}, o_4 \rangle$

# Combining Heuristics Admissibly: Example

Let  $h = h_1 + h_2 + h_3$ . Where is consistency violated?



Consider solution  $\langle o_1, o_2, o_3, o_4 \rangle$

## Inconsistency of $h_2$ and $h_3$

The reason that  $h_2$  and  $h_3$  are not additive is because the cost of  $o_2$  is considered in both.

Is there anything we can do about this?

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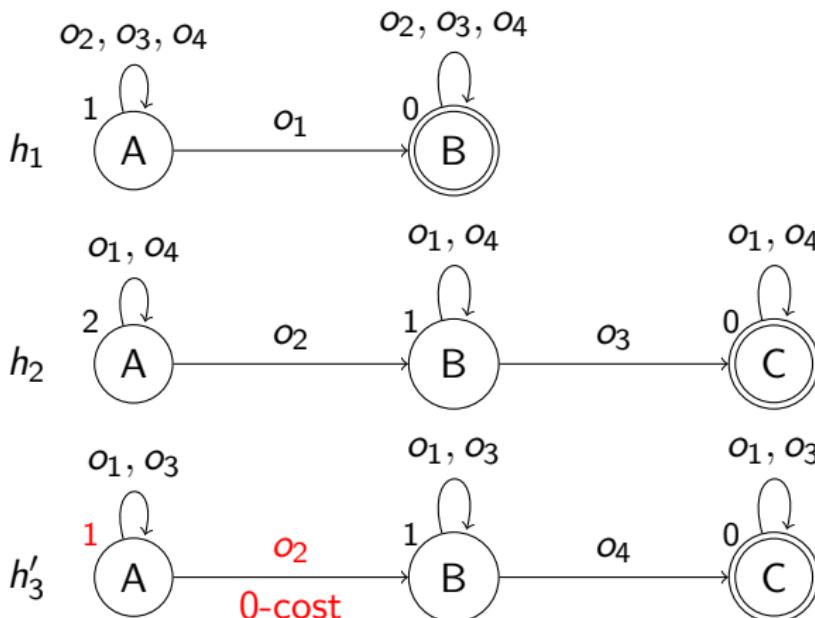
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Is there anything we can do about this?

**Solution:** We can ignore the cost of  $o_2$  in one heuristic by setting its cost to 0 (e.g.,  $cost_3(o_2) = 0$ ).

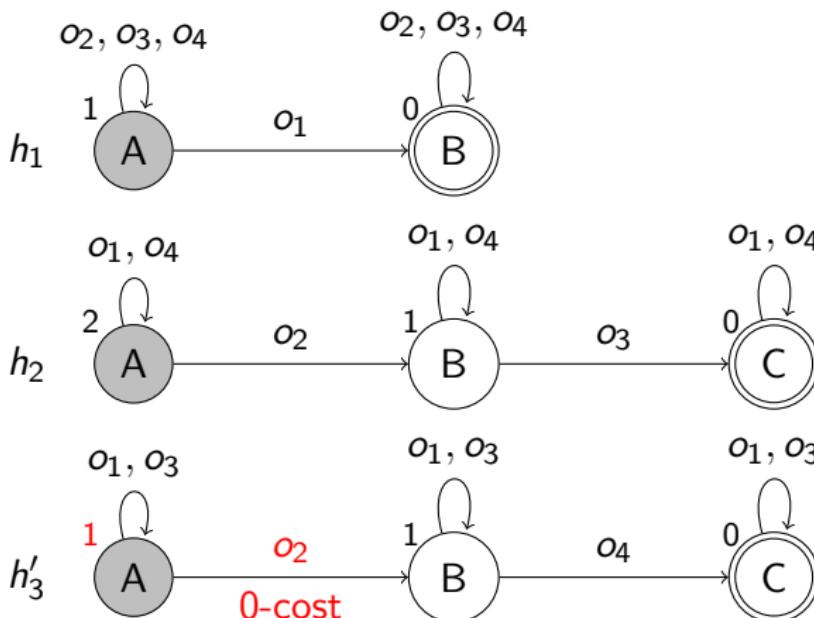
## Combining Heuristics Admissibly: Example

Let  $h' = h_1 + h_2 + h'_3$ , where  $h'_3 = h^{v_3}$  assuming  $cost_3(o_2) = 0$ .



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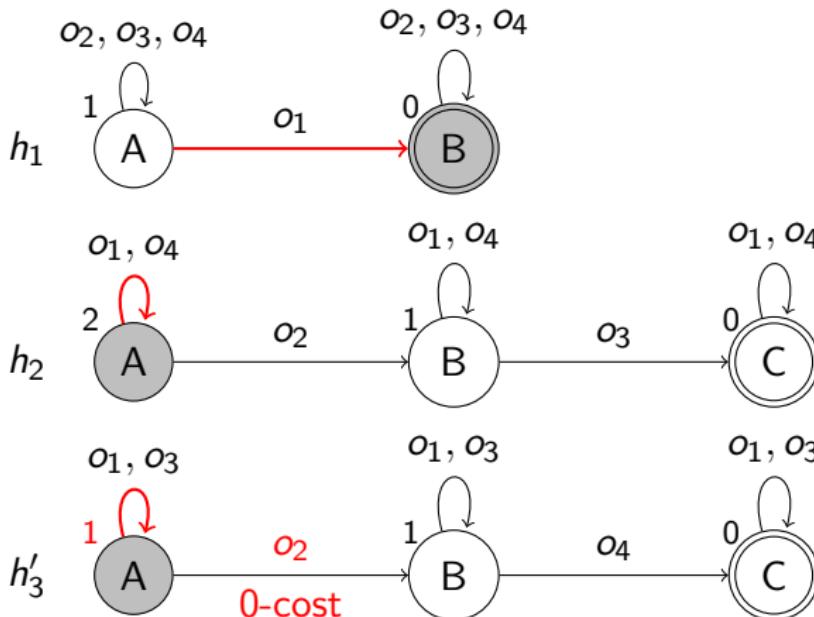
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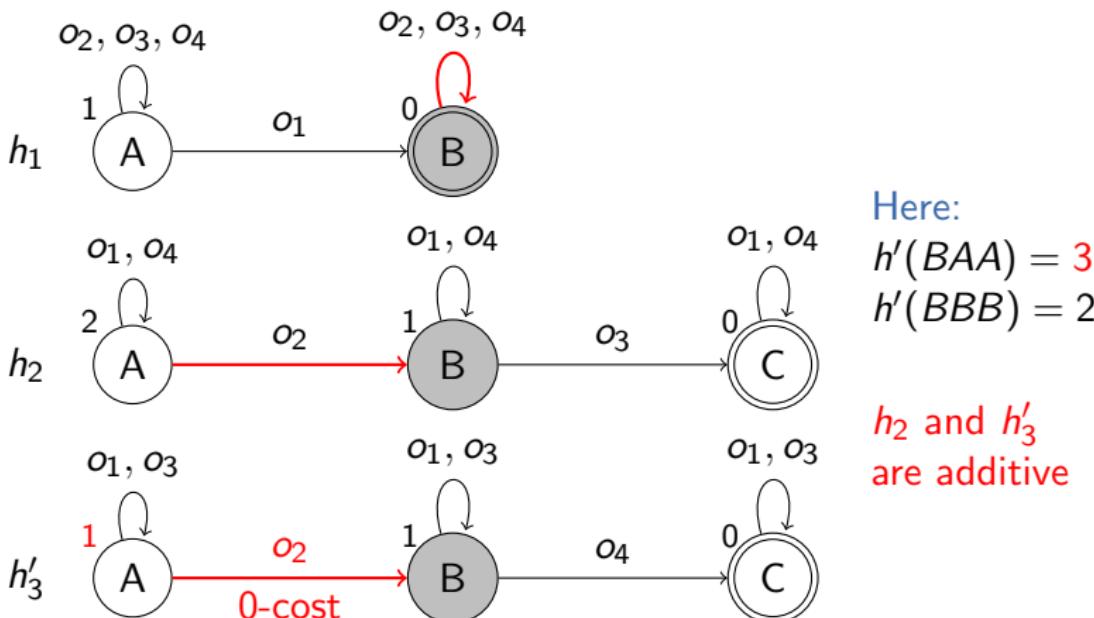
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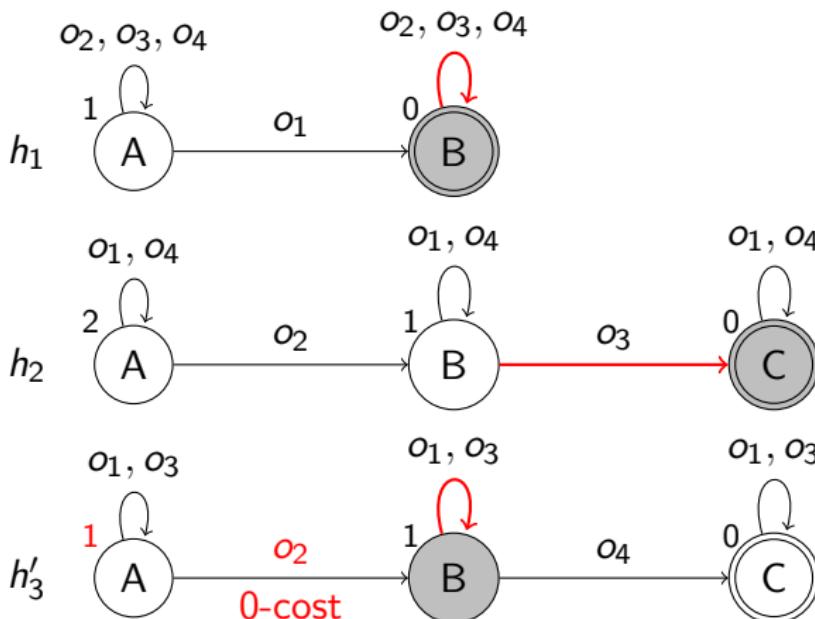
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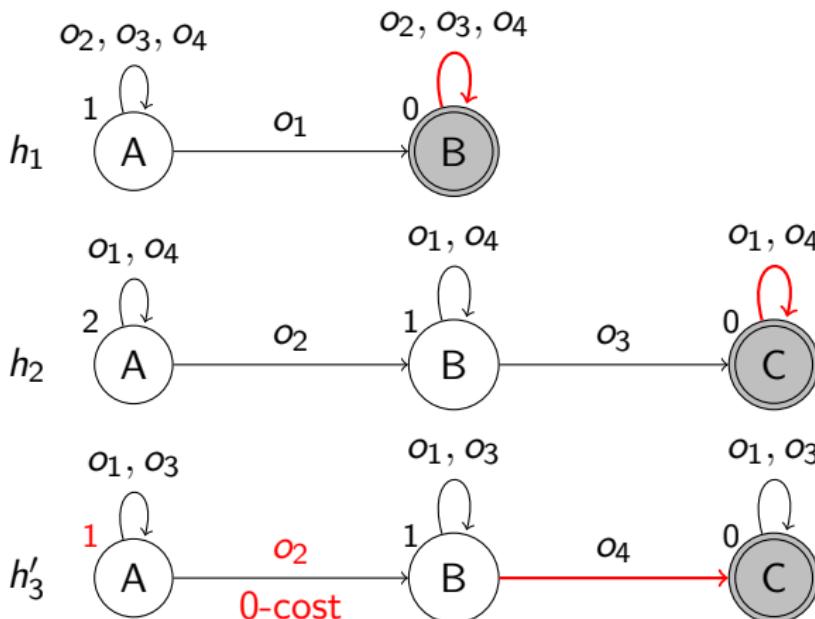
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Let  $h' = h_1 + h_2 + h'_3$ , where  $h'_3 = h^{v_3}$  assuming  $cost_3(o_2) = 0$ .



Consider solution  $\langle o_1, o_2, o_3, o_4 \rangle$

## Cost partitioning

Using the cost of every operator only in one heuristic is called a **zero-one cost partitioning**.

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Using the cost of every operator only in one heuristic is called a **zero-one cost partitioning**.

More generally, heuristics are additive if all operator costs are distributed in a way that the sum of the individual costs is no larger than the cost of the operator.

This can also be expressed as a constraint,  
the **cost partitioning constraint**:

$$\sum_{i=1}^n cost_i(o) \leq cost(o) \text{ for all } o \in O$$

(more details later)

Constraint-based Heuristics  
oooooooooooo

Multiple Heuristics  
oooooooo

Summary  
●○

# Summary

# Summary

- Landmarks and network flows are **constraints** that describe something that holds in every solution of the task.
- Heuristics can be summed up admissibly if the **cost partitioning constraint** is satisfied.