

# Discrete Mathematics in Computer Science

## B7. Operations on Relations

Malte Helmert, Gabriele Röger

University of Basel

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## Relations: Recap

- ▶ A **relation over sets**  $S_1, \dots, S_n$  is a set  $R \subseteq S_1 \times \dots \times S_n$ .
- ▶ A **binary** relation is a relation over two sets.
- ▶ A **homogeneous** relation  $R$  over set  $S$  is a binary relation  $R \subseteq S \times S$ .

## Set Operations

- ▶ Relations are sets of tuples, so we can build their union, intersection, complement, . . . .
- ▶ Let  $R$  be a relation over  $S_1, \dots, S_n$  and  $R'$  a relation over  $S'_1, \dots, S'_n$ . Then  $R \cup R'$  is a relation over  $S_1 \cup S'_1, \dots, S_n \cup S'_n$ .  
With the standard relations  $<, =$  and  $\leq$  for  $\mathbb{N}_0$ ,  
relation  $\leq$  corresponds to the union of relations  $<$  and  $=$ .
- ▶ Let  $R$  and  $R'$  be relations over  $n$  sets.  
Then  $R \cap R'$  is a relation.  
**Over which sets?**  
With the standard relations  $\leq, =$  and  $\geq$  for  $\mathbb{N}_0$ ,  
relation  $=$  corresponds to the intersection of  $\leq$  and  $\geq$ .
- ▶ If  $R$  is a relation over  $S_1, \dots, S_n$   
then so is the **complementary relation**  $\bar{R} = (S_1 \times \dots \times S_n) \setminus R$ .  
With the standard relations for  $\mathbb{N}_0$ , relation  $=$  is the  
complementary relation of  $\neq$  and  $>$  the one of  $\leq$ .

## Inverse of a Relation

### Definition

Let  $R \subseteq A \times B$  be a binary relation over  $A$  and  $B$ .

The **inverse relation** of  $R$  is the relation  $R^{-1} \subseteq B \times A$  given by  
 $R^{-1} = \{(b, a) \mid (a, b) \in R\}$ .

- ▶ The inverse of the  $<$  relation over  $\mathbb{N}_0$  is the  $>$  relation.
- ▶ Relation  $R$  with  $xRy$  iff person  $x$  has a key for  $y$ .  
Inverse:  $Q$  with  $aQb$  iff lock  $a$  can be opened by person  $b$ .

## Composition of Relations

### Definition (Composition of relations)

Let  $R_1$  be a relation over  $A$  and  $B$  and  
 $R_2$  be a relation over  $B$  and  $C$ .

The **composition** of  $R_1$  and  $R_2$  is the relation  $R_2 \circ R_1$  with:

$$R_2 \circ R_1 = \{(a, c) \mid \text{there is a } b \in B \text{ with} \\ (a, b) \in R_1 \text{ and } (b, c) \in R_2\}$$

How can we illustrate this graphically?

## Composition is Associative

### Theorem (Associativity of composition)

Let  $S_1, \dots, S_4$  be sets and  $R_1, R_2, R_3$  relations with  $R_i \subseteq S_i \times S_{i+1}$ .  
Then

$$R_3 \circ (R_2 \circ R_1) = (R_3 \circ R_2) \circ R_1.$$

### Proof.

It holds that  $(x_1, x_4) \in R_3 \circ (R_2 \circ R_1)$  iff there is an  $x_3$  with  
 $(x_1, x_3) \in R_2 \circ R_1$  and  $(x_3, x_4) \in R_3$ .

As  $(x_1, x_3) \in R_2 \circ R_1$  iff there is an  $x_2$  with  $(x_1, x_2) \in R_1$  and  
 $(x_2, x_3) \in R_2$ , we have overall that  $(x_1, x_4) \in R_3 \circ (R_2 \circ R_1)$  iff  
there are  $x_2, x_3$  with  $(x_1, x_2) \in R_1$ ,  $(x_2, x_3) \in R_2$  and  $(x_3, x_4) \in R_3$ .

This is the case iff there is an  $x_2$  with  $(x_1, x_2) \in R_1$  and  
 $(x_2, x_4) \in R_3 \circ R_2$ , which holds iff  $(x_1, x_4) \in (R_3 \circ R_2) \circ R_1$ .  $\square$

## Transitive Closure

### Definition (Transitive closure)

The **transitive closure**  $R^+$  of a relation  $R$  over set  $S$  is the **smallest relation over  $S$  that is transitive and has  $R$  as a subset.**

The transitive closure always exists. **Why?**

**Example:** If  $aRb$  specifies that block  $a$  lies on block  $b$ , what does  $R^+$  express?

## Transitive Closure II

Define the  $i$ -th power of a homogeneous relation  $R$  as

$$\begin{aligned} R^1 &= R && \text{if } i = 1 \text{ and} \\ R^i &= R \circ R^{i-1} && \text{for } i > 1 \end{aligned}$$

### Theorem

Let  $R$  be a relation over set  $S$ . Then  $R^+ = \bigcup_{i=1}^{\infty} R^i$ .

Without proof.

## Other Operators

- ▶ There are many more operators, also for general relations.
- ▶ Highly relevant for **queries over relational databases**.
- ▶ For example, **join operators** combine relations based on common entries.
- ▶ Example for a **natural join**:

Employee			Dept		Employee $\bowtie$ Dept			
Name	EmpId	DeptName	DeptName	Manager	Name	EmpId	DeptName	Manager
Harry	3415	Finance	Finance	George	Harry	3415	Finance	George
Sally	2241	Sales	Sales	Harriet	Sally	2241	Sales	Harriet
George	3401	Finance	Production	Charles	George	3401	Finance	George
Harriet	2202	Sales			Harriet	2202	Sales	Harriet
Mary	1257	Human Resources						

(Source: Wikipedia)

## Summary

- ▶ Relations: general, **binary**, **homogeneous**
- ▶ Properties: **reflexivity**, **symmetry**, **transitivity** (and related properties)
- ▶ Special relations: **equivalence relations**, **order relations**
- ▶ Operations: **inverse**, **composition**, **transitive closure**