Discrete Mathematics in Computer Science

B3. Cantor's Theorem

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B3.1 Cantor's Theorem

Countable Sets

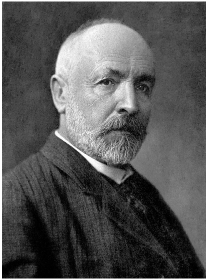
We already know:

- The cardinality of $\mathbb{N}_0$ is $\aleph_0$.
- All sets with cardinality $\aleph_0$ are called countably infinite.
- A countable set is finite or countably infinite.
- Every subset of a countable set is countable.
- The union of countably many countable sets is countable.

These questions were still open:

- Do all infinite sets have the same cardinality?
- Does the power set of infinite set S have the same cardinality as S?

Georg Cantor



- German mathematician (1845–1918)
- Proved that the rational numbers are countable.
- Proved that the real numbers are not countable.
- Cantor's Theorem: For every set S it holds that $|S| < |\mathcal{P}(S)|$.

Our Plan

- Understand Cantor's theorem
- Understand an important theoretical implication for computer science

Cantor's Diagonal Argument Illustrated on a Finite Set

$S = \{a, b, c\}$.

Consider an arbitrary injective function from S to $\mathcal{P}(S)$.
For example:

	a	b	c	
a	1	0	1	a mapped to $\{a, c\}$
b	1	1	0	b mapped to $\{a, b\}$
c	0	1	0	c mapped to $\{b\}$
	0	0	1	nothing was mapped to $\{c\}$.

We can identify an "unused" element of $\mathcal{P}(S)$.
Complement the entries on the main diagonal.

Works with every injective function from S to $\mathcal{P}(S)$.
$\rightarrow$ there cannot be a bijection from S to $\mathcal{P}(S)$.

Cantor's Diagonal Argument on a Countably Infinite Set

$S = \mathbb{N}_0$.

Consider an arbitrary injective function from $\mathbb{N}_0$ to $\mathcal{P}(\mathbb{N}_0)$.
For example:

	0	1	2	3	4	...
0	1	0	1	0	1	...
1	1	1	0	1	0	...
2	0	1	0	1	0	...
3	1	1	0	0	0	...
4	1	1	0	1	1	...
⋮	⋮	⋮	⋮	⋮	⋮	⋱
	0	0	1	1	0	...

Complementing the entries on the main diagonal again results in an "unused" element of $\mathcal{P}(\mathbb{N}_0)$.

Cantor's Theorem

Theorem (Cantor's Theorem)

For every set S it holds that $|S| < |\mathcal{P}(S)|$.

Proof.

We need to show that

1. There is an injective function from S to $\mathcal{P}(S)$.
2. There is no bijection from S to $\mathcal{P}(S)$.

For 1, consider function $f : S \to \mathcal{P}(S)$ with $f(x) = \{x\}$. Each element of S is paired with a unique element of $\mathcal{P}(S)$. ...

Cantor's Theorem

Proof (continued).

For 2, we show for every injective function $f : S \to \mathcal{P}(S)$ that it is not a bijection from S to $\mathcal{P}(S)$.
This is sufficient because every bijection is injective.

Let f be an arbitrary injective function with $f : S \to \mathcal{P}(S)$.

Consider $M = \{x \mid x \in S, x \notin f(x)\}$.

For every $x \in S$ it holds that $f(x) \neq M$ because
$x \in f(x)$ iff not $x \notin f(x)$ iff not $x \in M$ iff $x \notin M$.

Hence, there is no $x \in S$ with $f(x) = M$. As $M \in \mathcal{P}(S)$ this implies that f is not a bijection from S to $\mathcal{P}(S)$. □

B3.2 Consequences of Cantor's Theorem

Infinite Sets can Have Different Cardinalities

There are infinitely many different cardinalities of infinite sets:

- $|\mathbb{N}_0| < |\mathcal{P}(\mathbb{N}_0))| < |\mathcal{P}(\mathcal{P}(\mathbb{N}_0)))| < \dots$
- $|\mathbb{N}_0| = \aleph_0 = \beth_0$
- $|\mathcal{P}(\mathbb{N}_0)| = \beth_1 (= |\mathbb{R}|)$
- $|\mathcal{P}(\mathcal{P}(\mathbb{N}_0))| = \beth_2$
- $\dots$

Existence of Unsolvable Problems

There are more problems in computer science
than there are programs to solve them.

There are problems that cannot be solved by a computer program!

Why can we say so?

Decision Problems

"Intuitive Definition:" Decision Problem

A decision problem is a Yes-No question of the form "Does the given input have a certain property?"

- "Does the given binary tree have more than three leaves?"
- "Is the given integer odd?"
- "Given a train schedule, is there a connection from Basel to Belinzona that takes at most 2.5 hours?"

- Input can be encoded as some finite string.
- Problem can also be represented as the (possibly infinite) set of all input strings where the answer is "yes".
- A computer program solves a decision problem if it terminates on every input and returns the correct answer.

More Problems than Programs I

- A computer program is given by a finite string.
- A decision problem corresponds to a set of strings.

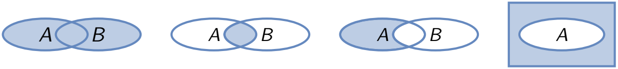
More Problems than Programs II

- Consider an arbitrary finite set of symbols (an alphabet) Σ.
- You can think of $\Sigma = \{0, 1\}$ as internally computers operate on binary representation.
- Let S be the set of all finite strings made from symbols in Σ.
- There are at most $|S|$ computer programs with this alphabet.
- There are at least $|\mathcal{P}(S)|$ problems with this alphabet.
  - every subset of S corresponds to a separate decision problem
- By Cantor's theorem $|S| < |\mathcal{P}(S)|$, so there are more problems than programs.

B3.3 Sets: Summary

Summary

- A set is an unordered collection of distinct objects.
- Set operations: union, intersection, set difference, complement



  - Commutativity, associativity and distributivity of union and intersection
  - De Morgan's law: $\overline{A \cup B} = \overline{A} \cap \overline{B}$ and $\overline{A \cap B} = \overline{A} \cup \overline{B}$.
- The cardinality measures the "size" of a set.
  - For finite sets, the cardinality equals the number of elements.
  - All sets with the same cardinality as $\mathbb{N}_0$ are countably infinite.
  - All sets with cardinality $\leq |\mathbb{N}_0|$ are countable.
- The power set $\mathcal{P}(S)$ of set S is the set of all subsets of S.
  - For finite sets S it holds that $|\mathcal{P}(S)| = 2^{|S|}$.
  - For all sets S it holds that $|S| < |\mathcal{P}(S)|$.