Discrete Mathematics in Computer Science B7. Operations on Relations

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- A relation over sets $S_{1}, \ldots, S_{n}$ is a set $R \subseteq S_{1} \times \cdots \times S_{n}$.
- A binary relation is a relation over two sets.
- A homogeneous relation $R$ over set $S$ is a binary relation $R \subseteq S \times S$.
- Relations are sets of tuples, so we can build their union, intersection, complement,
- Let $R$ be a relation over $S_{1}, \ldots, S_{n}$ and $R^{\prime}$ a relation over $S_{1}^{\prime}, \ldots, S_{n}^{\prime}$. Then $R \cup R^{\prime}$ is a relation over $S_{1} \cup S_{1}^{\prime}, \ldots, S_{n} \cup S_{n}^{\prime}$. With the standard relations $<,=$ and $\leq$ for $\mathbb{N}_{0}$, relation $\leq$ corresponds to the union of relations $<$ and $=$.
- Let $R$ and $R^{\prime}$ be relations over $n$ sets.

Then $R \cap R^{\prime}$ is a relation.
Over which sets?
With the standard relations $\leq,=$ and $\geq$ for $\mathbb{N}_{0}$,
relation $=$ corresponds to the intersection of $\leq$ and $\geq$.

- If $R$ is a relation over $S_{1}, \ldots, S_{n}$
then so is the complementary relation $\bar{R}=\left(S_{1} \times \cdots \times S_{n}\right) \backslash R$.
With the standard relations for $\mathbb{N}_{0}$, relation $=$ is the
complementary relation of $\neq$ and $>$ the one of $\leq$.


## Definition

Let $R \subseteq A \times B$ be a binary relation over $A$ and $B$.
The inverse relation of $R$ is the relation $R^{-1} \subseteq B \times A$ given by $R^{-1}=\{(b, a) \mid(a, b) \in R\}$.

- The inverse of the $<$ relation over $\mathbb{N}_{0}$ is the $>$ relation.
- Relation $R$ with $x R y$ iff person $x$ has a key for $y$. Inverse: $Q$ with $a Q b$ iff lock $a$ can be openened by person $b$.
B7. Operations on Relations
Composition of Relations

| Definition (Composition of relations) |
| :--- |
| Let $R_{1}$ be a relation over $A$ and $B$ and |
| $R_{2}$ be a relation over $B$ and $C$. |
| The composition of $R_{1}$ and $R_{2}$ is the relation $R_{2} \circ R_{1}$ with: |
| $\qquad$$R_{2} \circ R_{1}=\{(a, c) \mid$ there is a $b \in B$ with <br> $(a, b) \in R_{1}$ and $\left.(b, c) \in R_{2}\right\}$ |

## Composition is Associative

Theorem (Associativity of composition)
Let $S_{1}, \ldots, S_{4}$ be sets and $R_{1}, R_{2}, R_{3}$ relations with $R_{i} \subseteq S_{i} \times S_{i+1}$. Then

$$
R_{3} \circ\left(R_{2} \circ R_{1}\right)=\left(R_{3} \circ R_{2}\right) \circ R_{1}
$$

## Proof.

It holds that $\left(x_{1}, x_{4}\right) \in R_{3} \circ\left(R_{2} \circ R_{1}\right)$ iff there is an $x_{3}$ with $\left(x_{1}, x_{3}\right) \in R_{2} \circ R_{1}$ and $\left(x_{3}, x_{4}\right) \in R_{3}$.

As $\left(x_{1}, x_{3}\right) \in R_{2} \circ R_{1}$ iff there is an $x_{2}$ with $\left(x_{1}, x_{2}\right) \in R_{1}$ and $\left(x_{2}, x_{3}\right) \in R_{2}$, we have overall that $\left(x_{1}, x_{4}\right) \in R_{3} \circ\left(R_{2} \circ R_{1}\right)$ iff there are $x_{2}, x_{3}$ with $\left(x_{1}, x_{2}\right) \in R_{1},\left(x_{2}, x_{3}\right) \in R_{2}$ and $\left(x_{3}, x_{4}\right) \in R_{3}$.
This is the case iff there is an $x_{2}$ with $\left(x_{1}, x_{2}\right) \in R_{1}$ and
$\left(x_{2}, x_{4}\right) \in R_{3} \circ R_{2}$, which holds iff $\left(x_{1}, x_{4}\right) \in\left(R_{3} \circ R_{2}\right) \circ R_{1}$.

## Definition (Transitive closure)

The transitive closure $R^{*}$ of a relation $R$ over set $S$ is the smallest relation over $S$ that is transitive and has $R$ as a subset.

The transitive closure always exists. Why?
Example: If $a R b$ specifies that block $a$ lies on block $b$, what does $R^{*}$ express?

Define the $i$-th power of a homogeneous relation $R$ as

$$
\begin{array}{ll}
R^{1}=R & \text { if } i=1 \text { and } \\
R^{i}=R \circ R^{i-1} & \text { for } i>1
\end{array}
$$

Theorem
Let $R$ be a relation over set $S$. Then $R^{*}=\bigcup_{i=1}^{\infty} R^{i}$.
Without proof.

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