

Planning and Optimization

G4. Heuristics for Probabilistic Planning

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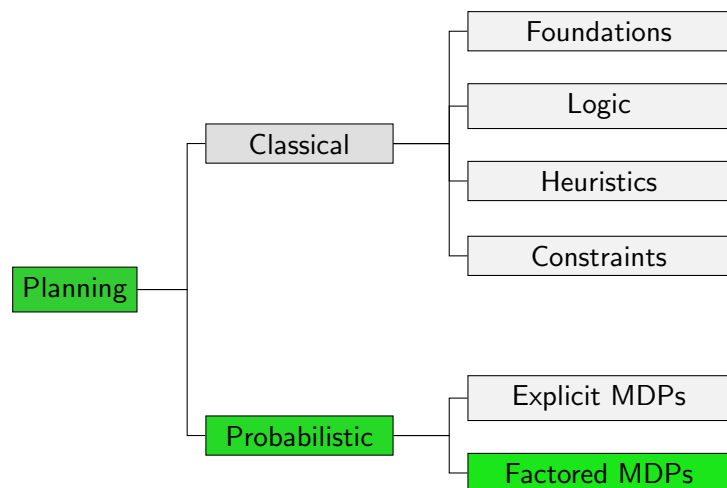
G4.1 Motivation

G4.2 Determinization

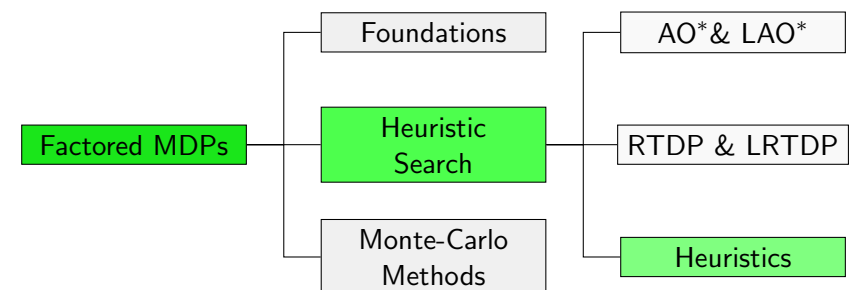
G4.3 Properties

G4.4 Summary

Content of this Course



Content of this Course: Factored MDPs



G4.1 Motivation

Heuristics

Heuristics for probabilistic planning

- ▶ most heuristics are based on **domain-specific** knowledge
- ▶ for a long time, **determinization-based** heuristics were the only notable **domain-independent** solution
- ▶ **Recent progress**: occupation measures (only covered in exercises)

G4.2 Determinization

What is a Determinization?

- ▶ Replace **SSP operators** with **deterministic** operators
- ▶ Results in **classical planning task**
- ▶ For SSPs, this is a **classical planning task**
- ▶ SSP and its determinization are **related** but **not equivalent**

How to Come up with a Determinization?

Typically, two **types** of determinization are distinguished:

- ▶ **All-outcomes determinization**
 - ▶ Create one deterministic transition to each outcome
- ▶ **Single-outcome determinization**
 - ▶ Pick **one outcome** of each probabilistic transition ...
 - ▶ ... and turn it into a deterministic transition
 - ▶ often, the **most likely** outcome is preserved
 - ▶ or one outcome is **sampled** according to its probability

All-outcomes Determinization

Definition (All-outcomes Determinization)

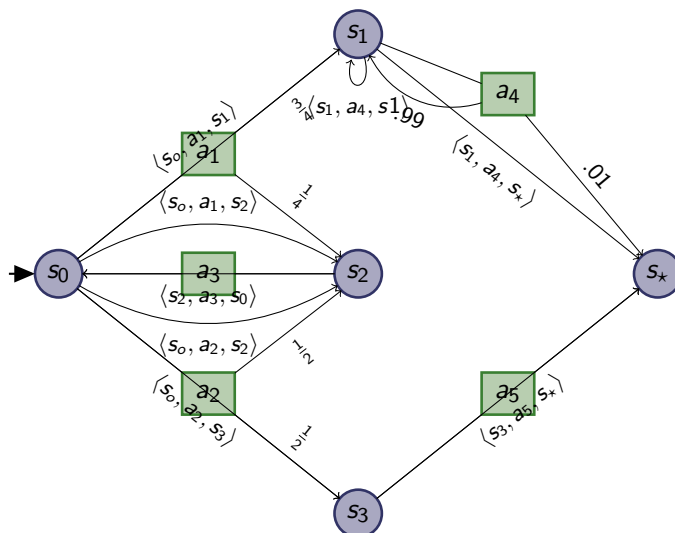
Let $\mathcal{T} = \langle S, L, c, T, s_0, S_* \rangle$ be an SSP. The **all-outcomes determinization** of \mathcal{T} is the (deterministic) transition system $\mathcal{T}^d = \langle S, L^d, c, T^d, s_0, S_* \rangle$ with

$$L^d = \{ \langle s, \ell, s' \rangle \mid s, s' \in S, \ell \in L \text{ and } T(s, \ell, s') > 0 \}$$

$$T^d = \{ \langle s, \ell^d, s' \rangle \mid \ell^d = \langle s, \ell, s' \rangle \in L^d \}.$$

All-outcomes Determinization: (Unit-cost) Example

Generate **one action for each outcome**



Single-outcome Determinization

Definition (Single-outcome Determinization)

Let $\mathcal{T} = \langle S, L, c, T, s_0, S_* \rangle$ be an SSP and let

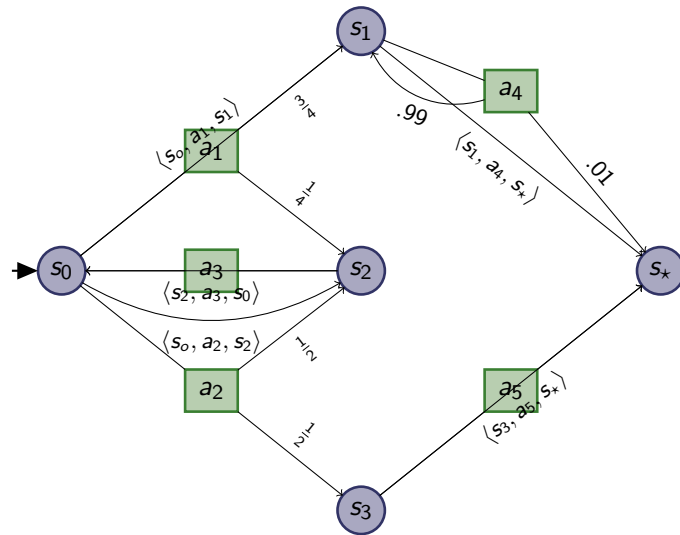
$$L^{aod} = \{ \langle s, \ell, s' \rangle \mid s, s' \in S, \ell \in L \text{ and } T(s, \ell, s') > 0 \}.$$

A (deterministic) transition system $\mathcal{T}^d = \langle S, L^d, c, T^d, s_0, S_* \rangle$ is a **single-outcome determinization** of \mathcal{T} if

- ▶ $L^d \subseteq L^{aod}$ s.t. for all $s \in S$ and $\ell \in L(s)$ there is exactly one $\langle s, \ell, s' \rangle \in L^d$ and
- ▶ $T^d = \{ \langle s, \ell^d, s' \rangle \mid \ell^d = \langle s, \ell, s' \rangle \in L^d \}$

Single-outcome Determinization: (Unit-cost) Example

Remove transition to **all outcomes but one**



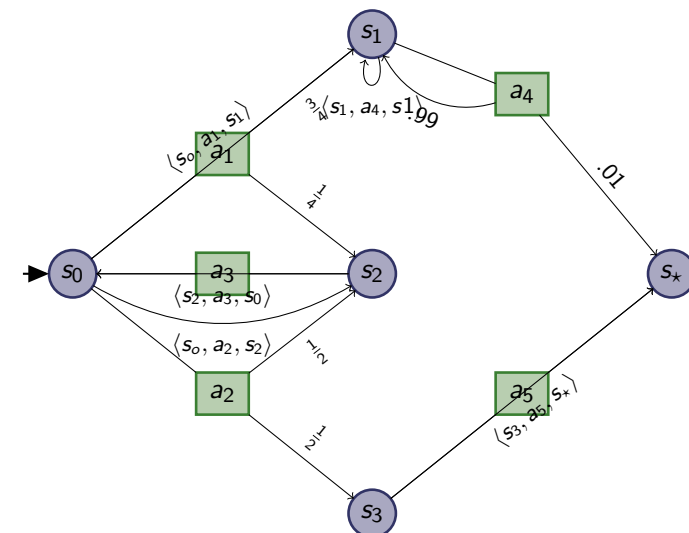
G4.3 Properties

Single-outcome Determinization: Properties

- ▶ single-outcome determinizations are **not well-suited** to be used as a heuristic:
 - ▶ can be inadmissible
 - ▶ and even unsafe

Single-outcome Determinization: Inadmissibility

Remove transition to **all outcomes but one**



Single-outcome Determinization: Properties

- ▶ single-outcome determinizations are **not well-suited** to be used as a heuristic:
 - ▶ can be inadmissible
 - ▶ and even unsafe
- ▶ often part of **domain-specific** solutions (e.g., by ignoring that some action may “fail”)
- ▶ and as part of algorithms that average over several **samples** (Chapter G5)

Min-min Heuristic

Definition (Min-min Heuristic)

Let \mathcal{T} be an SSP and let \mathcal{T}^d be the all-outcomes determinization of \mathcal{T} . The **min-min heuristic** h_{\min} maps each state $s \in S$ to the cost of the cheapest path from s to a goal state in \mathcal{T}^d .

Min-min Heuristic: Admissibility

Theorem (Admissibility)

The *min-min heuristic* is **admissible**.

Proof Sketch.

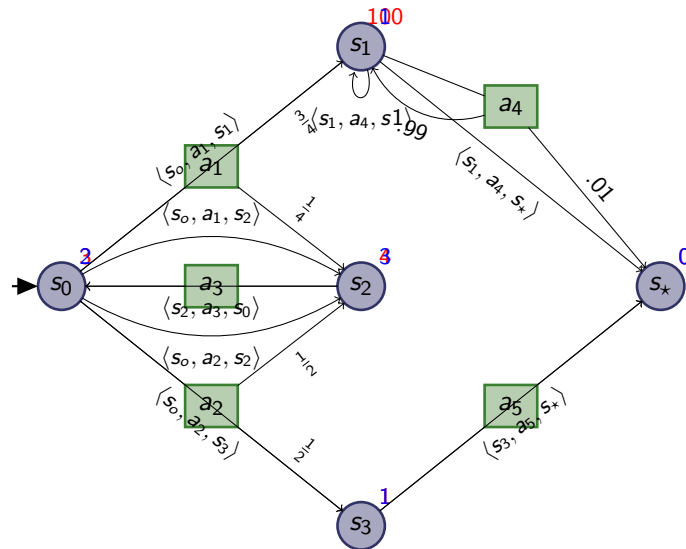
$$\begin{aligned}
 V^*(s) &\stackrel{(1)}{=} \min_{\ell \in L(s)} c(\ell) + \sum_{s' \in S} T(s, \ell, s') \cdot V^*(s') \\
 &\stackrel{(2)}{=} \min_{\ell \in L(s)} c(\ell) + \sum_{s' \in \text{succ}(s, \ell)} T(s, \ell, s') \cdot V^*(s') \\
 &\stackrel{(3)}{\geq} \min_{\ell \in L(s)} c(\ell) + \min_{s' \in \text{succ}(s, \ell)} V^*(s') = h_{\min}(s)
 \end{aligned}$$

(1) is the Bellman equation, (2) holds because $T(s, \ell, s') = 0$ for all $s' \notin \text{succ}(s, \ell)$ and (3) because the weight from more expensive outcomes is shifted to the cheapest one.

Min-min Heuristic: Properties

- ▶ The min-min heuristic is also called the **optimistic heuristic** ...
- ▶ ... because the planner may **choose its preferred outcome**
- ▶ Min-min heuristic can be **well informed** ...
- ▶ ... but can also be **utterly optimistic**

All-outcomes Determinization: Quality



Annotation: probabilistic state-values (in red) deterministic

Min-min Heuristic: Properties

- ▶ The min-min heuristic is also called the **optimistic heuristic** ...
- ▶ ... because the planner may **choose its preferred outcome**
- ▶ Min-min heuristic can be **well informed** ...
- ▶ ... but can also be **utterly optimistic**
- ▶ Min-min heuristic often **solvable in practice** even if SSP is not
- ▶ If still unsolvable: compute **classical heuristic** of all-outcomes determinization

G4.4 Summary

Summary

- ▶ Almost all heuristics in probabilistic planning are either **domain-specific** or based on a **determinization**
- ▶ A **single-outcome determinization** removes all outcomes from a transition **except for one**
- ▶ The **all-outcomes determinization** creates a deterministic transition **for each outcome**
- ▶ The **min-min heuristic** computes the shortest path in the all-outcomes determinization
- ▶ The **min-min heuristic** is admissible