

# Planning and Optimization

## G1. Heuristic Search: AO\* & LAO\* Part I

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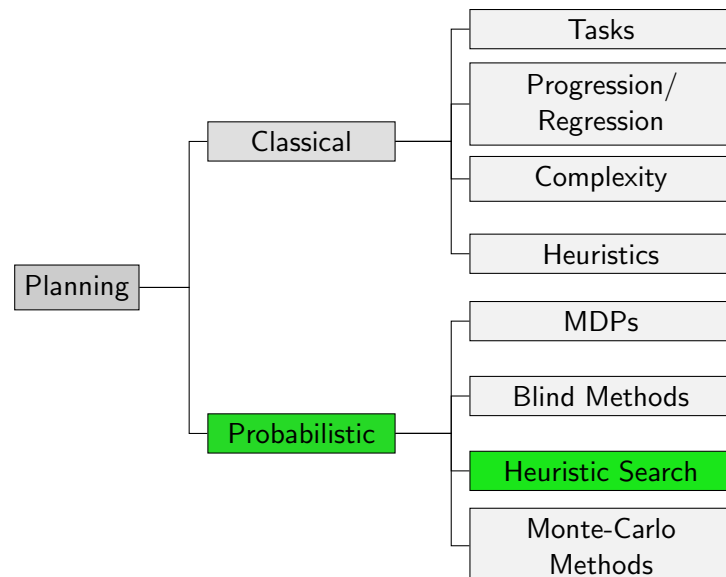
G1.1 Heuristic Search

G1.2 Motivation

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G1.4 Summary

## Content of this Course



# G1.1 Heuristic Search

## Heuristic Search: Recap

### Heuristic Search Algorithms

Heuristic search algorithms use heuristic functions to (partially or fully) determine the order of node expansion.

(From Lecture 15 of the AI course last semester)

## Best-first Search: Recap

### Best-first Search

A best-first search is a heuristic search algorithm that evaluates search nodes with an evaluation function  $f$  and always expands a node  $n$  with minimal  $f(n)$  value.

(From Lecture 15 of the AI course last semester)

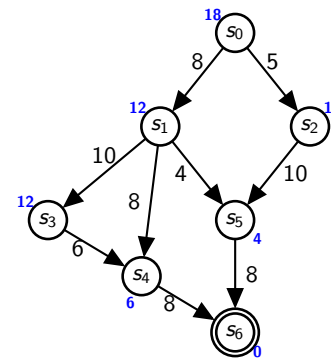
## A\* Search: Recap

### A\* Search

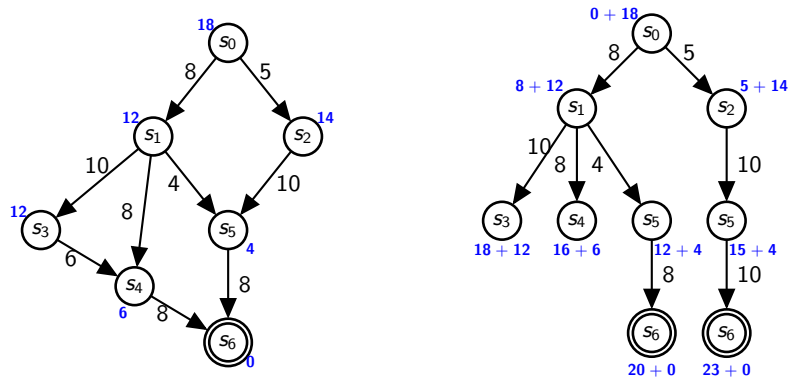
A\* is the best-first search algorithm with evaluation function  $f(n) = g(n) + h(n.state)$ .

(From Lecture 15 of the AI course last semester)

## A\* Search (With Reopening): Example



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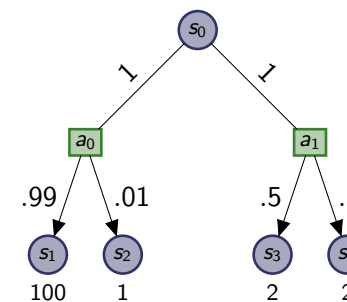
## G1.2 Motivation

## From A\* to AO\*

- ▶ Equivalent of A\* in (acyclic) probabilistic planning is **AO\***
- ▶ Even though we know A\* and foundations of probabilistic planning, the generalization is **far from straightforward**:
  - ▶ e.g., in A\*,  $g(n)$  is cost from root  $n_0$  to  $n$
  - ▶ equivalent in AO\* is **expected cost** from  $n_0$  to  $n$

## Expected Cost to Reach State

Consider the following expansion of state  $s_0$ :



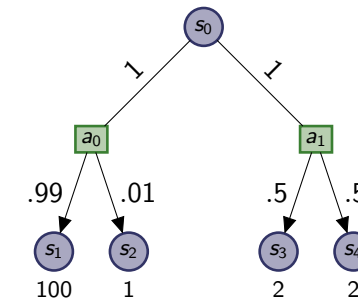
Expected cost to reach **any of the leaves** is **infinite** or undefined (neither is reached with probability 1).

## From A\* to AO\*

- ▶ Equivalent of A\* in (acyclic) probabilistic planning is **AO\***
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  - ▶ e.g., in A\*,  $g(n)$  is cost from root  $n_0$  to  $n$
  - ▶ equivalent in AO\* is **expected cost** from  $n_0$  to  $n$
  - ▶ alternative could be **expected cost from  $n_0$  to  $n$  given  $n$  is reached**

## Expected Cost to Reach State Given It Is Reached

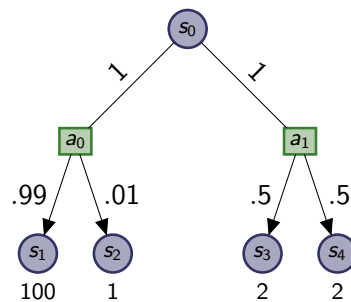
Consider the following expansion of state  $s_0$ :



Conditional probability is **misleading**:  $s_2$  would be expanded, which isn't part of the **best looking** option

## The Best Looking Action

Consider the following expansion of state  $s_0$ :



Conditional probability is **misleading**:  $s_2$  would be expanded, which isn't part of the **best looking** option:  
with **state-value estimate**  $\hat{V}(s) := h(s)$ , **greedy action**  $a_{\hat{V}}(s) = a_1$

## Expansion in Best Solution Graph

AO\* uses different idea:

- ▶ AO\* keeps track of **best solution graph**
- ▶ AO\* expands a state that can be **reached from  $s_0$**  by only **applying greedy actions**
- ▶  $\Rightarrow$  no  $g$ -value equivalent **required**

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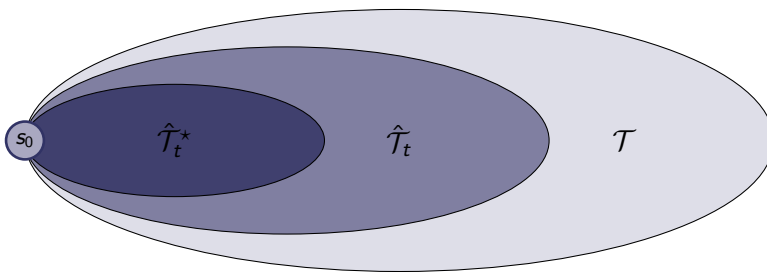
- ▶ AO\* keeps track of **best solution graph**
- ▶ AO\* expands a state that can be **reached from  $s_0$**  by only **applying greedy actions**
- ▶  $\Rightarrow$  no  $g$ -value equivalent **required**
- ▶ Equivalent version of A\* built on this idea can be derived  $\Rightarrow$  **A\* with backward induction**
- ▶ Since change is non-trivial, we focus on A\* variant now
- ▶ and generalize later to acyclic probabilistic tasks (AO\*)
- ▶ and probabilistic tasks in general (LAO\*)

## G1.3 A\* with Backward Induction

## Transition Systems

A\* with backward induction distinguishes **three** transition systems:

- ▶ The transition system  $\mathcal{T} = \langle S, L, c, T, s_0, S^* \rangle$   
 $\Rightarrow$  given **implicitly**
- ▶ The **explicated graph**  $\hat{\mathcal{T}}_t = \langle \hat{S}_t, L, c, \hat{T}_t, s_0, S^* \rangle$   
 $\Rightarrow$  the part of  $\mathcal{T}$  explicitly considered during search
- ▶ The **partial solution graph**  $\hat{\mathcal{T}}_t^* = \langle \hat{S}_t^*, L, c, \hat{T}_t^*, s_0, S^* \rangle$   
 $\Rightarrow$  The part of  $\hat{\mathcal{T}}_t$  that contains best solution



## Explicated Graph

- ▶ **Expanding** a state  $s$  at time step  $t$  **explicates** all successors  $s' \in \text{succ}(s)$  by adding them to **explicated graph**:

$$\hat{\mathcal{T}}_t = \langle \hat{S}_{t-1} \cup \text{succ}(s), L, c, \hat{T}_{t-1} \cup \{ \langle s, l, s' \rangle \in T \}, s_0, S^* \rangle$$

- ▶ Each explicated state is annotated with **state-value estimate**  $\hat{V}_t(s)$  that describes **estimated cost to a goal** at time step  $t$
- ▶ When state  $s'$  is explicated and  $s' \notin \hat{S}_{t-1}$ , its state-value estimate is **initialized** to  $\hat{V}_t(s') := h(s')$
- ▶ We call **leaf states** of  $\hat{\mathcal{T}}_t$  **fringe states**

## Partial Solution Graph

- ▶ The **partial solution graph**  $\hat{\mathcal{T}}_t^*$  is the subgraph of  $\hat{\mathcal{T}}_t$  that is spanned by the **smallest set** of states  $\hat{S}_t^*$  that satisfies:
  - ▶  $s_0 \in \hat{S}_t^*$
  - ▶ if  $s \in \hat{S}_t^*$ ,  $s' \in \hat{S}_t$  and  $\langle s, a_{\hat{V}_t(s)}(s), s' \rangle \in \hat{\mathcal{T}}_t$ , then  $s'$  in  $\hat{S}_t^*$
- ▶ The partial solution graph forms a **sequence of states**  $\langle s_0, \dots, s_n \rangle$ , starting with the initial state  $s_0$  and ending in the **greedy fringe state**  $s_n$

## Backward Induction

- ▶ A\* with backward induction does not maintain **static open list**
- ▶ **State-value estimates** determine **partial solution graph**
- ▶ **Partial solution graph** determines which state is expanded
- ▶ (Some) state-value estimates are **updated** in time step  $t$  by **backward induction**:

$$\hat{V}_t(s) = \min_{\langle s, l, s' \rangle \in \hat{\mathcal{T}}_t(s)} c(l) + \hat{V}_t(s')$$

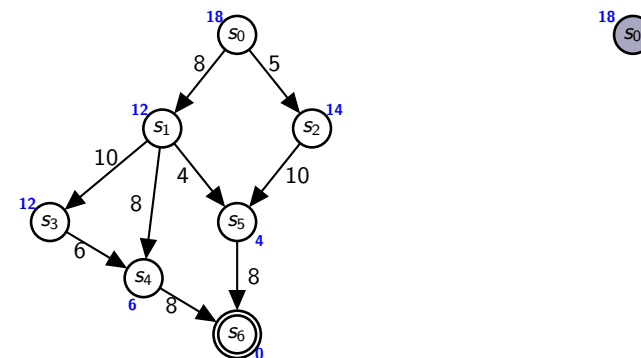
## A\* with backward induction

A\* with backward induction for classical planning task  $\mathcal{T}$

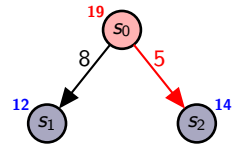
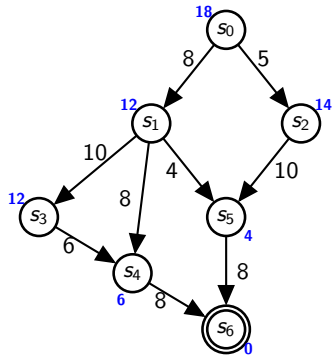
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explicate  $s_0$ 
while greedy fringe state  $s \notin S_*$ :
  expand  $s$ 
  perform backward induction of states in  $\hat{\mathcal{T}}_{t-1}^*$  in reverse order
return  $\hat{\mathcal{T}}_t^*$ 
  
```

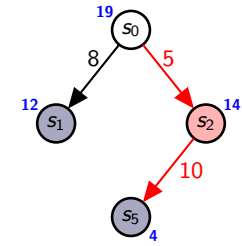
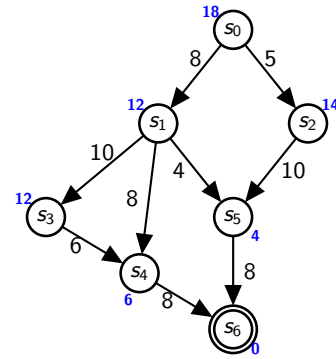
## A\* with backward induction



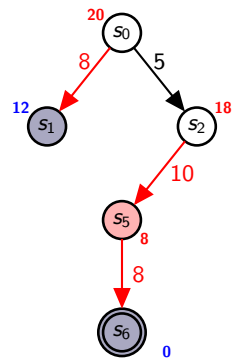
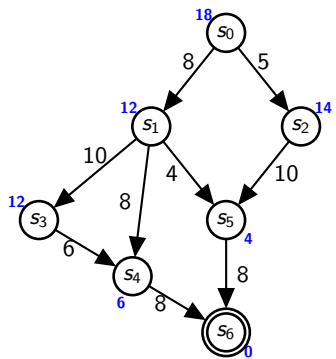
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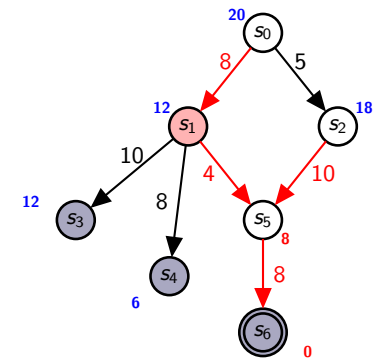
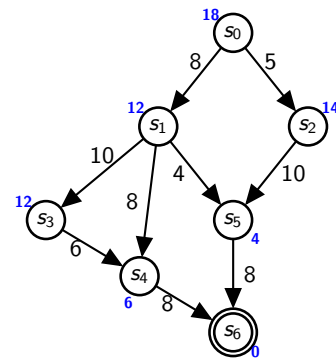
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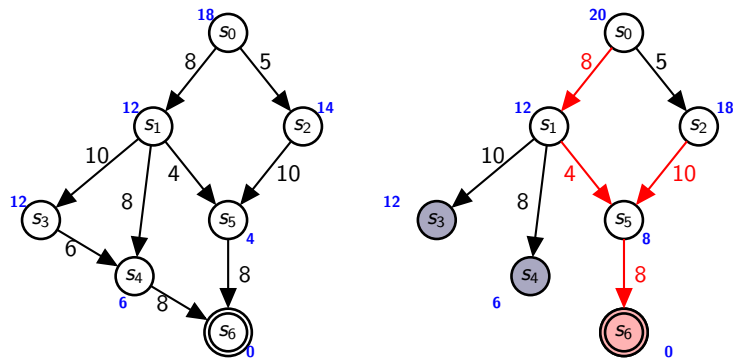
# A\* with backward induction



# A\* with backward induction



## A\* with backward induction



## Equivalence of A\* and A\* with Backward Induction

## Theorem

*A\* and A\* with Backward Induction expand the same set of states if run with identical admissible heuristic  $h$  and identical tie-breaking criterion.*

## Proof Sketch.

The proof shows that

- ▶ there is always a unique state  $s$  in greedy fringe of A\* with backward induction
- ▶  $f(s) = g(s) + h(s)$  is minimal among all fringe states
- ▶  $g(s)$  of fringe node  $s$  encoded in greedy action choices
- ▶  $h(s)$  of fringe node equal to  $\hat{V}_t(s)$

## G1.4 Summary

- ▶ Non-trivial to **generalize** A\* to probabilistic planning
- ▶ For better understanding of AO\*, we **change** A\* towards AO\*
- ▶ Derived **A\* with backward induction**, which is **similar** to AO\*
- ▶ and expands **identical states** as A\*

## Summary

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