

# Planning and Optimization

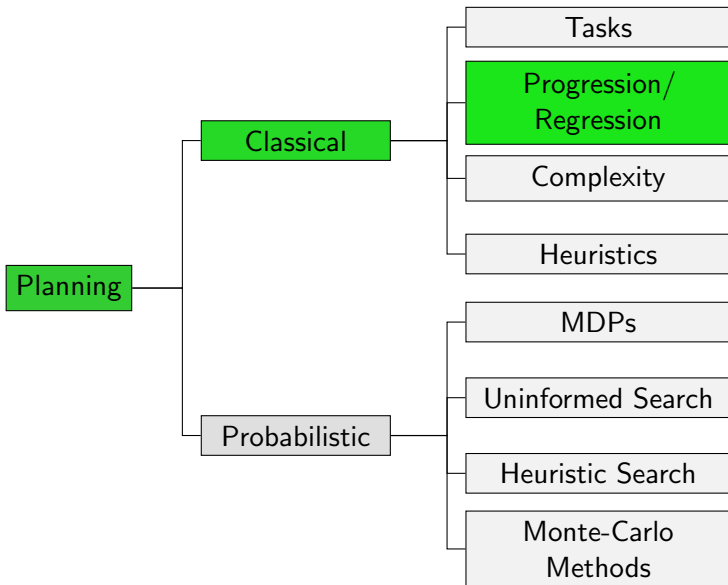
## B3. General Regression

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# Content of this Course



# Regression for General Planning Tasks

- With disjunctions and conditional effects, things become more tricky. How to regress  $a \vee (b \wedge c)$  with respect to  $\langle q, d \triangleright b \rangle$ ?
- In this chapter, we show how to regress **general sets of states** through **general operators**.
- We extensively use the idea of representing sets of states as formulas.

# Regressing State Variables

# Regressing State Variables: Motivation

Key question for general regression:

- Assume we are applying an operator with effect  $e$ .
- What must be true in the predecessor state for state variable  $v$  to be true in the successor state?

If we can answer this question, a general definition of regression is only a small additional step.

## Regressing State Variables: Key Idea

Assume we are in state  $s$  and apply effect  $e$  to obtain successor state  $s'$ .

State variable  $v$  is true in  $s'$  if

- effect  $e$  **makes it true**, or
- it **remains true**, i.e., it is true in  $s$  and not made false by  $e$ .

## Regressing a State Variable Through an Effect

### Definition (Regressing a State Variable Through an Effect)

Let  $e$  be an effect of a propositional planning task, and let  $v$  be a state variable.

The **regression of  $v$  through  $e$** , written  $\text{regr}(v, e)$ , is defined as the following logical formula:

$$\text{regr}(v, e) = \text{effcond}(v, e) \vee (v \wedge \neg \text{effcond}(\neg v, e)).$$

**Question:** Does this capture add-after-delete semantics correctly?

# Regressing State Variables: Example

## Example

Let  $e = (b \triangleright a) \wedge (c \triangleright \neg a) \wedge b \wedge \neg d$ .

$v$	$\text{regr}(v, e)$
$a$	$b \vee (a \wedge \neg c)$
$b$	$\top \vee (b \wedge \neg \perp) \equiv \top$
$c$	$\perp \vee (c \wedge \neg \perp) \equiv c$
$d$	$\perp \vee (d \wedge \neg \top) \equiv \perp$

Reminder:  $\text{regr}(v, e) = \text{effcond}(v, e) \vee (v \wedge \neg \text{effcond}(\neg v, e))$



# Regressing State Variables: Correctness (1)

## Lemma (Correctness of $\text{regr}(v, e)$ )

*In a propositional planning task, let  $s$  be a state,  $e$  be an effect and  $v$  be a state variable.*

*Then  $s \models \text{regr}(v, e)$  iff  $s[e] \models v$ .*

## Regressing State Variables: Correctness (2)

### Proof.

$(\Rightarrow)$ : We know  $s \models \text{regr}(v, e)$ , and hence  $s \models \text{effcond}(v, e) \vee (v \wedge \neg \text{effcond}(\neg v, e))$ .

Do a case analysis on the two disjuncts.

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 $s \models \text{effcond}(v, e) \vee (v \wedge \neg \text{effcond}(\neg v, e))$ .

Do a case analysis on the two disjuncts.

**Case 1:**  $s \models \text{effcond}(v, e)$ .

Then  $s[[e]] \models v$  by the first case in the definition of  $s[[e]]$ .

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### Proof.

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Do a case analysis on the two disjuncts.

**Case 1:**  $s \models \text{effcond}(v, e)$ .

Then  $s[e] \models v$  by the first case in the definition of  $s[e]$ .

**Case 2:**  $s \models (v \wedge \neg \text{effcond}(\neg v, e))$ .

Then  $s \models v$  and  $s \not\models \text{effcond}(\neg v, e)$ .

We may additionally assume  $s \not\models \text{effcond}(v, e)$

because otherwise we can apply Case 1 of this proof.

Then  $s[e] \models v$  by the third case in the definition of  $s[e]$ . ...

## Regressing State Variables: Correctness (3)

Proof (continued).

( $\Leftarrow$ ): Proof by contraposition.

We show that if  $\text{regr}(v, e)$  is **false** in  $s$ , then  $v$  is **false** in  $s[[e]]$ .

## Regressing State Variables: Correctness (3)

Proof (continued).

( $\Leftarrow$ ): Proof by contraposition.

We show that if  $\text{regr}(v, e)$  is **false** in  $s$ , then  $v$  is **false** in  $s[[e]]$ .

- By prerequisite,  $s \not\models \text{effcond}(v, e) \vee (v \wedge \neg \text{effcond}(\neg v, e))$ .

## Regressing State Variables: Correctness (3)

Proof (continued).

( $\Leftarrow$ ): Proof by contraposition.

We show that if  $regr(v, e)$  is **false** in  $s$ , then  $v$  is **false** in  $s[[e]]$ .

- By prerequisite,  $s \not\models effcond(v, e) \vee (v \wedge \neg effcond(\neg v, e))$ .
- Hence  $s \models \neg effcond(v, e) \wedge (\neg v \vee effcond(\neg v, e))$ .

## Regressing State Variables: Correctness (3)

### Proof (continued).

( $\Leftarrow$ ): Proof by contraposition.

We show that if  $\text{regr}(v, e)$  is **false** in  $s$ , then  $v$  is **false** in  $s[[e]]$ .

- By prerequisite,  $s \not\models \text{effcond}(v, e) \vee (v \wedge \neg \text{effcond}(\neg v, e))$ .
- Hence  $s \models \neg \text{effcond}(v, e) \wedge (\neg v \vee \text{effcond}(\neg v, e))$ .
- From the first conjunct, we get  $s \models \neg \text{effcond}(v, e)$  and hence  $s \not\models \text{effcond}(v, e)$ .



## Regressing State Variables: Correctness (3)

### Proof (continued).

( $\Leftarrow$ ): Proof by contraposition.

We show that if  $\text{regr}(v, e)$  is **false** in  $s$ , then  $v$  is **false** in  $s[[e]]$ .

- By prerequisite,  $s \not\models \text{effcond}(v, e) \vee (v \wedge \neg \text{effcond}(\neg v, e))$ .
- Hence  $s \models \neg \text{effcond}(v, e) \wedge (\neg v \vee \text{effcond}(\neg v, e))$ .
- From the first conjunct, we get  $s \models \neg \text{effcond}(v, e)$  and hence  $s \not\models \text{effcond}(v, e)$ .
- From the second conjunct, we get  $s \models \neg v \vee \text{effcond}(\neg v, e)$ .

## Regressing State Variables: Correctness (3)

### Proof (continued).

( $\Leftarrow$ ): Proof by contraposition.

We show that if  $regr(v, e)$  is **false** in  $s$ , then  $v$  is **false** in  $s[[e]]$ .

- By prerequisite,  $s \not\models effcond(v, e) \vee (v \wedge \neg effcond(\neg v, e))$ .
- Hence  $s \models \neg effcond(v, e) \wedge (\neg v \vee effcond(\neg v, e))$ .
- From the first conjunct, we get  $s \models \neg effcond(v, e)$  and hence  $s \not\models effcond(v, e)$ .
- From the second conjunct, we get  $s \models \neg v \vee effcond(\neg v, e)$ .
- **Case 1:**  $s \models \neg v$ . Then  $v$  is false before applying  $e$  and remains false, so  $s[[e]] \not\models v$ .

## Regressing State Variables: Correctness (3)

### Proof (continued).

( $\Leftarrow$ ): Proof by contraposition.

We show that if  $\text{regr}(v, e)$  is **false** in  $s$ , then  $v$  is **false** in  $s[[e]]$ .

- By prerequisite,  $s \not\models \text{effcond}(v, e) \vee (v \wedge \neg \text{effcond}(\neg v, e))$ .
- Hence  $s \models \neg \text{effcond}(v, e) \wedge (\neg v \vee \text{effcond}(\neg v, e))$ .
- From the first conjunct, we get  $s \models \neg \text{effcond}(v, e)$  and hence  $s \not\models \text{effcond}(v, e)$ .
- From the second conjunct, we get  $s \models \neg v \vee \text{effcond}(\neg v, e)$ .
- **Case 1:**  $s \models \neg v$ . Then  $v$  is false before applying  $e$  and remains false, so  $s[[e]] \not\models v$ .
- **Case 2:**  $s \models \text{effcond}(\neg v, e)$ . Then  $v$  is deleted by  $e$  and not simultaneously added, so  $s[[e]] \not\models v$ .



# Regressing Formulas Through Effects

## Regressing Formulas Through Effects: Idea

- We can now generalize regression from state variables to general formulas over state variables.
- The basic idea is to replace **every occurrence** of every state variable  $v$  by  $regr(v, e)$  as defined in the previous section.
- The following definition makes this more formal.

# Regressing Formulas Through Effects: Definition

## Definition (Regressing a Formula Through an Effect)

In a propositional planning task, let  $e$  be an effect, and let  $\varphi$  be a formula over state variables.

The **regression of  $\varphi$  through  $e$** , written  $\text{regr}(\varphi, e)$ , is defined as the following logical formula:

$$\text{regr}(\top, e) = \top$$

$$\text{regr}(\perp, e) = \perp$$

$$\text{regr}(v, e) = \text{effcond}(v, e) \vee (v \wedge \neg \text{effcond}(\neg v, e))$$

$$\text{regr}(\neg\psi, e) = \neg \text{regr}(\psi, e)$$

$$\text{regr}(\psi \vee \chi, e) = \text{regr}(\psi, e) \vee \text{regr}(\chi, e)$$

$$\text{regr}(\psi \wedge \chi, e) = \text{regr}(\psi, e) \wedge \text{regr}(\chi, e).$$

# Regressing Formulas Through Effects: Example

## Example

Let  $e = (b \triangleright a) \wedge (c \triangleright \neg a) \wedge b \wedge \neg d$ .

Recall:

- $\text{regr}(a, e) \equiv b \vee (a \wedge \neg c)$
- $\text{regr}(b, e) \equiv \top$
- $\text{regr}(c, e) \equiv c$
- $\text{regr}(d, e) \equiv \perp$

We get:

$$\begin{aligned}\text{regr}((a \vee d) \wedge (c \vee d), e) &\equiv ((b \vee (a \wedge \neg c)) \vee \perp) \wedge (c \vee \perp) \\ &\equiv (b \vee (a \wedge \neg c)) \wedge c \\ &\equiv b \wedge c\end{aligned}$$

# Regressing Formulas Through Effects: Correctness (1)

## Lemma (Correctness of $\text{regr}(\varphi, e)$ )

Let  $\varphi$  be a logical formula,  $e$  an effect and  $s$  a state.

Then  $s \models \text{regr}(\varphi, e)$  iff  $s[e] \models \varphi$ .



## Regressing Formulas Through Effects: Correctness (2)

Proof.

The proof is by structural induction on  $\varphi$ .

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**Induction hypothesis:**  $s \models \text{regr}(\psi, e)$  iff  $s[[e]] \models \psi$   
for all proper subformulas  $\psi$  of  $\varphi$ .

## Regressing Formulas Through Effects: Correctness (2)

### Proof.

The proof is by structural induction on  $\varphi$ .

**Induction hypothesis:**  $s \models \text{regr}(\psi, e)$  iff  $s[[e]] \models \psi$   
for all proper subformulas  $\psi$  of  $\varphi$ .

**Base case  $\varphi = \top$ :**

We have  $\text{regr}(\top, e) = \top$ , and  $s \models \top$  iff  $s[[e]] \models \top$  is correct.

## Regressing Formulas Through Effects: Correctness (2)

### Proof.

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**Induction hypothesis:**  $s \models \text{regr}(\psi, e)$  iff  $s[[e]] \models \psi$   
for all proper subformulas  $\psi$  of  $\varphi$ .

**Base case  $\varphi = \top$ :**

We have  $\text{regr}(\top, e) = \top$ , and  $s \models \top$  iff  $s[[e]] \models \top$  is correct.

**Base case  $\varphi = \perp$ :**

We have  $\text{regr}(\perp, e) = \perp$ , and  $s \models \perp$  iff  $s[[e]] \models \perp$  is correct.

## Regressing Formulas Through Effects: Correctness (2)

### Proof.

The proof is by structural induction on  $\varphi$ .

**Induction hypothesis:**  $s \models \text{regr}(\psi, e)$  iff  $s[[e]] \models \psi$   
for all proper subformulas  $\psi$  of  $\varphi$ .

**Base case  $\varphi = \top$ :**

We have  $\text{regr}(\top, e) = \top$ , and  $s \models \top$  iff  $s[[e]] \models \top$  is correct.

**Base case  $\varphi = \perp$ :**

We have  $\text{regr}(\perp, e) = \perp$ , and  $s \models \perp$  iff  $s[[e]] \models \perp$  is correct.

**Base case  $\varphi = v$ :**

We have  $s \models \text{regr}(v, e)$  iff  $s[[e]] \models v$  from the previous lemma. ...

## Regressing Formulas Through Effects: Correctness (3)

Proof (continued).

Inductive case  $\varphi = \neg\psi$ :

$$\begin{aligned} s \models \text{regr}(\neg\psi, e) &\text{ iff } s \models \neg\text{regr}(\psi, e) \\ &\text{ iff } s \not\models \text{regr}(\psi, e) \\ &\text{ iff } s[[e]] \not\models \psi \\ &\text{ iff } s[[e]] \models \neg\psi \end{aligned}$$

# Regressing Formulas Through Effects: Correctness (3)

Proof (continued).

Inductive case  $\varphi = \neg\psi$ :

$$\begin{aligned} s \models \text{regr}(\neg\psi, e) &\text{ iff } s \models \neg\text{regr}(\psi, e) \\ &\text{ iff } s \not\models \text{regr}(\psi, e) \\ &\text{ iff } s[[e]] \not\models \psi \\ &\text{ iff } s[[e]] \models \neg\psi \end{aligned}$$

Inductive case  $\varphi = \psi \vee \chi$ :

$$\begin{aligned} s \models \text{regr}(\psi \vee \chi, e) &\text{ iff } s \models \text{regr}(\psi, e) \vee \text{regr}(\chi, e) \\ &\text{ iff } s \models \text{regr}(\psi, e) \text{ or } s \models \text{regr}(\chi, e) \\ &\text{ iff } s[[e]] \models \psi \text{ or } s[[e]] \models \chi \\ &\text{ iff } s[[e]] \models \psi \vee \chi \end{aligned}$$

# Regressing Formulas Through Effects: Correctness (3)

## Proof (continued).

Inductive case  $\varphi = \neg\psi$ :

$$\begin{aligned} s \models \text{regr}(\neg\psi, e) &\text{ iff } s \models \neg\text{regr}(\psi, e) \\ &\text{ iff } s \not\models \text{regr}(\psi, e) \\ &\text{ iff } s[[e]] \not\models \psi \\ &\text{ iff } s[[e]] \models \neg\psi \end{aligned}$$

Inductive case  $\varphi = \psi \vee \chi$ :

$$\begin{aligned} s \models \text{regr}(\psi \vee \chi, e) &\text{ iff } s \models \text{regr}(\psi, e) \vee \text{regr}(\chi, e) \\ &\text{ iff } s \models \text{regr}(\psi, e) \text{ or } s \models \text{regr}(\chi, e) \\ &\text{ iff } s[[e]] \models \psi \text{ or } s[[e]] \models \chi \\ &\text{ iff } s[[e]] \models \psi \vee \chi \end{aligned}$$

Inductive case  $\varphi = \psi \wedge \chi$ :

Like previous case, replacing “ $\vee$ ” by “ $\wedge$ ”  
and replacing “or” by “and”.





# Regressing Formulas Through Operators

# Regressing Formulas Through Operators: Idea

- We can now regress arbitrary formulas through arbitrary effects.
- The last missing piece is a definition of regression through **operators**, describing exactly in which states  $s$  applying a given operator  $o$  leads to a state satisfying a given formula  $\varphi$ .
- There are two requirements:
  - The operator  $o$  must be **applicable** in the state  $s$ .
  - The **resulting state**  $s[[o]]$  must **satisfy**  $\varphi$ .

# Regressing Formulas Through Operators: Definition

## Definition (Regressing a Formula Through an Operator)

In a propositional planning task, let  $o$  be an operator, and let  $\varphi$  be a formula over state variables.

The **regression of  $\varphi$  through  $o$** , written  $\text{regr}(\varphi, o)$ , is defined as the following logical formula:

$$\text{regr}(\varphi, o) = \text{pre}(o) \wedge \text{regr}(\varphi, \text{eff}(o)).$$

# Regressing Formulas Through Operators: Correctness (1)

## Theorem (Correctness of $\text{regr}(\varphi, o)$ )

*Let  $\varphi$  be a logical formula,  $o$  an operator and  $s$  a state.*

*Then  $s \models \text{regr}(\varphi, o)$  iff  $o$  is applicable in  $s$  and  $s[o] \models \varphi$ .*

## Regressing Formulas Through Operators: Correctness (2)

Reminder:  $\text{regr}(\varphi, o) = \text{pre}(o) \wedge \text{regr}(\varphi, \text{eff}(o))$

Proof.

Case 1:  $s \models \text{pre}(o)$ .

Then  $o$  is applicable in  $s$  and the statement we must prove simplifies to:  $s \models \text{regr}(\varphi, \text{eff}(o))$  iff  $s[[e]] \models \varphi$ .

This was proved in the previous lemma.

## Regressing Formulas Through Operators: Correctness (2)

Reminder:  $\text{regr}(\varphi, o) = \text{pre}(o) \wedge \text{regr}(\varphi, \text{eff}(o))$

Proof.

Case 1:  $s \models \text{pre}(o)$ .

Then  $o$  is applicable in  $s$  and the statement we must prove simplifies to:  $s \models \text{regr}(\varphi, \text{eff}(o))$  iff  $s[[e]] \models \varphi$ .

This was proved in the previous lemma.

Case 2:  $s \not\models \text{pre}(o)$ .

Then  $s \not\models \text{regr}(\varphi, o)$  and  $o$  is not applicable in  $s$ .

Hence both statements are false and therefore equivalent. □

# Regression Examples (1)

Examples: compute regression and simplify to DNF

- $\text{regr}(b, \langle a, b \rangle)$   
 $\equiv a \wedge (\top \vee (b \wedge \neg \perp))$   
 $\equiv a$
- $\text{regr}(b \wedge c \wedge d, \langle a, b \rangle)$   
 $\equiv a \wedge (\top \vee (b \wedge \neg \perp)) \wedge (\perp \vee (c \wedge \neg \perp)) \wedge (\perp \vee (d \wedge \neg \perp))$   
 $\equiv a \wedge c \wedge d$
- $\text{regr}(b \wedge \neg c, \langle a, b \wedge c \rangle)$   
 $\equiv a \wedge (\top \vee (b \wedge \neg \perp)) \wedge \neg(\top \vee (c \wedge \neg \perp))$   
 $\equiv a \wedge \top \wedge \perp$   
 $\equiv \perp$

## Regression Examples (2)

Examples: compute regression and simplify to DNF

- $\text{regr}(b, \langle a, c \triangleright b \rangle)$ 
  - $\equiv a \wedge (c \vee (b \wedge \neg \perp))$
  - $\equiv a \wedge (c \vee b)$
  - $\equiv (a \wedge c) \vee (a \wedge b)$
- $\text{regr}(b, \langle a, (c \triangleright b) \wedge ((d \wedge \neg c) \triangleright \neg b) \rangle)$ 
  - $\equiv a \wedge (c \vee (b \wedge \neg(d \wedge \neg c)))$
  - $\equiv a \wedge (c \vee (b \wedge (\neg d \vee c)))$
  - $\equiv a \wedge (c \vee (b \wedge \neg d) \vee (b \wedge c))$
  - $\equiv a \wedge (c \vee (b \wedge \neg d))$
  - $\equiv (a \wedge c) \vee (a \wedge b \wedge \neg d)$



# Summary

# Summary

- Regressing a **state variable** through an (arbitrary) operator must consider two cases:
  - state variables **made true** (by add effects)
  - state variables **remaining true** (by absence of delete effects)
- Regression of state variables can be generalized to arbitrary formulas  $\varphi$  by replacing each occurrence of a state variable in  $\varphi$  by its regression.
- **Regressing a formula  $\varphi$**  through an **operator** involves regressing  $\varphi$  through the effect and enforcing the precondition.