

# Theory of Computer Science

## A3. Proof Techniques

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A3. Proof Techniques

Proofs

## A3.1 Proofs

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A3. Proof Techniques

Proofs

## What is a Proof?

A **mathematical proof** is

- ▶ a sequence of logical steps
- ▶ starting with one set of statements
- ▶ that comes to the conclusion  
that some statement must be true.

What is a **statement**?

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## Mathematical Statements

### Mathematical Statement

A **mathematical statement** is a declarative sentence that is either true or false (but not both).

Examples (some true, some false):

- ▶ Let  $p \in \mathbb{N}_0$  be a prime number. Then  $p$  is odd.
- ▶ There exists an even prime number.
- ▶ The equation  $a^k + b^k = c^k$  has infinitely many solutions with  $a, b, c, k \in \mathbb{N}_1$  and  $k \geq 2$ .

German: Mathematische Aussage

## Mathematical Statements: Quantification

Statements often use **quantification**.

- ▶ **Universal quantification:**  
 “For all  $x$  in set  $S$  it holds that  $\langle$ sub-statement on  $x$  $\rangle$ .”  
 This is **true** if the sub-statement is true for every  $x$  in  $S$ .
- ▶ **Existential quantification:**  
 “There is an  $x$  in set  $S$  such that  $\langle$ sub-statement on  $x$  $\rangle$ .”  
 This is **true** if there exists at least one  $x$  in  $S$  for which the sub-statement is true.

Examples (some true, some false):

- ▶ For all  $x \in \mathbb{N}_1$  it holds that  $x + 1$  is in  $\mathbb{N}_1$ .
- ▶ For all  $x \in \mathbb{N}_1$  it holds that  $x - 1$  is in  $\mathbb{N}_1$ .
- ▶ There is an  $x \in \mathbb{N}_1$  such that  $x = \sqrt{x}$ .

## Mathematical Statements: Preconditions and Conclusions

We can identify **preconditions** and **conclusions**.

“If  $\langle$ preconditions $\rangle$  then  $\langle$ conclusions $\rangle$ .”

The statement is **true** if the conclusions are true whenever the preconditions are true.

Not every statement has preconditions. Preconditions are often used in universally quantified sub-statements.

Examples (some true, some false):

- ▶ If 4 is a prime number then  $2 \cdot 3 = 4$ .
- ▶ If  $n$  is a prime number with  $n > 2$  then  $n$  is odd.
- ▶ For all  $p \in \mathbb{N}_1$  it holds that if  $p$  is a prime number then  $p$  is odd.

## Different Statements with the same Meaning

The following statements have the same meaning, we just move preconditions into the quantification, make some aspects implicit, and change the structure.

- ▶ For all  $p \in \mathbb{N}_1$  it holds that if  $p$  is a prime number with  $p > 2$  then  $p$  is odd.
- ▶ For all prime numbers  $p$  it holds that if  $p > 2$  then  $p$  is odd.
- ▶ Let  $p$  be a natural number with  $p > 2$ . Then  $p$  is odd if  $p$  is prime.
- ▶ If  $p$  is a prime number with  $p > 2$  then  $p$  is odd.
- ▶ All prime numbers  $p > 2$  are odd.

A single mathematical statement can be expressed in different ways, as long as the meaning stays the same.

Like paraphrasing a sentence in everyday language.

## On what Statements can we Build the Proof?

A mathematical proof is

- ▶ a sequence of logical steps
- ▶ **starting with one set of statements**
- ▶ that comes to the conclusion that some statement must be true.

We can use:

- ▶ **axioms**: statements that are assumed to always be true in the current context
- ▶ **theorems** and **lemmas**: statements that were already proven
  - ▶ lemma: an intermediate tool
  - ▶ theorem: itself a relevant result
- ▶ **premises**: assumptions we make to see what consequences they have

German: Axiom, Theorem/Satz, Lemma, Prämisse/Annahme

## What is a Logical Step?

A mathematical proof is

- ▶ **a sequence of logical steps**
- ▶ starting with one set of statements
- ▶ that comes to the conclusion that some statement must be true.

Each step **directly follows**

- ▶ from the axioms,
- ▶ premises,
- ▶ previously proven statements and
- ▶ the preconditions of the statement we want to prove.

For a formal definition, we would need formal logics.

## The Role of Definitions

### Definition

A **set** is an unordered collection of distinct objects.

The objects in a set are called the **elements** of the set. A set is said to **contain** its elements.

We write  $x \in S$  to indicate that  $x$  is an element of set  $S$ , and  $x \notin S$  to indicate that  $S$  does not contain  $x$ .

The set that does not contain any objects is the **empty set**  $\emptyset$ .

- ▶ A definition introduces an abbreviation.
- ▶ Whenever we say “set”, we could instead say “an unordered collection of distinct objects” and vice versa.
- ▶ Definitions can also introduce notation.

German: Definition

## Disproofs

- ▶ A **disproof** (**refutation**) shows that a given mathematical statement is **false** by giving an example where the preconditions are true, but the conclusion is false.
- ▶ This requires deriving, in a sequence of proof steps, the opposite (negation) of the conclusion.

### Example (False statement)

“If  $p \in \mathbb{N}_0$  is a prime number then  $p$  is odd.”

### Refutation.

Consider natural number 2 as a counter example. It is prime because it has exactly 2 divisors, 1 and itself. It is not odd, because it is divisible by 2. □

German: Widerlegung

## Exercise

You want to disprove the following statement with a counterexample:

If the sun is shining then all kids eat ice cream.

What properties must your counterexample have?

[Discuss with your neighbour; 2 minutes]



## A Word on Style

A proof should help the reader to see why the result must be true.

- ▶ A proof should be easy to follow.
- ▶ Omit unnecessary information.
- ▶ Move self-contained parts into separate lemmas.
- ▶ In complicated proofs, reveal the overall structure in advance.
- ▶ Have a clear line of argument.

→ Writing a proof is like writing an essay.

## A3.2 Proof Strategies

## Proof Strategies

typical proof/disproof strategies:

- 1 “All  $x \in S$  with the property  $P$  also have the property  $Q$ .”  
 “For all  $x \in S$ : if  $x$  has property  $P$ , then  $x$  has property  $Q$ .”
  - ▶ To prove, assume you are given an arbitrary  $x \in S$  that has the property  $P$ .  
Give a sequence of proof steps showing that  $x$  must have the property  $Q$ .
  - ▶ To disprove, find a **counterexample**, i. e., find an  $x \in S$  that has property  $P$  but not  $Q$  and prove this.

## Proof Strategies

typical proof/disproof strategies:

- ② “ $A$  is a subset of  $B$ .”
  - ▶ To prove, assume you have an arbitrary element  $x \in A$  and prove that  $x \in B$ .
  - ▶ To disprove, find an element in  $x \in A \setminus B$  and prove that  $x \in A \setminus B$ .

## Proof Strategies

typical proof/disproof strategies:

- ③ “For all  $x \in S$ :  $x$  has property  $P$  iff  $x$  has property  $Q$ .”  
(“iff”: “if and only if”)
  - ▶ To prove, separately prove “if  $P$  then  $Q$ ” and “if  $Q$  then  $P$ ”.
  - ▶ To disprove, disprove “if  $P$  then  $Q$ ” or disprove “if  $Q$  then  $P$ ”.

## Proof Strategies

typical proof/disproof strategies:

- ④ “ $A = B$ ”, where  $A$  and  $B$  are sets.
  - ▶ To prove, separately prove “ $A \subseteq B$ ” and “ $B \subseteq A$ ”.
  - ▶ To disprove, disprove “ $A \subseteq B$ ” or disprove “ $B \subseteq A$ ”.

## Proof Techniques

proof techniques we use in this course:

- ▶ direct proof
- ▶ indirect proof (proof by contradiction)

## A3.3 Direct Proof

## Direct Proof

### Direct Proof

Direct derivation of the statement by deducing or rewriting.

German: Direkter Beweis

## Direct Proof: Example

### Theorem (distributivity)

For all sets  $A, B, C$ :  $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$ .

### Proof.

We first show that  $x \in A \cap (B \cup C)$  implies  $x \in (A \cap B) \cup (A \cap C)$  ( $\subseteq$  part):

Let  $x \in A \cap (B \cup C)$ . Then by the definition of  $\cap$  it holds that  $x \in A$  and  $x \in B \cup C$ .

We make a case distinction between  $x \in B$  and  $x \notin B$ :

If  $x \in B$  then, because  $x \in A$  is true,  $x \in A \cap B$  must be true.

Otherwise, because  $x \in B \cup C$  we know that  $x \in C$  and thus with  $x \in A$ , that  $x \in A \cap C$ .

In both cases  $x \in A \cap B$  or  $x \in A \cap C$ , and we conclude  $x \in (A \cap B) \cup (A \cap C)$ . ...

## Direct Proof: Example

### Theorem (distributivity)

For all sets  $A, B, C$ :  $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$ .

### Proof (continued).

$\supseteq$  part: we must show that  $x \in (A \cap B) \cup (A \cap C)$  implies  $x \in A \cap (B \cup C)$ .

Let  $x \in (A \cap B) \cup (A \cap C)$ .

We make a case distinction between  $x \in A \cap B$  and  $x \notin A \cap B$ :

If  $x \in A \cap B$  then  $x \in A$  and  $x \in B$ .

The latter implies  $x \in B \cup C$  and hence  $x \in A \cap (B \cup C)$ .

If  $x \notin A \cap B$  we know  $x \in A \cap C$  due to  $x \in (A \cap B) \cup (A \cap C)$ .

This (analogously) implies  $x \in A$  and  $x \in C$ , and hence  $x \in B \cup C$  and thus  $x \in A \cap (B \cup C)$ .

In both cases we conclude  $x \in A \cap (B \cup C)$ . ...

## Direct Proof: Example

### Theorem (distributivity)

For all sets  $A, B, C$ :  $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$ .

### Proof (continued).

We have shown that every element of  $A \cap (B \cup C)$  is an element of  $(A \cap B) \cup (A \cap C)$  and vice versa. Thus, both sets are equal.  $\square$

## Direct Proof: Example

### Theorem (distributivity)

For all sets  $A, B, C$ :  $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$ .

### Proof.

Alternative:

$$\begin{aligned} A \cap (B \cup C) &= \{x \mid x \in A \text{ and } x \in B \cup C\} \\ &= \{x \mid x \in A \text{ and } (x \in B \text{ or } x \in C)\} \\ &= \{x \mid (x \in A \text{ and } x \in B) \text{ or } (x \in A \text{ and } x \in C)\} \\ &= \{x \mid x \in A \cap B \text{ or } x \in A \cap C\} \\ &= (A \cap B) \cup (A \cap C) \end{aligned}$$

 $\square$ 

## A3.4 Indirect Proof

## Indirect Proof

### Indirect Proof (Proof by Contradiction)

- ▶ Make an **assumption** that the statement is false.
- ▶ Use the assumption to derive a **contradiction**.
- ▶ This shows that the assumption must be false and hence the original statement must be true.

German: Indirekter Beweis, Beweis durch Widerspruch

## Indirect Proof: Example 1

### Theorem

Let  $A$  and  $B$  be sets. If  $A \setminus B = \emptyset$  then  $A \subseteq B$ .

### Proof.

We prove the theorem by contradiction.

Assume that there are sets  $A$  and  $B$  with  $A \setminus B = \emptyset$  and  $A \not\subseteq B$ .

Let  $A$  and  $B$  be such sets.

Since  $A \not\subseteq B$  there is some  $x \in A$  such that  $x \notin B$ .

For this  $x$  it holds that  $x \in A \setminus B$ .

This is a contradiction to  $A \setminus B = \emptyset$ .

We conclude that the assumption was false and thus the theorem is true.  $\square$

## Indirect Proof: Example 2

### Theorem

There are infinitely many prime numbers.

### Proof.

**Assumption:** There are only finitely many prime numbers.

Let  $P = \{p_1, \dots, p_n\}$  be the set of all prime numbers.

Define  $m = p_1 \cdot \dots \cdot p_n + 1$ .

Since  $m \geq 2$ , it must have a prime factor.

Let  $p$  be such a prime factor.

Since  $p$  is a prime number,  $p$  has to be in  $P$ .

The number  $m$  is not divisible without remainder by any of the numbers in  $P$ . Hence  $p$  is no factor of  $m$ .

$\rightsquigarrow$  **Contradiction**  $\square$

## A3.5 Summary

## Summary

- ▶ A **proof** is based on axioms and previously proven statements.
- ▶ Individual **proof steps** must be obvious derivations.
- ▶ **direct proof**: sequence of derivations or rewriting
- ▶ **indirect proof**: refute that the statement is false.