

Knowledge, Reasoning and Planning

1. Organization, Dates & Topics

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Capacity Limit

Capacity: ≤ 20 students in this seminar

Priority rules:

- 1st priority: students who have already successfully completed the course “Planning and Optimization”
- 2nd priority: students who are enrolled in the MSc computer science programme with major in machine intelligence
- provided that they have registered for the seminar until February 20, 2026
- otherwise first-come first-served

Format

Seminar Format

- theoretical part + project
- 6 ECTS points
- assessment: graded (1.0–6.0)

Grading

Grading

- Presentation: 25%
 - Written Report: 25%
 - Peer Review: 15%
 - Project: 25%
 - Project Presentation: 10%
-
- each part is individually graded (1.0–6.0)

Plagiarism

Plagiarism

- **plagiarism:** presenting work or ideas of others as your own
 - this includes AI-generated content!
- consequence: failing the seminar
- if in doubt: **ask us in advance!**

Repeat offenders can be excluded from the study program.

Course Links

Seminar Homepage

<https://dmi.unibas.ch/en/studies/computer-science/course-offer-spring-2026/77779-seminar-knowledge-reasoning-and-planning/>

- seminar description
- slides
- examples for good talk slides, reports, and reviews
- link to ADAM workspace

ADAM Workspace

- additional material that we cannot share publicly
- link to Discord server

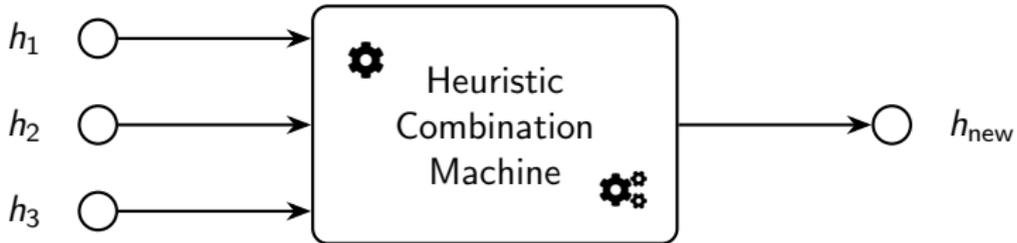
Discord Server

- place to interact with us and with each other

Questions about the Organization

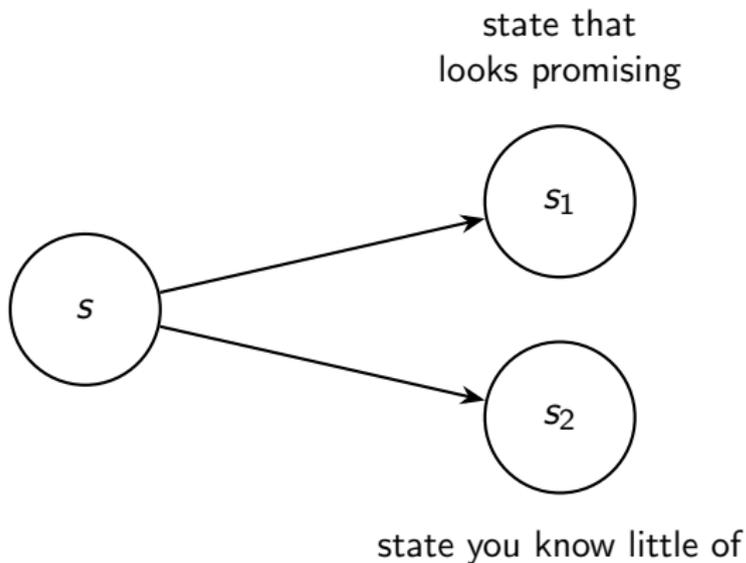
Questions?

3. Satisficing Search with Multiple Heuristics



- different heuristics good for different search space regions
- Scorpion Maidu, winner of the last international planning competition alternates between six open lists

4. Width- and Novelty-Based Search



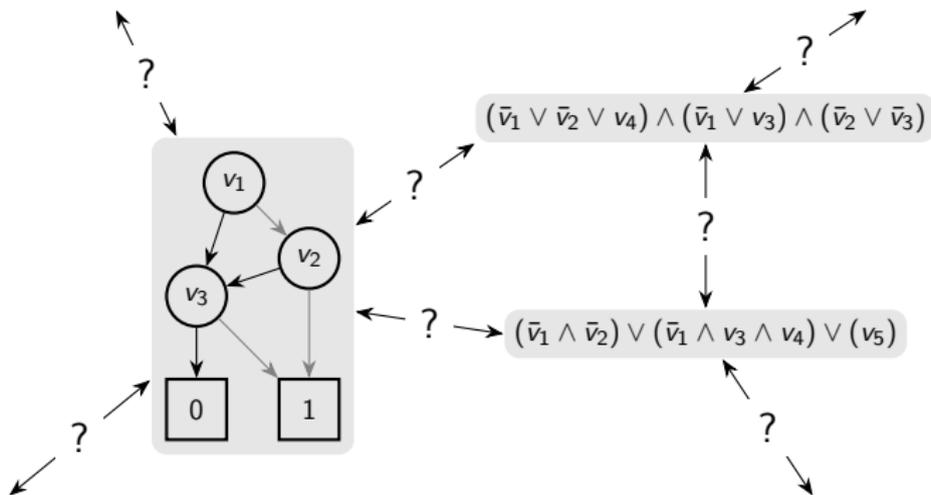
- exploration vs. exploitation tradeoff
- “novelty” in planning

5. Planning as Symbolic Search

```
while goal-states  $\cap$  visited-states  $\neq \emptyset$ :  
    if no applicable operators:  
        return unsolvable  
    visited-states.apply(operators)  
return plan
```

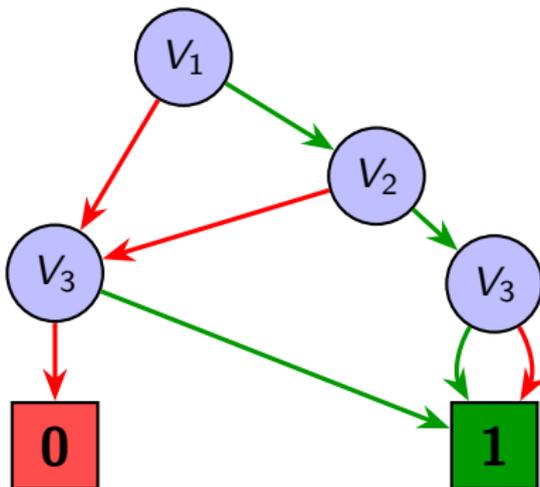
- keep track of all visited states using a compact representation
- simulate a standard search algorithm like BFS

6. Knowledge Compilation



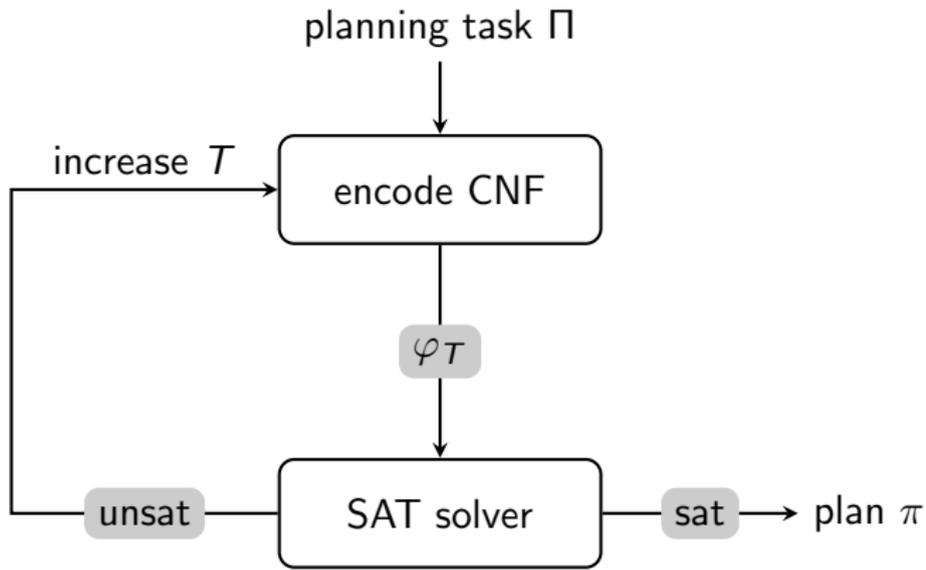
- How can knowledge be represented succinctly?
- Which queries can be answered efficiently?
- Which operations can be performed efficiently?

7. Decision Diagrams



- compact representation of logical formulas
- BDDs, ZDDs, relaxed decision diagrams, restricted decision diagrams

8. Planning as SAT



- encode plan existence within horizon T as Boolean satisfiability problem
- different encodings: sequential, \forall -step, \exists -step, ...

9. Planning with Property Directed Reachability

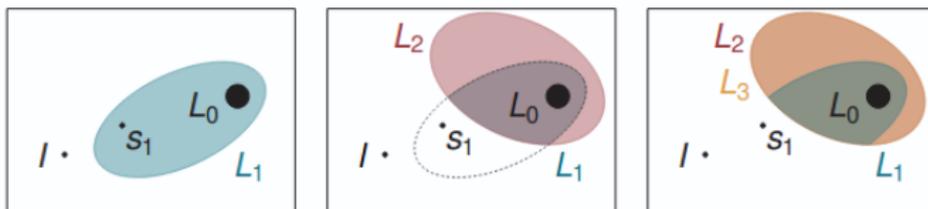


Figure: Eriksson and Helmert (2020)

- Use propositional formulas to overapproximate set of states within k steps from the goal.
- Refine overapproximation using counterexamples.

10. Tractable Boolean Satisfaction Problems

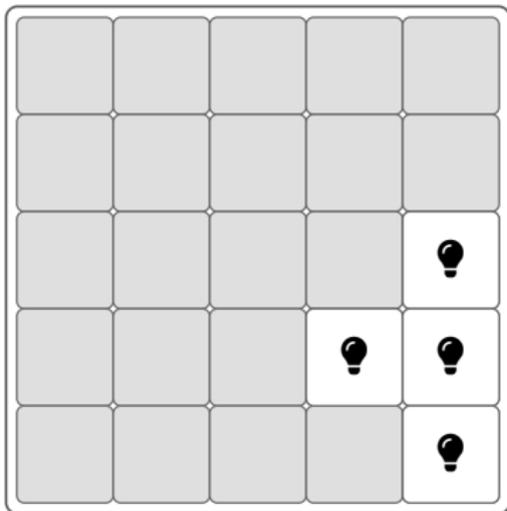


Figure: Lights Out puzzle

- some boolean constraint satisfaction problems are tractable
- Which ones? Why?
- relationship to SAT

11. SAT Solving

Is the following formula satisfiable?

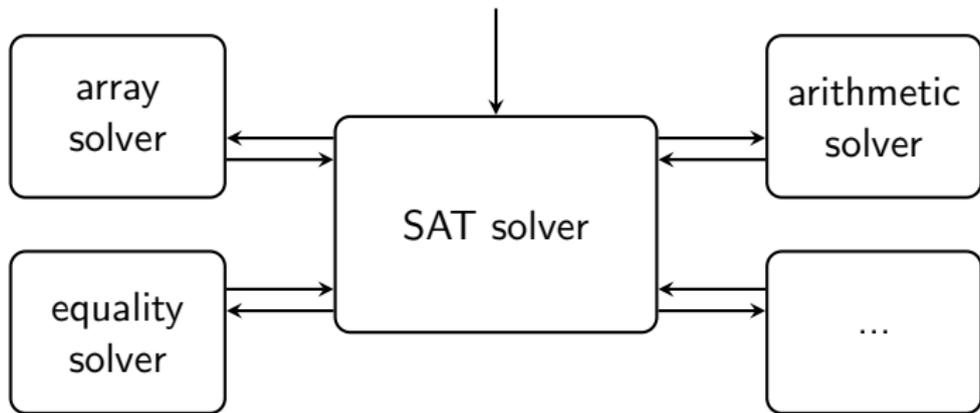
$$(p \vee q \vee \neg r) \wedge (\neg p \vee r) \wedge (\neg q) \wedge (p \vee r)$$

- SAT is NP-complete
- SAT solvers work well in practice
- basic algorithm DPLL/CDCL

12. SMT Solving

Is the following formula satisfiable **modulo theory T** ?

$$(a = 0) \wedge (b = 0 \vee b = c) \wedge (c = a - 2) \wedge (f(a) \neq f(b))$$



- basic algorithm DPLL(T)/CDCL(T)
- decision procedures

13. Plan Improvement

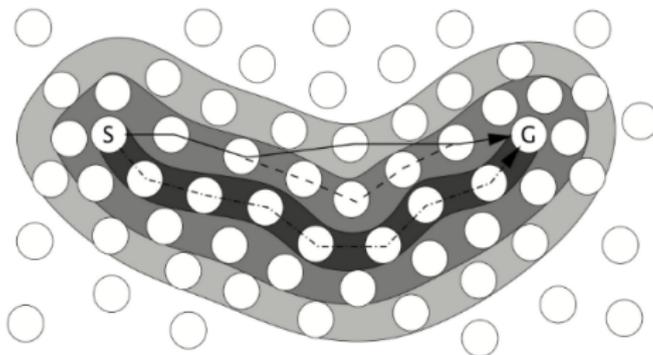
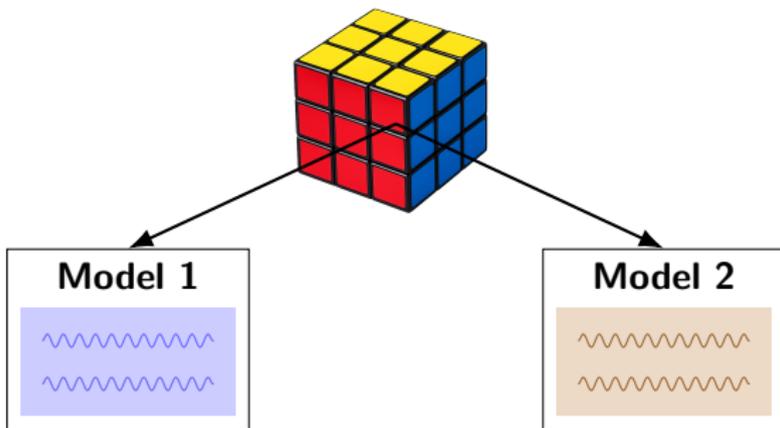


Figure: Basil Kohler (2012)

- start with a suboptimal plan
- apply an iterative algorithm that improves its quality

14. PDDL Modeling



- impact on performance and heuristic values
- compactness of representation
- axioms, conditional effects
- pros and cons of redundancy

15. Grounding PDDL Models

$$\text{PDDL} \rightarrow \langle \mathcal{V}, \mathcal{O}, s_0, \gamma \rangle$$

- planning problems are often defined using the first-order language PDDL
- most planners translate (i.e., ground) PDDL problems into the type of planning tasks you are familiar with
- How does this work?

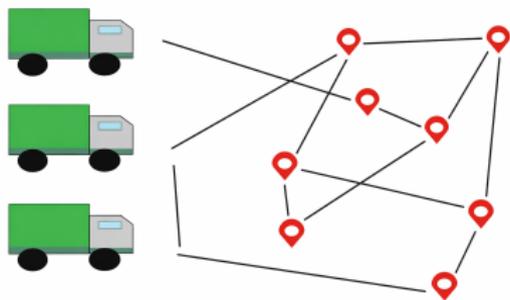
16. Logic Programs and ASP

```
topic(programming; logic; ai).  
  
prereq(ai, programming).  
prereq(ai, logic).  
  
tired(logic).  
  
skip(T) :- tired(T).  
skip(T) :- prereq(T,P), not study(P).  
  
study(T) :- has-book(T), not skip(T).
```

- model complex domains with rules and constraints
- reason about recursive conditions and transitive closure

17. Invariant Synthesis for Classical Planning

invariant synthesis: automatically find things that are true in every reachable state



In every reachable state:

- no two trucks are in the same position
- all trucks are in exactly one position
- there are exactly three trucks
- ...

18. Proof Calculi

***54·43.** $\vdash \therefore \alpha, \beta \in 1. \supset : \alpha \wedge \beta = \Lambda. \equiv . \alpha \vee \beta \in 2$

Dem.

$\vdash . *54·26. \supset \vdash \therefore \alpha = \iota'x. \beta = \iota'y. \supset : \alpha \vee \beta \in 2. \equiv . x \neq y.$
 $[*51·231] \qquad \qquad \qquad \equiv . \iota'x \wedge \iota'y = \Lambda .$
 $[*13·12] \qquad \qquad \qquad \equiv . \alpha \wedge \beta = \Lambda \qquad (1)$

$\vdash . (1). *11·11·35. \supset$
 $\vdash \therefore (\exists x, y). \alpha = \iota'x. \beta = \iota'y. \supset : \alpha \vee \beta \in 2. \equiv . \alpha \wedge \beta = \Lambda \qquad (2)$

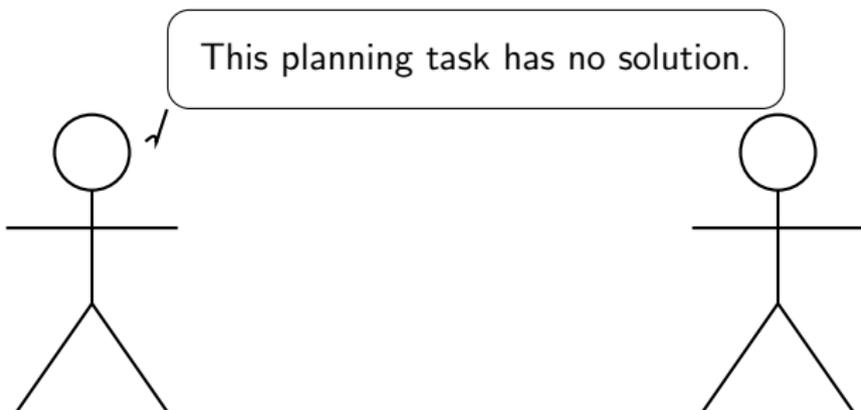
$\vdash . (2). *11·54. *52·1. \supset \vdash . \text{Prop}$

From this proposition it will follow, when arithmetical addition has been defined, that $1 + 1 = 2$.

Figure: Principia Mathematica, p. 379

- What is a proof?
- Are there algorithms that can automatically prove theorems for us?

19. Unsolvability Certificates for Planning

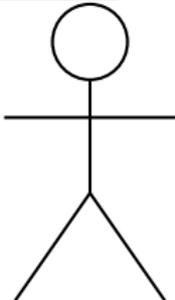
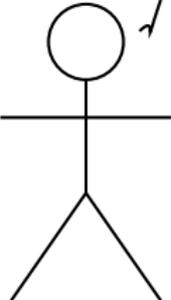


19. Unsolvability Certificates for Planning

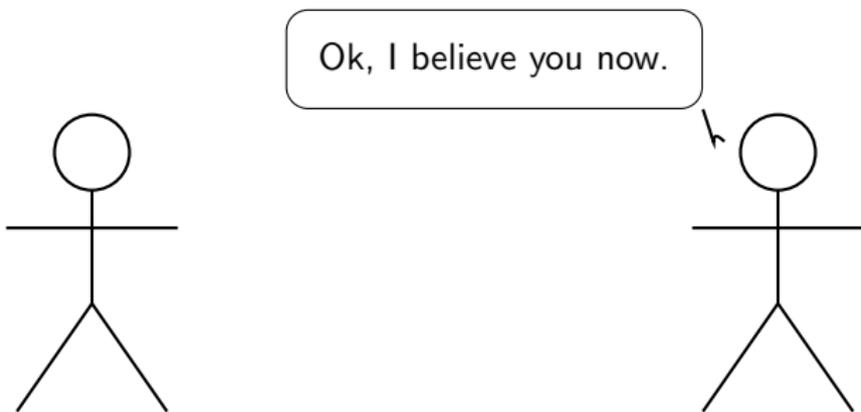


19. Unsolvability Certificates for Planning

1. $I_s \models \varphi$ (MO)
2. $\varphi \models \bigvee_{v \in G} \neg v$ (CE)
3. $\varphi' := \varphi[V \rightarrow V']$ (RN[↖])
4. $\beta_o := \bigwedge_{v \in \text{add}(o)} v'$ (CL)
5. $\eta_o := \bigwedge_{v \in V \setminus (\text{add}(o) \cup \text{del}(o))} (v \leftrightarrow v')$ (BI[↖])
6. $\chi_o \models \varphi'$ (SE)
- ⋮



19. Unsolvability Certificates for Planning



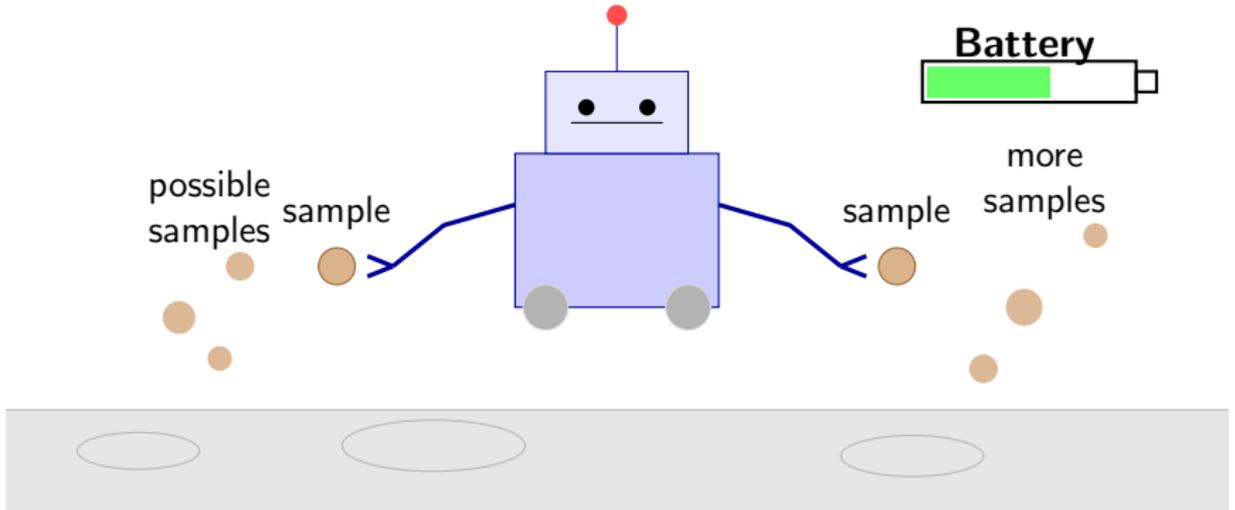
20. Unsolvability with Parity Constraint Invariants and Monovariants



Figure: The Fifteen Puzzle

- algorithms that automatically detect unsolvability using parity and monotonicity arguments

21. Oversubscription Planning



- e.g., maximize the amount of samples collected without running out of battery
- oversubscription planning = find a plan that maximizes utility given a budget

22. FOND Planning

Classical Planning

start-engine

pre: engine-off \wedge has-key
eff: engine-on

FOND Planning

start-engine

pre: engine-off \wedge has-key
eff: engine-on **or** key-jammed

- fully **observable**, **non-deterministic** planning
- non-deterministic operator effects

23. Numeric Planning with MIP Compilation

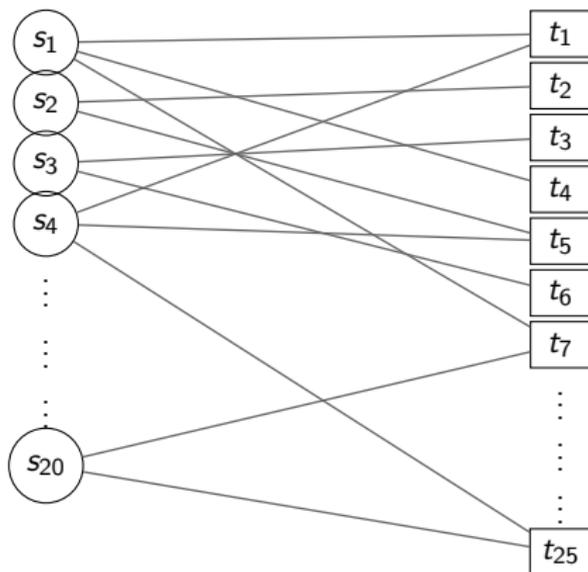
drive-Basel-Freiburg

pre: $\text{in-Basel} \wedge \text{fuel} \geq 30 \wedge \text{has-passport}$

eff: $\neg \text{in-Basel} \wedge \text{in-Freiburg} \wedge \text{fuel} := \text{fuel} - 30$

- planning with numeric variables
- compile the problem into a mixed integer program (MIP)
- solve it with a MIP solver

25. LP Solvers for Bipartite Graph Matching



- 20 students, ≥ 3 topics each: can you find a matching such that no two students get the same topic?
- IP and LP bipartite graph matching
- Complexity and comparison to standard algorithms