# Foundations of Artificial Intelligence F4. Automated Planning: Delete Relaxation Heuristics

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# Foundations of Artificial Intelligence

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### Automated Planning: Overview

#### Chapter overview: automated planning

- ► F1. Introduction
- ► F2. Planning Formalisms
- ► F3. Delete Relaxation
- ► F4. Delete Relaxation Heuristics
- ► F5. Abstraction
- ► F6. Abstraction Heuristics

# F4.1 Relaxed Planning Graphs

# Relaxed Planning Graphs

- relaxed planning graphs: represent which variables in Π<sup>+</sup> can be reached and how
- ightharpoonup graphs with variable layers  $V^i$  and action layers  $A^i$ 
  - ightharpoonup variable layer  $V^0$  contains the variable vertex  $v^0$  for all  $v \in I$
  - ▶ action layer  $A^{i+1}$  contains the action vertex  $a^{i+1}$  for action a if  $V^i$  contains the vertex  $v^i$  for all  $v \in pre(a)$
  - variable layer  $V^{i+1}$  contains the variable vertex  $v^{i+1}$  if previous variable layer contains  $v^i$ , or previous action layer contains  $a^{i+1}$  with  $v \in add(a)$

German: relaxierter Planungsgraph, Variablenknoten, Aktionsknoten

# Relaxed Planning Graphs (Continued)

- ▶ a goal vertex g if  $v^n \in V^n$  for all  $v \in G$ , where n is last layer
- ▶ graph can be constructed for arbitrary many layers but stabilizes after a bounded number of layers  $\rightsquigarrow V^{i+1} = V^i$  and  $A^{i+1} = A^i$  (Why?)
- directed edges:
  - from  $v^i$  to  $a^{i+1}$  if  $v \in pre(a)$  (precondition edges)
  - from  $a^i$  to  $v^i$  if  $v \in add(a)$  (effect edges)
  - ightharpoonup from  $v^i$  to  $v^{i+1}$  (no-op edges)
  - ▶ from  $v^n$  to g if  $v \in G$  (goal edges)

German: Zielknoten, Vorbedingungskanten, Effektkanten, Zielkanten, No-Op-Kanten

# Illustrative Example

We write actions a with 
$$pre(a) = \{p_1, ..., p_k\}$$
,  $add(a) = \{q_1, ..., q_l\}$ ,  $del(a) = \emptyset$  and  $cost(a) = c$  as  $p_1, ..., p_k \stackrel{c}{\rightarrow} q_1, ..., q_l$ 

$$V = \{m, n, o, p, q, r, s, t\}$$

$$I = \{m\}$$

$$G = \{o, p, q, r, s\}$$

$$A = \{a_1, a_2, a_3, a_4, a_5, a_6\}$$

$$a_1 = m \xrightarrow{3} n, o$$

$$a_2 = m, o \xrightarrow{1} p$$

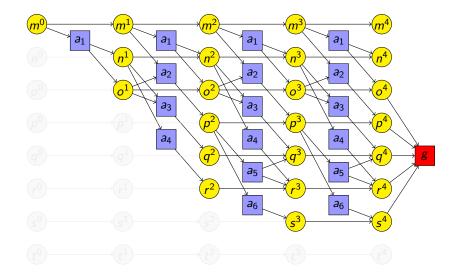
$$a_3 = n, o \xrightarrow{1} q$$

$$a_4 = n \xrightarrow{1} r$$

$$a_5 = p \xrightarrow{1} q, r$$

$$a_6 = p \xrightarrow{1} s$$

### Illustrative Example: Relaxed Planning Graph



# Generic Relaxed Planning Graph Heuristic

```
Heuristic Values from Relaxed Planning Graph
function generic-rpg-heuristic(\langle V, I, G, A \rangle, s):
     \Pi^+ := \langle V, s, G, A^+ \rangle
     for k \in \{0, 1, 2, \dots\}:
           rpg := RPG_k(\Pi^+) [relaxed planning graph to layer k]
           if rpg contains a goal node:
                Annotate nodes of rpg.
                if termination criterion is true:
                     return heuristic value from annotations
           else if graph has stabilized:
                return \infty
```

- → general template for RPG heuristics
- → to obtain concrete heuristic: instantiate highlighted elements

### Concrete Examples for Generic RPG Heuristic

Many planning heuristics fit this general template.

#### In this course:

- maximum heuristic h<sup>max</sup> (Bonet & Geffner, 1999)
- ▶ additive heuristic h<sup>add</sup> (Bonet, Loerincs & Geffner, 1997)
- Keyder & Geffner's (2008) variant of the FF heuristic h<sup>FF</sup> (Hoffmann & Nebel, 2001)

German: Maximum-Heuristik, additive Heuristik, FF-Heuristik

#### remark:

► The most efficient implementations of these heuristics do not use explicit planning graphs, but rather alternative (equivalent) definitions.

# F4.2 Maximum and Additive Heuristics

#### Maximum and Additive Heuristics

- $\blacktriangleright$   $h^{\text{max}}$  and  $h^{\text{add}}$  are the simplest RPG heuristics.
- Vertex annotations are numerical values.
- The vertex values estimate the costs
  - to make a given variable true
  - to reach and apply a given action
  - to reach the goal

### Maximum and Additive Heuristics: Filled-in Template

#### $h^{\text{max}}$ and $h^{\text{add}}$

#### computation of annotations:

- costs of variable vertices:0 in layer 0;otherwise minimum of the costs of predecessor vertices
- costs of action and goal vertices: maximum (h<sup>max</sup>) or sum (h<sup>add</sup>) of predecessor vertex costs; for action vertices a<sup>i</sup>, also add cost(a)

#### termination criterion:

**stability**: terminate if  $V^i = V^{i-1}$  and costs of all vertices in  $V^i$  equal corresponding vertex costs in  $V^{i-1}$ 

#### heuristic value:

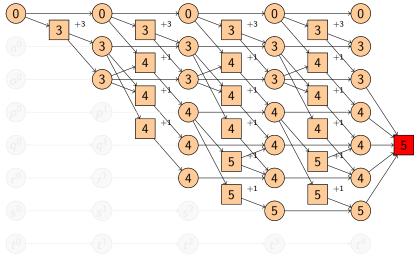
value of goal vertex

#### Maximum and Additive Heuristics: Intuition

#### intuition:

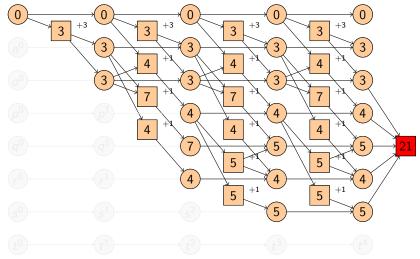
- variable vertices:
  - choose cheapest way of reaching the variable
- ► action/goal vertices:
  - h<sup>max</sup> is optimistic: assumption: when reaching the most expensive precondition variable, we can reach the other precondition variables in parallel (hence maximization of costs)
  - h<sup>add</sup> is pessimistic: assumption: all precondition variables must be reached completely independently of each other (hence summation of costs)

# Illustrative Example: $h^{max}$



 $h^{\max}(\{m\})=5$ 

# Illustrative Example: $h^{add}$



 $h^{\mathrm{add}}(\{m\}) = 21$ 

### $h^{\text{max}}$ and $h^{\text{add}}$ : Remarks

#### comparison of $h^{\text{max}}$ and $h^{\text{add}}$ :

- both are safe and goal-aware
- $ightharpoonup h^{\text{max}}$  is admissible and consistent;  $h^{\text{add}}$  is neither.
- $\rightarrow$   $h^{\text{add}}$  not suited for optimal planning
- ► However,  $h^{\text{add}}$  is usually much more informative than  $h^{\text{max}}$ . Greedy best-first search with  $h^{\text{add}}$  is a decent algorithm.
- ▶ Apart from not being admissible, h<sup>add</sup> often vastly overestimates the actual costs because positive synergies between subgoals are not recognized.
- → FF heuristic

# F4.3 FF Heuristic

#### FF Heuristic

#### The FF Heuristic

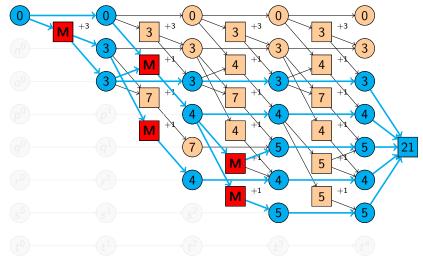
identical to  $h^{add}$ , but additional steps at the end:

- ► Mark goal vertex.
- Apply the following marking rules until nothing more to do:
  - marked action or goal vertex?
    - → mark all predecessors
  - - (tie-breaking: prefer variable vertices; otherwise arbitrary)

#### heuristic value:

- ► The actions corresponding to the marked action vertices build a relaxed plan.
- ► The cost of this plan is the heuristic value.

# Illustrative Example: hFF



$$h^{\mathsf{FF}}(\{m\}) = 3 + 1 + 1 + 1 + 1 = 7$$

#### FF Heuristic: Remarks

- ▶ Like h<sup>add</sup>, h<sup>FF</sup> is safe and goal-aware, but neither admissible nor consistent.
- $\triangleright$  approximation of  $h^+$  which is always at least as good as  $h^{add}$
- usually significantly better
- can be computed in almost linear time (O(n log n)) in the size of the description of the planning task
- computation of heuristic value depends on tie-breaking of marking rules (h<sup>FF</sup> not well-defined)
- ▶ one of the most successful planning heuristics

# Comparison of Relaxation Heuristics

#### Relationships of Relaxation Heuristics

Let s be a state in the STRIPS planning task  $\langle V, I, G, A \rangle$ .

#### Then

- $h^{\max}(s) \le h^+(s) \le h^*(s)$
- $h^{\max}(s) \le h^+(s) \le h^{\mathsf{FF}}(s) \le h^{\mathsf{add}}(s)$
- $\blacktriangleright$   $h^*$  and  $h^{FF}$  are incomparable
- $\blacktriangleright$   $h^*$  and  $h^{\text{add}}$  are incomparable

#### further remarks:

- For non-admissible heuristics, it is generally neither good nor bad to compute higher values than another heuristic.
- ▶ For relaxation heuristics, the objective is to approximate  $h^+$  as closely as possible.

# F4.4 Summary

# Summary

- Many delete relaxation heuristics can be viewed as computations on relaxed planning graphs (RPGs).
- examples: h<sup>max</sup>, h<sup>add</sup>, h<sup>FF</sup>
- $\blacktriangleright$   $h^{\text{max}}$  and  $h^{\text{add}}$  propagate numeric values in the RPGs
  - difference:  $h^{\text{max}}$  computes the maximum of predecessor costs for action and goal vertices;  $h^{\text{add}}$  computes the sum
- h<sup>FF</sup> marks vertices and sums the costs of marked action vertices.
- generally:  $h^{\max}(s) \le h^+(s) \le h^{\mathsf{FF}}(s) \le h^{\mathsf{add}}(s)$