## Foundations of Artificial Intelligence

B5. State-Space Search: Tree Search and Graph Search

Malte Helmert

University of Basel

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## State-Space Search: Overview

### Chapter overview: state-space search

- B1–B3. Foundations
- B4-B8. Basic Algorithms
  - B4. Data Structures for Search Algorithms
  - B5. Tree Search and Graph Search
  - B6. Breadth-first Search
  - B7. Uniform Cost Search
  - B8. Depth-first Search and Iterative Deepening
- B9–B15. Heuristic Algorithms

Introduction

# Introduction

## Search Algorithms

Introduction

### General Search Algorithm

iteratively create a search tree:

- starting with the initial state,
- repeatedly expand a state by generating its successors (which state depends on the used search algorithm)
- stop when a goal state is expanded (sometimes: generated)
- or all reachable states have been considered

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In this chapter, we study two essential classes of search algorithms:

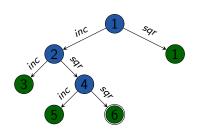
- tree search
- graph search

Each class consists of a large number of concrete algorithms.

German: expandieren, erzeugen, Baumsuche, Graphensuche

# Tree Search

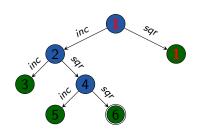
### Tree Search: General Idea



- possible paths to be explored organized in a tree (search tree)
- search nodes correspond 1:1 to paths from initial state
- duplicates a.k.a. transpositions (i.e., multiple nodes with identical state) possible
- search tree can have unbounded depth

German: Suchbaum, Duplikate, Transpositionen

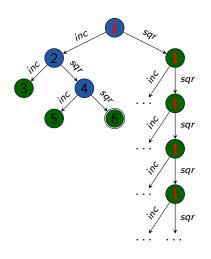
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### Generic Tree Search Algorithm

```
open := \mathbf{new} \ \mathsf{OpenList}
open.\mathsf{insert}(\mathsf{make\_root\_node}())
\mathbf{while} \ \mathbf{not} \ open.\mathsf{is\_empty}():
n := open.\mathsf{pop}()
\mathbf{if} \ \mathsf{is\_goal}(n.\mathsf{state}):
\mathbf{return} \ \mathsf{extract\_path}(n)
\mathbf{for} \ \mathbf{each} \ \langle a,s' \rangle \in \mathsf{succ}(n.\mathsf{state}):
n' := \mathsf{make\_node}(n,a,s')
open.\mathsf{insert}(n')
\mathbf{return} \ \mathsf{unsolvable}
```

## Generic Tree Search Algorithm: Discussion

#### discussion:

- generic template for tree search algorithms
- → for concrete algorithm, we must (at least) decide how to implement the open list
  - concrete algorithms often conceptually follow template, (= generate the same search tree), but deviate from details for efficiency reasons

# Graph Search

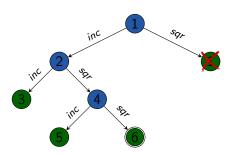
Graph Search

Graph Search

## Graph Search

#### differences to tree search:

- recognize duplicates: when a state is reached on multiple paths, only keep one search node
- search nodes correspond 1:1 to reachable states
- depth of search tree bounded



#### remarks:

- some graph search algorithms do not immediately eliminate all duplicates (→ later)
- one possible reason: find optimal solutions when a path to state s found later is cheaper than one found earlier

### Generic Graph Search Algorithm

```
open := new OpenList
open.insert(make_root_node())
closed := new ClosedList
while not open.is_empty():
     n := open.pop()
     if closed.lookup(n.state) = none:
           closed.insert(n)
          if is_goal(n.state):
                return extract_path(n)
          for each \langle a, s' \rangle \in \text{succ}(n.\text{state}):
                n' := \mathsf{make\_node}(n, a, s')
                open.insert(n')
return unsolvable
```

#### discussion:

- same comments as for generic tree search apply
- in "pure" algorithm, closed list does not actually need to store the search nodes
  - sufficient to implement closed as set of states
  - advanced algorithms often need access to the nodes, hence we show this more general version here
- some variants perform goal and duplicate tests elsewhere (earlier) → following chapters

# **Evaluating Search Algorithms**

four criteria for evaluating search algorithms:

### Completeness

Is the algorithm guaranteed to find a solution if one exists?

Does it terminate if no solution exists?

first property: semi-complete both properties: complete

German: Vollständigkeit, semi-vollständig, vollständig

## Criteria: Optimality

four criteria for evaluating search algorithms:

### Optimality

Are the solutions returned by the algorithm always optimal?

German: Optimalität

## Criteria: Time Complexity

four criteria for evaluating search algorithms:

### Time Complexity

How much time does the algorithm need until termination?

- usually worst case analysis
- usually measured in generated nodes

often a function of the following quantities:

- b: (branching factor) of state space (max. number of successors of a state)
- d: search depth (length of longest path in generated search tree)

German: Zeitaufwand, Verzweigungsgrad, Suchtiefe

## Criteria: Space Complexity

four criteria for evaluating search algorithms:

### Space Complexity

How much memory does the algorithm use?

- usually worst case analysis
- usually measured in (concurrently) stored nodes

often a function of the following quantities:

- b: (branching factor) of state space (max. number of successors of a state)
- d: search depth (length of longest path in generated search tree)

German: Speicheraufwand

## Analyzing the Generic Search Algorithms

### Generic Tree Search Algorithm

- Is it complete? Is it semi-complete?
- Is it optimal?
- What is its worst-case time complexity?
- What is its worst-case space complexity?

### Generic Graph Search Algorithm

- Is it complete? Is it semi-complete?
- Is it optimal?
- What is its worst-case time complexity?
- What is its worst-case space complexity?

# Summary

Summary

## Summary (1)

#### tree search:

• search nodes correspond 1:1 to paths from initial state

### graph search:

- search nodes correspond 1:1 to reachable states
- → duplicate elimination

generic methods with many possible variants

## Summary (2)

### evaluating search algorithms:

- completeness and semi-completeness
- optimality
- time complexity and space complexity