Algorithms and Data Structures A15. Sorting: Overview & Outlook

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Algorithms and Data Structures

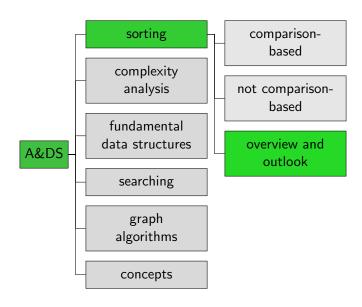
March 20, 2025 — A15. Sorting: Overview & Outlook

A15.1 Overview

A15.2 Outlook

A15.3 Quiz

Content of the Course



A15.1 Overview

Comparison-based Sorting: Overview

Algorithm	Running time $O(\cdot)$	Memory $O(\cdot)$	stable
	best/avg./worst	best/avg./worst	
Selection sort	n^2	1	no
Insertion sort	$n/n^2/n^2$	1	yes
Merge sort	$n \log n$	n	yes
Quicksort	$n \log n / n \log n / n^2$	$\log n/\log n/n$	no
Heapsort	$n \log n$	1	no

Very nice visualization of the algorithms at https://www.toptal.com/developers/sorting-algorithms/

Comparison-based Algorithms: Comments

- Insertion sort is very fast on short sequences and can be used to improve merge sort or quicksort for short ranges.
- Quicksort has a very short (= fast) inner loop. With randomization, the worst case almost never happens.
- Merge sort has the advantage of being stable.
 The merge step is also relevant for external sorting.
 Gets for example often used for database applications.
- Heapsort is in practise slightly slower than merge sort, but interesting because it is an in-place approach.
 e.g. for embedded systems.
- ▶ Equal asymptotic running time does not mean that algorithms take equally long (different hidden constants in $O(\cdot)$). Heapsort needs twice as many comparisons as merge sort.

A15.2 Outlook

Partially Sorted Data

- Often some subsequences of the input are already sorted (so-called runs).
- Insertion sort directly benefits from this.
- ► For some other approaches, there are variants that exploit runs, e.g. natural merge sort.

Many Equivalent Keys

- Quite common in practical applications.
 e.g. sorting students by place of residence
- There are special variants for some algorithms.
- ► For example, 3-way partitioning in quicksort

Sorting Complex Objects

- Most of the time, we do not want to sort numbers but complex objects.
- It would be extremely expensive to move them in memory for every swap.
- ▶ Instead: Sort elements that only consist of the key and a pointer/reference to the actual object.

Not So Correct Algorithms

INEFFECTIVE SORTS

```
DEFINE HALFHEARTED MERGESORT (LIST):

IF LENGTH (LIST) < 2:

RETIORN LIST

PYOT = INT (LENGTH (LIST) / 2)

A = HALFHEARTED MERGESORT (LIST[:PYOT])

B = HALFHEARTED MERGESORT (LIST[PWOT:])

// UMMMMMM

RETURN [A, B] // HERE. SORRY.
```

```
DEFINE FASTBOGOSORT(LIST):

// AN OPTIMIZED BOGOSORT

// RUNS IN O(NLOSN)
FOR N FROM 1 TO LOG(LENGTH(LIST));
SHUFFLE(LIST):
IF ISSORTED(LIST):
RETURN LIST

RETURN LIST

RETURN *KERNEL PAGE FAULT (ERROR CODE: 2)*
```

```
DEFINE JOBINTERVIEW QUICKSORT (LIST):
OK SO YOU CHOOSE A PIVOT
```

```
DEFINE PANICSORT(LIST):
```

```
full comic at https://xkcd.com/1185/
(CC BY-NC 2.5)
```

Solve other Problems by Sorting

k-smallest element

- For example, identifying the median $(k = \lfloor n/2 \rfloor)$.
- ► Use quicksort but only perform the recursive call for the relevant range (→ quickselect).

Duplicates

- ► How many different keys are there? Which value is most common? Are there duplicate keys?
- Can be solved directly with quadratic algorithms.
- Or more clever sort first and then use a single scan.

A15.3 Quiz

Quiz



kahoot.it