Foundations of Artificial Intelligence

E2. Propositional Logic: Equivalence and Normal Forms

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Propositional Logic: Overview

Chapter overview: propositional logic

- E1. Syntax and Semantics
- E2. Equivalence and Normal Forms
- E3. Reasoning and Resolution
- E4. DPLL Algorithm
- E5. Local Search and Outlook

Equivalence

Logical Equivalance

Definition (logically equivalent)

Formulas φ and ψ are called logically equivalent $(\varphi \equiv \psi)$ if for all interpretations $I: I \models \varphi$ iff $I \models \psi$.

German: logisch äquivalent

Equivalences

Logical Equivalences

Let φ , ψ , and η be formulas.

•
$$(\varphi \land \psi) \equiv (\psi \land \varphi)$$
 and $(\varphi \lor \psi) \equiv (\psi \lor \varphi)$ (commutativity)

•
$$((\varphi \land \psi) \land \eta) \equiv (\varphi \land (\psi \land \eta))$$
 and $((\varphi \lor \psi) \lor \eta) \equiv (\varphi \lor (\psi \lor \eta))$ (associativity)

•
$$((\varphi \land \psi) \lor \eta) \equiv ((\varphi \lor \eta) \land (\psi \lor \eta))$$
 and $((\varphi \lor \psi) \land \eta) \equiv ((\varphi \land \eta) \lor (\psi \land \eta))$ (distributivity)

•
$$\neg(\varphi \land \psi) \equiv (\neg \varphi \lor \neg \psi)$$
 and $\neg(\varphi \lor \psi) \equiv (\neg \varphi \land \neg \psi)$ (De Morgan)

•
$$\neg\neg\varphi\equiv\varphi$$
 (double negation)

•
$$(\varphi \to \psi) \equiv (\neg \varphi \lor \psi)$$
 $((\to)$ -elimination)

Commutativity and associativity are often used implicitly \rightsquigarrow We write $(X_1 \land X_2 \land X_3 \land X_4)$ instead of $(X_1 \land (X_2 \land (X_3 \land X_4)))$

Normal Forms

Normal Forms: Terminology

Definition (literal)

If $P \in \Sigma$, then the formulas P and $\neg P$ are called literals.

P is called positive literal, $\neg P$ is called negative literal.

The complementary literal to P is $\neg P$ and vice versa.

For a literal ℓ , the complementary literal to ℓ is denoted with $\bar{\ell}$.

German: Literal, positives/negatives/komplementäres Literal

Question: What is the difference between $\bar{\ell}$ and $\neg \ell$?

Normal Forms: Terminology

Definition (clause)

A disjunction of 0 or more literals is called a clause.

The empty clause (with 0 literals) is \perp .

Clauses consisting of exactly one literal are called unit clauses.

German: Klausel, leere Klausel, Einheitsklausel

Definition (monomial)

A conjunction of 0 or more literals is called a monomial.

German: Monom

Normal Forms

Definition (normal forms)

A formula φ is in conjunctive normal form (CNF, clause form) if φ is a conjunction of 0 or more clauses:

$$\varphi = \bigwedge_{i=1}^{n} \left(\bigvee_{j=1}^{m_i} \ell_{i,j} \right)$$

A formula φ is in disjunctive normal form (DNF) if φ is a disjunction of 0 or more monomials:

$$\varphi = \bigvee_{i=1}^{n} \left(\bigwedge_{j=1}^{m_i} \ell_{i,j} \right)$$

German: konjunktive Normalform, disjunktive Normalform

Normal Forms

For every propositional formula, there exists a logically equivalent propositional formula in CNF and in DNF.

Conversion to CNF with equivalences

• eliminate implications $(\varphi \to \psi) \equiv (\neg \varphi \lor \psi)$

$$((\rightarrow)$$
-elimination)

2 move negations inside

$$\neg(\varphi \land \psi) \equiv (\neg \varphi \lor \neg \psi) \qquad \text{(De Morgan)}
\neg(\varphi \lor \psi) \equiv (\neg \varphi \land \neg \psi) \qquad \text{(De Morgan)}
\neg \neg \varphi \equiv \varphi \qquad \text{(double negation)}$$

- \bullet simplify constant subformulas (\top, \bot)

There are formulas φ for which every logically equivalent formula in CNF and DNF is exponentially longer than φ .

Summary

Summary

- two formulas are logically equivalent if they have the same models
- different kinds of formulas:
 - atomic formulas and literals
 - clauses and monomials
 - conjunctive normal form (CNF) and disjunctive normal form (DNF)
- for every formula, there is a logically equivalent formula in CNF and a logically equivalent formula in DNF