# Foundations of Artificial Intelligence B3. State-Space Search: Examples of State Spaces

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## State-Space Search: Overview

### Chapter overview: state-space search

- B1–B3. Foundations
  - B1. State Spaces
  - B2. Representation of State Spaces
  - B3. Examples of State Spaces
- B4-B8. Basic Algorithms
- B9-B15. Heuristic Algorithms

## Three Examples

In this chapter we introduce three state spaces that we will use as illustrating examples:

- route planning in Romania
- a blocks world
- missionaries and cannibals

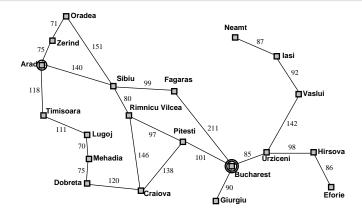
# Route Planning in Romania

## Route Planning in Romania

Route Planning in Romania

### Setting: Route Planning in Romania

We are on holiday in Romania and are currently located in Arad. Our flight home leaves from Bucharest. How to get there?



## Romania Formally

Route Planning in Romania

## State Space Route Planning in Romania

- states S: {arad, bucharest, craiova, . . . , zerind}
- actions A:  $move_{c,c'}$  for any two cities c and c' connected by a single road segment
- action costs cost: see figure, e.g.,  $cost(move_{iasi,vaslui}) = 92$
- transitions  $T: s \xrightarrow{a} s'$  iff  $a = move_{s,s'}$
- initial state:  $s_l = arad$
- goal states:  $S_G = \{ bucharest \}$

Blocks World

## Blocks World

## Blocks World

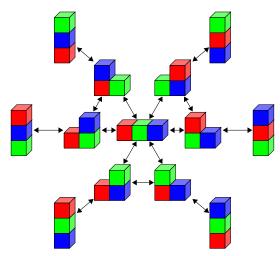
Blocks world is a traditional example problem in Al.

### Setting: Blocks World

- Colored blocks lie on a table.
- They can be stacked into towers, moving one block at a time.
- Our task is to create a given goal configuration.

## Example: Blocks World with Three Blocks

Action names omitted for readability. All actions cost 1. Initial state and goal can be arbitrary.



state space  $\langle S, A, cost, T, s_1, S_G \rangle$  for blocks world with n blocks

## State Space Blocks World

#### states S:

partitions of  $\{1, 2, ..., n\}$  into nonempty ordered lists

example n=3:

- $\{\langle 1, 2, 3 \rangle\}, \{\langle 1, 3, 2 \rangle\}, \{\langle 2, 1, 3 \rangle\},$  $\{\langle 2, 3, 1 \rangle\}, \{\langle 3, 1, 2 \rangle\}, \{\langle 3, 2, 1 \rangle\}$
- $\{\langle 1,2\rangle,\langle 3\rangle\},\{\langle 2,1\rangle,\langle 3\rangle\},\{\langle 1,3\rangle,\langle 2\rangle\},$  $\{\langle 3,1\rangle,\langle 2\rangle\},\{\langle 2,3\rangle,\langle 1\rangle\},\{\langle 3,2\rangle,\langle 1\rangle\}$
- $\{\langle 1 \rangle, \langle 2 \rangle, \langle 3 \rangle\}$

state space  $\langle S, A, cost, T, s_1, S_G \rangle$  for blocks world with *n* blocks

### State Space Blocks World

#### actions A:

- $\{move_{u,v} \mid u,v \in \{1,\ldots,n\} \text{ with } u \neq v\}$ 
  - move block u onto block v.
  - both must be uppermost blocks in their towers
- $\{to\text{-table}_u \mid u \in \{1, ..., n\}\}$ 
  - move block u onto the table (→ forming a new tower)
  - must be uppermost block in its tower

#### action costs cost:

cost(a) = 1 for all actions  $a \in A$ 

state space  $\langle S, A, cost, T, s_1, S_G \rangle$  for blocks world with *n* blocks

### State Space Blocks World

#### transitions:

- transition  $s \xrightarrow{a} s'$  with  $a = move_{u,v}$  exists iff
  - $s = \{\langle b_1, \dots, b_k, u \rangle, \langle c_1, \dots, c_m, v \rangle\} \cup X$  and
  - if k > 0:  $s' = \{\langle b_1, \dots, b_k \rangle, \langle c_1, \dots, c_m, v, u \rangle\} \cup X$
  - if k = 0:  $s' = \{\langle c_1, \dots, c_m, v, u \rangle\} \cup X$
- transition  $s \xrightarrow{a} s'$  with a = to-table, exists iff
  - $s = \{\langle b_1, \dots, b_k, u \rangle\} \cup X$  and
  - if k > 0:  $s' = \{\langle b_1, \ldots, b_k \rangle, \langle u \rangle\} \cup X$
  - if k = 0: s' = s

state space  $\langle S, A, cost, T, s_1, S_G \rangle$  for blocks world with *n* blocks

### State Space Blocks World

initial state  $s_1$  and goal states  $S_G$ :

one possible scenario for n = 3:

- $s_1 = \{\langle 1, 3 \rangle, \langle 2 \rangle\}$
- $S_G = \{\{\langle 3, 2, 1 \rangle\}\}$

(in general can have arbitrary scenarios)

## Blocks World: Properties

blocks	states	blocks	states
1	1	10	58941091
2	3	11	824073141
3	13	12	12470162233
4	73	13	202976401213
5	501	14	3535017524403
6	4051	15	65573803186921
7	37633	16	1290434218669921
8	394353	17	26846616451246353
9	4596553	18	588633468315403843

- For every given initial and goal state with n blocks, simple algorithms find a solution in time O(n). (How?)
- Finding optimal solutions is NP-complete (with a compact problem description).

## Missionaries and Cannibals

### Missionaries and Cannibals

## Setting: Missionaries and Cannibals

- Six people must cross a river.
- Their rowing boat can carry one or two people across the river at a time.
   (It is too small for three.)
- Three people are missionaries, three are cannibals.
- Missionaries may never stay with a majority of cannibals.



## Missionaries and Cannibals Formally

## State Space Missionaries and Cannibals

#### states S:

triples of numbers  $(m, c, b) \in \{0, 1, 2, 3\} \times \{0, 1, 2, 3\} \times \{0, 1\}$ :

- number of missionaries m,
- cannibals c and
- boats b

on the left river bank

initial state:  $s_l = \langle 3, 3, 1 \rangle$ 

goal:  $S_G = \{(0, 0, 0), (0, 0, 1)\}$ 

actions, action costs, transitions: ?

# Summary

## Summary

#### illustrating examples for state spaces:

- route planning in Romania:
  - small example of explicitly representable state space
- blocks world:
  - family of tasks where n blocks on a table must be rearranged
  - traditional example problem in AI
  - number of states explodes quickly as n grows
- missionaries and cannibals:
  - traditional brain teaser with small state space (32 states, of which many unreachable)