

# Algorithms and Data Structures

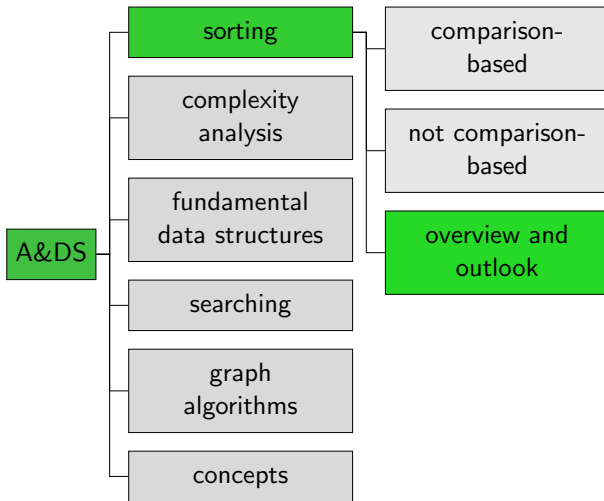
## A15. Sorting: Overview & Outlook

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March 21, 2024

# Content of the Course



# Overview

# Comparison-based Sorting: Overview

Algorithm	Running time $O(\cdot)$	Memory $O(\cdot)$	stable
	best/avg./worst	best/avg./worst	
Selection sort	$n^2$	1	no
Insertion sort	$n/n^2/n^2$	1	yes
Merge sort	$n \log n$	$n$	yes
Quicksort	$n \log n/n \log n/n^2$	$\log n/\log n/n$	no
Heap sort	$n \log n$	1	no

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Very nice [visualization of the algorithms](https://www.toptal.com/developers/sorting-algorithms/) at

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- Equal asymptotic running time does not mean that algorithms take equally long (different hidden constants in  $O(\cdot)$ ).  
Heapsort needs twice as many comparisons as merge sort.

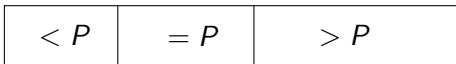
# Outlook

# Partially Sorted Data

- Often some subsequences of the input are already sorted (so-called runs).
- Insertion sort directly benefits from this.
- For some other approaches, there are variants that exploit runs, e.g. [natural merge sort](#).

## Many Equivalent Keys

- Quite common in practical applications.  
e.g. *sorting students by place of residence*
- There are special variants for some algorithms.
- For example, 3-way partitioning in quicksort



# Sorting Complex Objects

- Most of the time, we do not want to sort numbers but **complex objects**.
- It would be extremely expensive to move them in memory for every swap.
- **Instead:** Sort elements that only consist of the key and a pointer/reference to the actual object.

# Not So Correct Algorithms

## INEFFECTIVE SORTS

```
DEFINE HALFHEARTEDMERGESORT(LIST):  
  IF LENGTH(LIST) < 2:  
    RETURN LIST  
  PIVOT = INT(LENGTH(LIST) / 2)  
  A = HALFHEARTEDMERGESORT(LIST[:PIVOT])  
  B = HALFHEARTEDMERGESORT(LIST[PIVOT:])  
  // UMMMMMM  
  RETURN [A, B] // HERE. SORRY.
```

```
DEFINE FASTBOGOSORT(LIST):  
  // AN OPTIMIZED BOGOSORT  
  // RUNS IN  $O(N \log N)$   
  FOR N FROM 1 TO LOG(LENGTH(LIST)):  
    SHUFFLE(LIST):  
    IF ISSORTED(LIST):  
      RETURN LIST  
  RETURN "KERNEL PAGE FAULT (ERROR CODE: 2)"
```

```
DEFINE JOBIINTERVIEWQUICKSORT(LIST):  
  OK SO YOU CHOOSE A PIVOT  
  ...
```

```
DEFINE PANICSORT(LIST):  
  IF ISSORTED(LIST):
```

full comic at <https://xkcd.com/1185/>  
(CC BY-NC 2.5)

## Solve other Problems by Sorting

### *k*-smallest element

- For example, identifying the median ( $k = \lfloor n/2 \rfloor$ ).
- Use quicksort but only perform the recursive call for the relevant range ( $\rightarrow$  quickselect).



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### Duplicates

- How many different keys are there? Which value is most common? Are there duplicate keys?
- Can be solved directly with quadratic algorithms.
- Or – more clever – sort first and then use a single scan.

# Quiz

# Quiz



kahoot.it