Theory of Computer Science

B9. Context-free Languages: Closure & Decidability

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Pumping Lemma

Pumping Lemma for Context-free Languages

We used the pumping lemma from chapter B6 to show that a language is not regular. Is there a similar lemma for context-free languages?

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We used the pumping lemma from chapter B6 to show that a language is not regular. Is there a similar lemma for context-free languages?

Yes!

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Pumping Lemma for Context-free Languages

Pumping lemma for context-free languages:

- It is possible to prove a variant of the pumping lemma for context-free languages.
- Pumping is more complex than for regular languages:
  - word is decomposed into the form $uvwxy$ with $|vx| \geq 1$, $|vwx| \leq p$
  - pumped words have the form uv^iwx^iy
- This allows us to prove that certain languages are not context-free.
- example: $\{a^n b^n c^n \mid n \geq 1\}$ is not context-free (we will later use this without proof)

Closure Properties

Closure under Union, Concatenation, Star

Theorem

The context-free languages are closed under:

- *union*
- *concatenation*
- *star*

Closure under Union, Concatenation, Star: Proof

Proof.

Closed under union:

Let $G_1 = \langle V_1, \Sigma_1, R_1, S_1 \rangle$ and $G_2 = \langle V_2, \Sigma_2, R_2, S_2 \rangle$ be context-free grammars. W.l.o.g., $V_1 \cap V_2 = \emptyset$.

Then $\langle V_1 \cup V_2 \cup \{S\}, \Sigma_1 \cup \Sigma_2, R_1 \cup R_2 \cup \{S \to S_1, S \to S_2\}, S \rangle$ (where $S \notin V_1 \cup V_2$) is a context-free grammar for $\mathcal{L}(G_1) \cup \mathcal{L}(G_2)$.

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Closure under Union, Concatenation, Star: Proof

Proof (continued).

Closed under concatenation:

Let $G_1 = \langle V_1, \Sigma_1, R_1, S_1 \rangle$ and $G_2 = \langle V_2, \Sigma_2, R_2, S_2 \rangle$ be context-free grammars. W.l.o.g., $V_1 \cap V_2 = \emptyset$.

Then $\langle V_1 \cup V_2 \cup \{S\}, \Sigma_1 \cup \Sigma_2, R_1 \cup R_2 \cup \{S \to S_1 S_2\}, S \rangle$ (where $S \notin V_1 \cup V_2$) is a context-free grammar for $\mathcal{L}(G_1)\mathcal{L}(G_2)$.

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Closure under Union, Concatenation, Star: Proof

Proof (continued).

Closed under star:

Let $G = \langle V, \Sigma, R, S \rangle$ be a context-free grammar
where w.l.o.g. S never occurs on the right-hand side of a rule.

Then $G' = \langle V \cup \{S'\}, \Sigma, R', S' \rangle$ with $S' \notin V$ and
$R' = R \cup \{S' \to \varepsilon, S' \to S, S' \to SS'\}$ is a context-free grammar
for $\mathcal{L}(G)^*$. □

No Closure under Intersection or Complement

Theorem

The context-free languages are not closed under:

- *intersection*
- *complement*

No Closure under Intersection or Complement: Proof

Proof.

Not closed under intersection:

The languages $L_1 = \{\mathrm{a}^i \mathrm{b}^j \mathrm{c}^j \mid i, j \geq 1\}$
and $L_2 = \{\mathrm{a}^i \mathrm{b}^j \mathrm{c}^i \mid i, j \geq 1\}$ are context-free.

- For example, $G_1 = \langle \{\mathrm{S}, \mathrm{A}, \mathrm{X}\}, \{\mathtt{a}, \mathtt{b}, \mathtt{c}\}, R, \mathrm{S} \rangle$ with $R = \{\mathrm{S} \to \mathrm{AX}, \mathrm{A} \to \mathtt{a}, \mathrm{A} \to \mathtt{a}\mathrm{A}, \mathrm{X} \to \mathtt{bc}, \mathrm{X} \to \mathtt{b}\mathrm{X}\mathtt{c}\}$ is a context-free grammar for L_1.
- For example, $G_2 = \langle \{\mathrm{S}, \mathrm{B}\}, \{\mathtt{a}, \mathtt{b}, \mathtt{c}\}, R, \mathrm{S} \rangle$ with $R = \{\mathrm{S} \to \mathtt{a}\mathrm{S}\mathtt{c}, \mathrm{S} \to \mathrm{B}, \mathrm{B} \to \mathtt{b}, \mathrm{B} \to \mathtt{b}\mathrm{B}\}$ is a context-free grammar for L_2.

Their intersection is $L_1 \cap L_2 = \{\mathrm{a}^n \mathrm{b}^n \mathrm{c}^n \mid n \geq 1\}$.
We have remarked before that this language is not context-free.

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No Closure under Intersection or Complement: Proof

Proof (continued).

Not closed under complement:

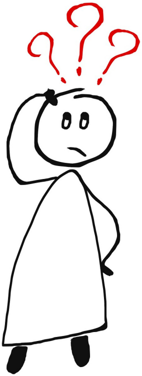
By contradiction: assume they were closed under complement.

Then they would also be closed under intersection
because they are closed under union and

$$L_1 \cap L_2 = \overline{\overline{L_1} \cup \overline{L_2}}.$$

This is a contradiction because we showed
that they are not closed under intersection. □

Questions



Questions?

Decidability

Word Problem

Definition (Word Problem for Context-free Languages)

The word problem $\mathrm{P}_\in$ for context-free languages is:

Given: context-free grammar G with alphabet Σ
and word $w \in \Sigma^*$

Question: Is $w \in \mathcal{L}(G)$?

Decidability: Word Problem

Theorem

The word problem $\mathrm{P}_\in$ for context-free languages is decidable.

Proof.

If $w = \varepsilon$, then $w \in \mathcal{L}(G)$ iff $S \to \varepsilon$ with start variable S is a rule of G.

Since for all other rules $w_\mathsf{l} \to w_\mathsf{r}$ of G we have $|w_\mathsf{l}| \leq |w_\mathsf{r}|$, the intermediate results when deriving a non-empty word never get shorter.

So it is possible to systematically consider all (finitely many) derivations of words up to length $|w|$ and test whether they derive the word w. □

Note: This is a terribly inefficient algorithm.

Emptiness Problem

Definition (Emptiness Problem for Context-free Languages)

The emptiness problem $\mathrm{P}_{\emptyset}$ for context-free languages is:

Given: context-free grammar G
Question: Is $\mathcal{L}(G) = \emptyset$?

Decidability: Emptiness Problem

Theorem

The emptiness problem for context-free languages is decidable.

Proof.

Given a grammar G, determine all variables in G that allow deriving words that only consist of terminal symbols:

- First mark all variables A for which a rule $A \to w$ exists such that w only consists of terminal symbols or $w = \varepsilon$.
- Then mark all variables A for which a rule $A \to w$ exists such that all nonterminal systems in w are already marked.
- Repeat this process until no further markings are possible.

$\mathcal{L}(G)$ is empty iff the start variable is unmarked at the end of this process. □

Finiteness Problem

Definition (Finiteness Problem for Context-free Languages)

The finiteness problem $\mathrm{P}_{\not\infty}$ for context-free languages is:

Given: context-free grammar G

Question: Is $|\mathcal{L}(G)| < \infty$?

Decidability: Finiteness Problem

Theorem

The finiteness problem for context-free languages is decidable.

We omit the proof. A possible proof uses the pumping lemma for context-free languages.

Proof sketch:

- We can compute certain bounds $l, u \in \mathbb{N}_0$ for a given context-free grammar G such that $\mathcal{L}(G)$ is infinite iff there exists $w \in \mathcal{L}(G)$ with $l \leq |w| \leq u$.
- Hence we can decide finiteness by testing all (finitely many) such words by using an algorithm for the word problem.

Intersection Problem

Definition (Intersection Problem for Context-free Languages)

The intersection problem $\mathrm{P}_\cap$ for context-free languages is:

Given: context-free grammars G and G'

Question: Is $\mathcal{L}(G) \cap \mathcal{L}(G') = \emptyset$?

Equivalence Problem

Definition (Equivalence Problem for Context-free Languages)

The equivalence problem $\mathrm{P}_=$ for context-free languages is:

Given: context-free grammars G and G'

Question: Is $\mathcal{L}(G) = \mathcal{L}(G')$?

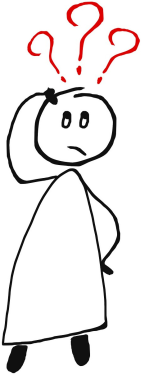
Undecidability: Equivalence and Intersection Problem

Theorem

The equivalence problem for context-free languages and the intersection problem for context-free languages are not decidable.

We cannot show this with the means currently available,
but we will get back to this in Part C (computability theory).

Questions



Questions?

Summary

Summary

- The context-free languages are closed under union, concatenation and star.
- The context-free languages are not closed under intersection or complement.
- The word problem, emptiness problem and finiteness problem for the class of context-free languages are decidable.
- The equivalence problem and intersection problem for the class of context-free languages are not decidable.

Further Topics on Context-free Languages and PDAs

- With the CYK-algorithm (Cocke, Younger and Kasami) it is possible to decide $w \in \mathcal{L}(G)$ in time $O(|w|^3)$ for a grammar in Chomsky normal form and a word w.
- Deterministic push-down automata have the restriction $|\delta(q, a, A)| + |\delta(q, \varepsilon, A)| \leq 1$ for all $q \in Q, a \in \Sigma, A \in \Gamma$.
- The languages recognized by deterministic PDAs are called deterministic context-free languages. They form a strict superset of the regular languages and a strict subset of the context-free languages.