

Theory of Computer Science

A3. Proof Techniques

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Introduction

What is a Proof?

A **mathematical proof** is

- a sequence of logical steps
- starting with one set of statements
- that comes to the conclusion
that some statement must be true.

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What is a **statement**?

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The statement is **true** if the conclusions are true whenever the preconditions are true.

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Notes:

- set of preconditions is sometimes empty
- often, “assumptions” is used instead of “preconditions”; slightly unfortunate because “assumption” is also used with another meaning (\rightsquigarrow cf. indirect proofs)

Examples of Mathematical Statements

Examples (some true, some false):

- “Let $p \in \mathbb{N}_0$ be a prime number. Then p is odd.”
- “There exists an even prime number.”
- “Let $p \in \mathbb{N}_0$ with $p \geq 3$ be a prime number. Then p is odd.”
- “All prime numbers $p \geq 3$ are odd.”
- “For all sets A, B, C : $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$ ”

What are the preconditions, what are the conclusions?

On what Statements can we Build the Proof?

A mathematical proof is

- a sequence of logical steps
- **starting with one set of statements**
- that comes to the conclusion
that some statement must be true.

We can use:

- **axioms**: statements that are assumed to always be true
in the current context
- **theorems** and **lemmas**: statements that were already proven
 - lemma: an intermediate tool
 - theorem: itself a relevant result
- **premises**: assumptions we make to see
what consequences they have

What is a Logical Step?

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Each step **directly follows**

- from the axioms,
- premises,
- previously proven statements and
- the preconditions of the statement we want to prove.

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For a formal definition, we would need formal logics.

The Role of Definitions

Definition

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- A definition introduces an abbreviation.
- Whenever we say “set”, we could instead say “an unordered collection of distinct objects” and vice versa.
- Definitions can also introduce notation.

Disproofs

- A **disproof** (**refutation**) shows that a given mathematical statement is **false** by giving an example where the preconditions are true, but the conclusion is false.
- This requires deriving, in a sequence of proof steps, the opposite (negation) of the conclusion.
- Formally, disproofs are proofs of modified (“negated”) statements.
- Be careful about how to negate a statement!

Proof Strategies

typical proof/disproof strategies:

- 1 “All $x \in S$ with the property P also have the property Q .”
“For all $x \in S$: if x has property P , then x has property Q .”
 - To prove, assume you are given an arbitrary $x \in S$ that has the property P .
Give a sequence of proof steps showing that x must have the property Q .
 - To disprove, find a **counterexample**, i. e., find an $x \in S$ that has property P but not Q and prove this.

Proof Strategies

typical proof/disproof strategies:

- ② “ A is a subset of B .”
 - To prove, assume you have an arbitrary element $x \in A$ and prove that $x \in B$.
 - To disprove, find an element in $x \in A \setminus B$ and prove that $x \in A \setminus B$.

Proof Strategies

typical proof/disproof strategies:

- ③ “For all $x \in S$: x has property P iff x has property Q .”
 (“iff”: “if and only if”)
 - To prove, separately prove “if P then Q ” and “if Q then P ”.
 - To disprove, disprove “if P then Q ” or disprove “if Q then P ”.

Proof Strategies

typical proof/disproof strategies:

- ④ “ $A = B$ ”, where A and B are sets.
 - To prove, separately prove “ $A \subseteq B$ ” and “ $B \subseteq A$ ”.
 - To disprove, disprove “ $A \subseteq B$ ” or disprove “ $B \subseteq A$ ”.

Proof Techniques

most common proof techniques:

- direct proof
- indirect proof (proof by contradiction)
- proof by contrapositive
- mathematical induction

Exercise

You want to disprove the following statement with a counterexample:

If the sun is shining then all kids eat ice cream.

What properties must your counterexample have?



Direct Proof

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Direct Proof

Direct derivation of the statement by deducing or rewriting.

Direct Proof: Example

Theorem (distributivity)

For all sets A, B, C : $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$.

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We first show that $x \in A \cap (B \cup C)$ implies
 $x \in (A \cap B) \cup (A \cap C)$ (\subseteq part):

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Let $x \in A \cap (B \cup C)$. Then by the definition of \cap it holds that
 $x \in A$ and $x \in B \cup C$.

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Let $x \in A \cap (B \cup C)$. Then by the definition of \cap it holds that
 $x \in A$ and $x \in B \cup C$.

We make a case distinction between $x \in B$ and $x \notin B$:

If $x \in B$ then, because $x \in A$ is true, $x \in A \cap B$ must be true.

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If $x \in B$ then, because $x \in A$ is true, $x \in A \cap B$ must be true.

Otherwise, because $x \in B \cup C$ we know that $x \in C$ and thus with
 $x \in A$, that $x \in A \cap C$.

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Let $x \in A \cap (B \cup C)$. Then by the definition of \cap it holds that $x \in A$ and $x \in B \cup C$.

We make a case distinction between $x \in B$ and $x \notin B$:

If $x \in B$ then, because $x \in A$ is true, $x \in A \cap B$ must be true.

Otherwise, because $x \in B \cup C$ we know that $x \in C$ and thus with $x \in A$, that $x \in A \cap C$.

In both cases $x \in A \cap B$ or $x \in A \cap C$,
and we conclude $x \in (A \cap B) \cup (A \cap C)$

Direct Proof: Example

Theorem (distributivity)

For all sets A, B, C : $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$.

Proof (continued).

\supseteq part: we must show that $x \in (A \cap B) \cup (A \cap C)$ implies $x \in A \cap (B \cup C)$.

Let $x \in (A \cap B) \cup (A \cap C)$.

Direct Proof: Example

Theorem (distributivity)

For all sets A, B, C : $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$.

Proof (continued).

\supseteq part: we must show that $x \in (A \cap B) \cup (A \cap C)$ implies $x \in A \cap (B \cup C)$.

Let $x \in (A \cap B) \cup (A \cap C)$.

We make a case distinction between $x \in A \cap B$ and $x \notin A \cap B$:

If $x \in A \cap B$ then $x \in A$ and $x \in B$.

The latter implies $x \in B \cup C$ and hence $x \in A \cap (B \cup C)$.

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Let $x \in (A \cap B) \cup (A \cap C)$.

We make a case distinction between $x \in A \cap B$ and $x \notin A \cap B$:

If $x \in A \cap B$ then $x \in A$ and $x \in B$.

The latter implies $x \in B \cup C$ and hence $x \in A \cap (B \cup C)$.

If $x \notin A \cap B$ we know $x \in A \cap C$ due to $x \in (A \cap B) \cup (A \cap C)$.

This (analogously) implies $x \in A$ and $x \in C$, and hence $x \in B \cup C$ and thus $x \in A \cap (B \cup C)$.

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Let $x \in (A \cap B) \cup (A \cap C)$.

We make a case distinction between $x \in A \cap B$ and $x \notin A \cap B$:

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The latter implies $x \in B \cup C$ and hence $x \in A \cap (B \cup C)$.

If $x \notin A \cap B$ we know $x \in A \cap C$ due to $x \in (A \cap B) \cup (A \cap C)$.

This (analogously) implies $x \in A$ and $x \in C$, and hence $x \in B \cup C$ and thus $x \in A \cap (B \cup C)$.

In both cases we conclude $x \in A \cap (B \cup C)$.

...

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For all sets A, B, C : $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$.

Proof (continued).

We have shown that every element of $A \cap (B \cup C)$ is an element of $(A \cap B) \cup (A \cap C)$ and vice versa. Thus, both sets are equal. □

Direct Proof: Example

Theorem (distributivity)

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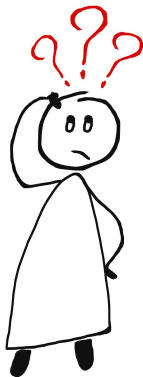
Proof.

Alternative:

$$\begin{aligned} A \cap (B \cup C) &= \{x \mid x \in A \text{ and } x \in B \cup C\} \\ &= \{x \mid x \in A \text{ and } (x \in B \text{ or } x \in C)\} \\ &= \{x \mid (x \in A \text{ and } x \in B) \text{ or } (x \in A \text{ and } x \in C)\} \\ &= \{x \mid x \in A \cap B \text{ or } x \in A \cap C\} \\ &= (A \cap B) \cup (A \cap C) \end{aligned}$$



Questions



Questions?

Indirect Proof

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Indirect Proof (Proof by Contradiction)

- Make an **assumption** that the statement is false.
- Derive a **contradiction** from the assumption together with the preconditions of the statement.
- This shows that the assumption must be false given the preconditions of the statement, and hence the original statement must be true.

Indirect Proof: Example

Theorem

There are infinitely many prime numbers.

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Proof.

Assumption: There are only finitely many prime numbers.



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Let $P = \{p_1, \dots, p_n\}$ be the set of all prime numbers.

Define $m = p_1 \cdot \dots \cdot p_n + 1$.



Indirect Proof: Example

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Assumption: There are only finitely many prime numbers.

Let $P = \{p_1, \dots, p_n\}$ be the set of all prime numbers.

Define $m = p_1 \cdot \dots \cdot p_n + 1$.

Since $m \geq 2$, it must have a prime factor.

Let p be such a prime factor.



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Since p is a prime number, p has to be in P .



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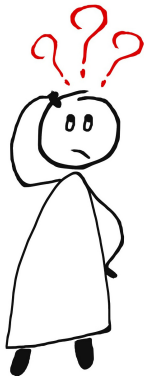
Since p is a prime number, p has to be in P .

The number m is not divisible without remainder by any of the numbers in P . Hence p is no factor of m .

↪ **Contradiction**



Questions



Questions?

Contrapositive

Proof by Contrapositive

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Prove “If A , then B ” by proving “If not B , then not A .”

Proof by Contrapositive

Proof by Contrapositive

Prove “If A , then B ” by proving “If not B , then not A .”

Examples:

- Prove “For all $n \in \mathbb{N}_0$: if n^2 is odd, then n is odd” by proving “For all $n \in \mathbb{N}_0$, if n is even, then n^2 is even.”
- Prove “For all $n \in \mathbb{N}_0$: if n is not a square number, then \sqrt{n} is irrational” by proving “For all $n \in \mathbb{N}_0$: if \sqrt{n} is rational, then n is a square number.”

Exercise

How would you prove the following statement by contrapositive:

If the sun is shining then all kids eat ice cream.



Mathematical Induction

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Mathematical Induction

Proof of a statement for all natural numbers n with $n \geq m$

- **basis**: proof of the statement for $n = m$
- **induction hypothesis (IH)**:
suppose that the statement is true for all k with $m \leq k \leq n$
- **inductive step**: proof of the statement for $n + 1$
using the induction hypothesis

Mathematical Induction: Example

Theorem

For all $n \in \mathbb{N}_0$ with $n \geq 1$: $\sum_{k=1}^n (2k - 1) = n^2$

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Proof.

Mathematical induction over n :

basis $n = 1$: $\sum_{k=1}^1 (2k - 1) = 2 - 1 = 1 = 1^2$



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Proof.

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IH: $\sum_{k=1}^m (2k - 1) = m^2$ for all $1 \leq m \leq n$



Mathematical Induction: Example

Theorem

For all $n \in \mathbb{N}_0$ with $n \geq 1$: $\sum_{k=1}^n (2k - 1) = n^2$

Proof.

Mathematical induction over n :

basis $n = 1$: $\sum_{k=1}^1 (2k - 1) = 2 - 1 = 1 = 1^2$

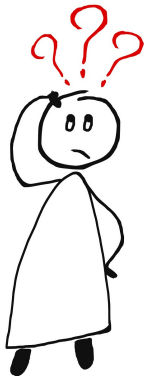
IH: $\sum_{k=1}^m (2k - 1) = m^2$ for all $1 \leq m \leq n$

inductive step $n \rightarrow n + 1$:

$$\begin{aligned}\sum_{k=1}^{n+1} (2k - 1) &= \left(\sum_{k=1}^n (2k - 1) \right) + 2(n + 1) - 1 \\ &\stackrel{\text{IH}}{=} n^2 + 2(n + 1) - 1 \\ &= n^2 + 2n + 1 = (n + 1)^2\end{aligned}$$



Questions



Questions?

Summary

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- A **proof** is based on axioms and previously proven statements.
- Individual **proof steps** must be obvious derivations.
- **direct proof**: sequence of derivations or rewriting
- **indirect proof**: refute the negated statement
- **contrapositive**: prove " $A \Rightarrow B$ " as " $\text{not } B \Rightarrow \text{not } A$ "
- **mathematical induction**: prove statement for a starting point and show that it always carries over to the next number