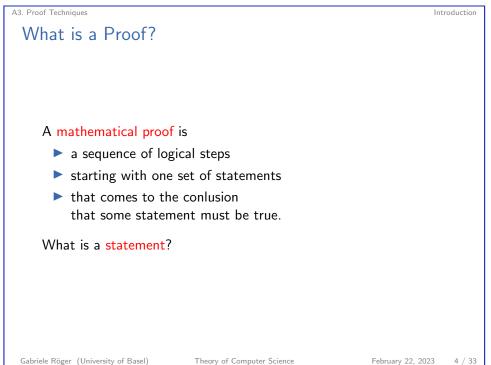


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# Theory of Computer Science February 22, 2023 — A3. Proof Techniques A3.1 Introduction A3.2 Direct Proof A3.3 Indirect Proof A3.4 Contrapositive A3.5 Mathematical Induction A3.6 Summary Gabriele Röger (University of Basel) Theory of Computer Science February 22, 2023 2 / 33



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### Mathematical Statements

Mathematical Statement

A mathematical statement consists of a set of preconditions and a set of conclusions.

The statement is **true** if the conclusions are true whenever the preconditions are true.

#### Notes:

- set of preconditions is sometimes empty
- often, "assumptions" is used instead of "preconditions"; slightly unfortunate because "assumption" is also used with another meaning (~> cf. indirect proofs)

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A3. Proof Techniques

## On what Statements can we Build the Proof?

### A mathematical proof is

- ► a sequence of logical steps
- starting with one set of statements
- that comes to the conlusion that some statement must be true.

### We can use:

- axioms: statements that are assumed to always be true in the current context
- theorems and lemmas: statements that were already proven
  - lemma: an intermediate tool
  - theorem: itself a relevant result
- premises: assumptions we make to see what consequences they have



### A3. Proof Techniques

Introduction

## Examples of Mathematical Statements

### Examples (some true, some false):

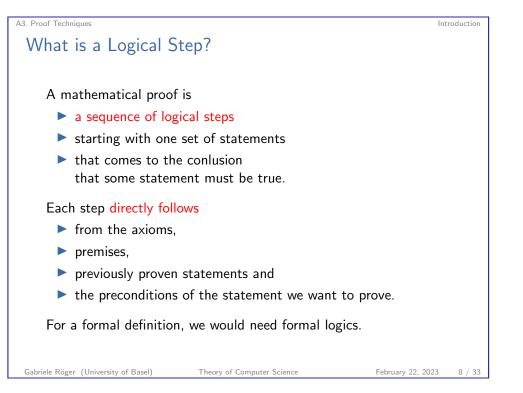
- "Let  $p \in \mathbb{N}_0$  be a prime number. Then p is odd."
- "There exists an even prime number."
- "Let  $p \in \mathbb{N}_0$  with  $p \ge 3$  be a prime number. Then p is odd."
- "All prime numbers  $p \ge 3$  are odd."
- ▶ "For all sets A, B, C:  $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$ "

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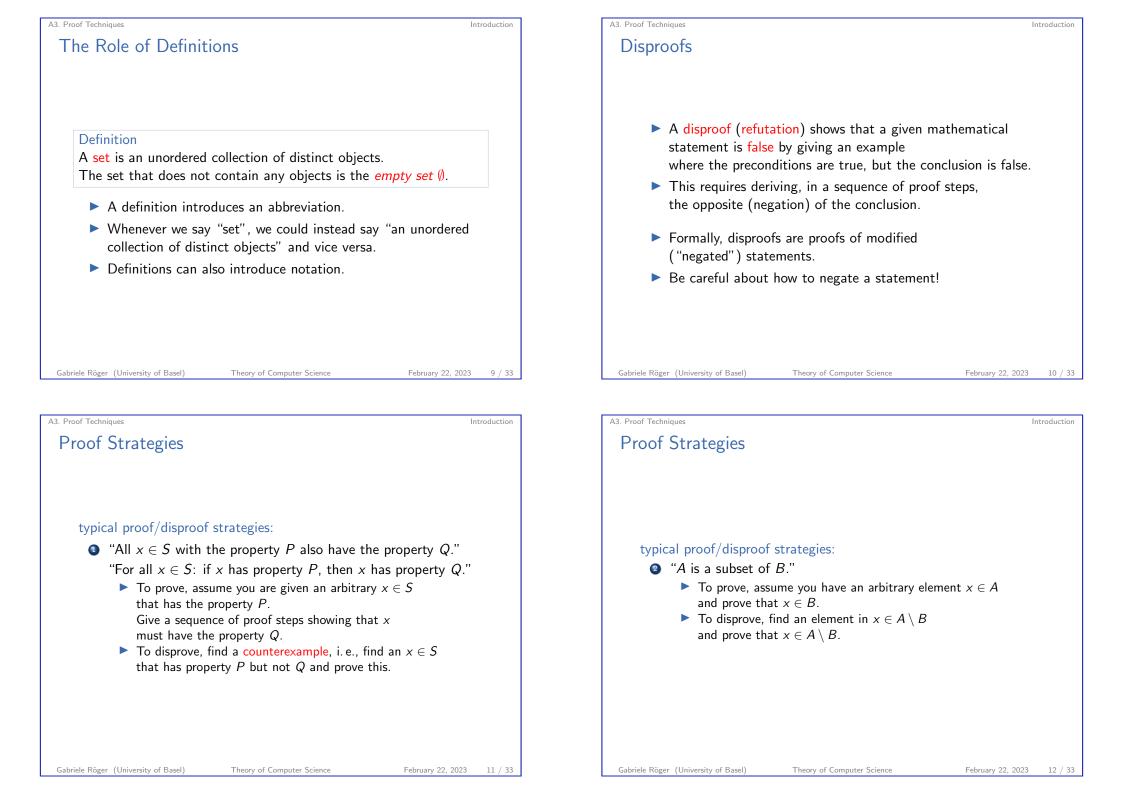
What are the preconditions, what are the conclusions?

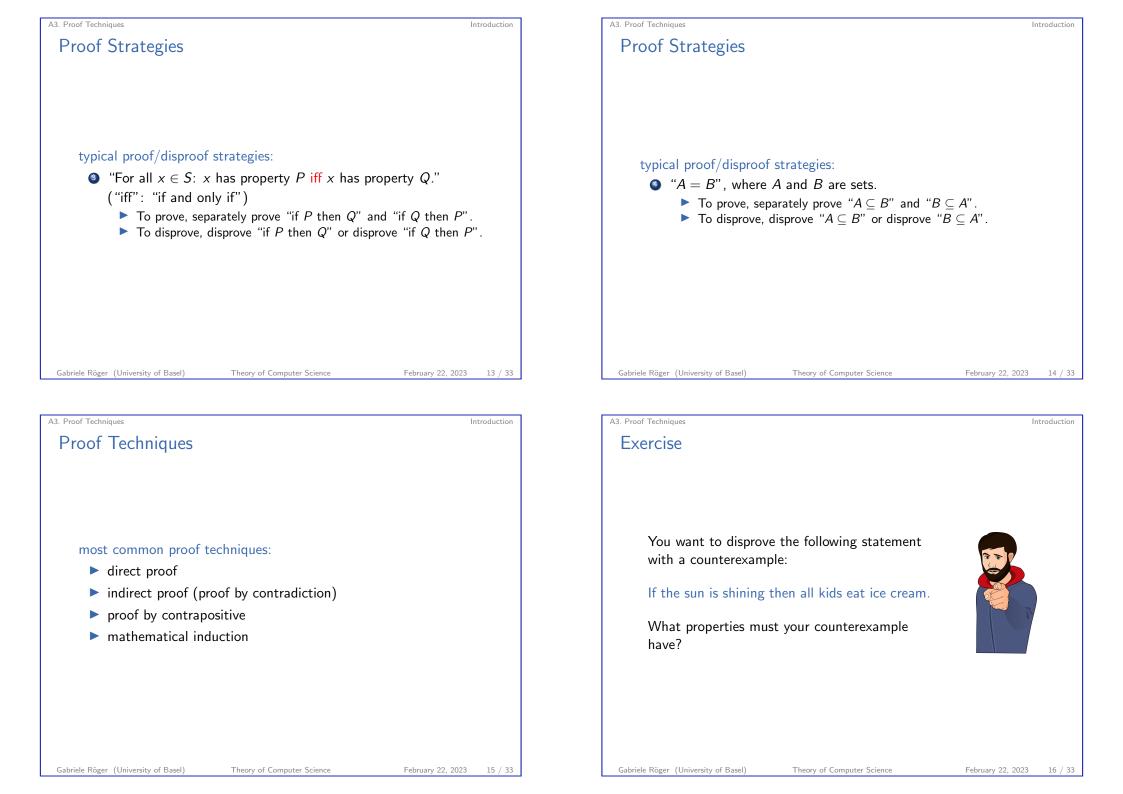
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A3.2 Direct	Proof		
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#### A3. Proof Techniques

Direct Proof: Example

Theorem (distributivity)

For all sets A, B, C:  $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$ .

### Proof.

We first show that  $x \in A \cap (B \cup C)$  implies  $x \in (A \cap B) \cup (A \cap C) (\subseteq part)$ :

Let  $x \in A \cap (B \cup C)$ . Then by the definition of  $\cap$  it holds that  $x \in A$  and  $x \in B \cup C$ .

We make a case distinction between  $x \in B$  and  $x \notin B$ :

If  $x \in B$  then, because  $x \in A$  is true,  $x \in A \cap B$  must be true.

Otherwise, because  $x \in B \cup C$  we know that  $x \in C$  and thus with  $x \in A$ , that  $x \in A \cap C$ .

```
In both cases x \in A \cap B or x \in A \cap C,
and we conclude x \in (A \cap B) \cup (A \cap C).
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Direct Proof

. . .

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A3. Proof Techniques		D	irect Proof
Direct Proof			
Direct Proof			
Direct derivation of th	e statement by deducing o	or rewriting.	
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## Direct Proof: Example

Theorem (distributivity) For all sets A, B, C:  $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$ .

### Proof (continued).

 $\supseteq$  part: we must show that  $x \in (A \cap B) \cup (A \cap C)$  implies  $x \in A \cap (B \cup C).$ 

Let  $x \in (A \cap B) \cup (A \cap C)$ .

We make a case distinction between  $x \in A \cap B$  and  $x \notin A \cap B$ :

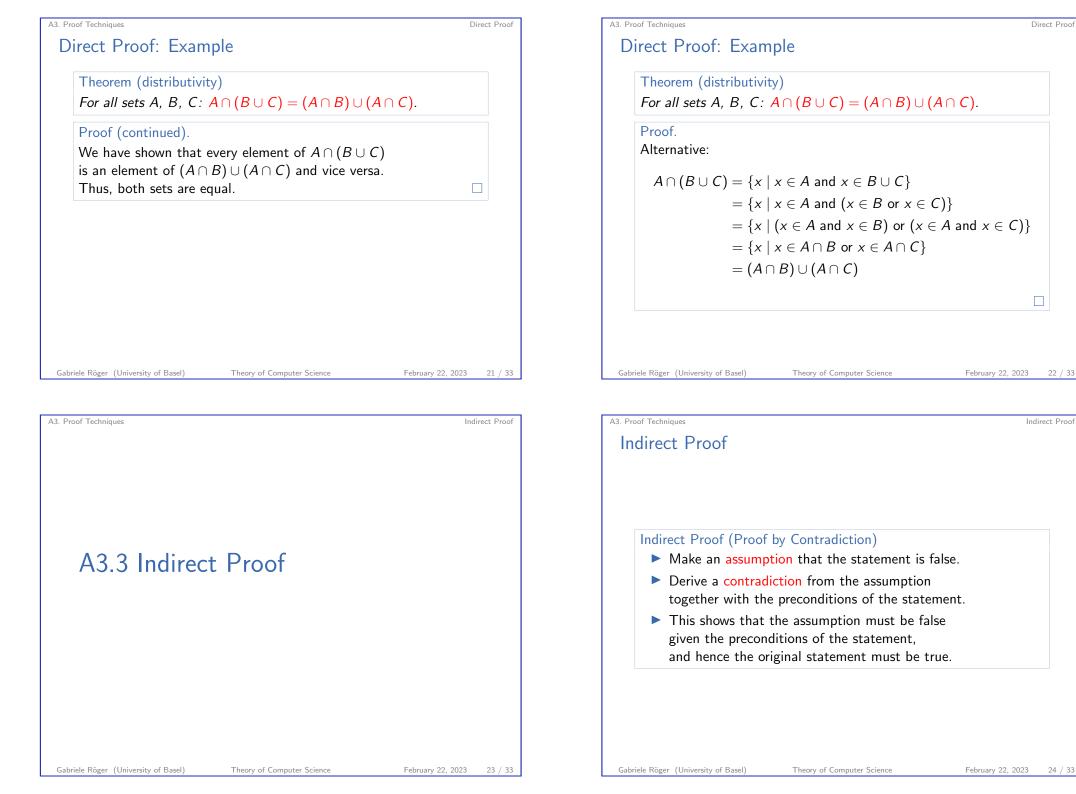
If  $x \in A \cap B$  then  $x \in A$  and  $x \in B$ .

The latter implies  $x \in B \cup C$  and hence  $x \in A \cap (B \cup C)$ .

If  $x \notin A \cap B$  we know  $x \in A \cap C$  due to  $x \in (A \cap B) \cup (A \cap C)$ . This (analogously) implies  $x \in A$  and  $x \in C$ , and hence  $x \in B \cup C$ and thus  $x \in A \cap (B \cup C)$ .

In both cases we conclude  $x \in A \cap (B \cup C)$ .

. . .



Indirect Proof

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Contrapositive

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### Indirect Proof: Example

#### Theorem

There are infinitely many prime numbers.

### Proof.

Assumption: There are only finitely many prime numbers. Let  $P = \{p_1, \ldots, p_n\}$  be the set of all prime numbers. Define  $m = p_1 \cdot \ldots \cdot p_n + 1$ . Since  $m \ge 2$ , it must have a prime factor. Let p be such a prime factor. Since p is a prime number, p has to be in P. The number *m* is not divisible without remainder by any of the numbers in P. Hence p is no factor of m. → Contradiction

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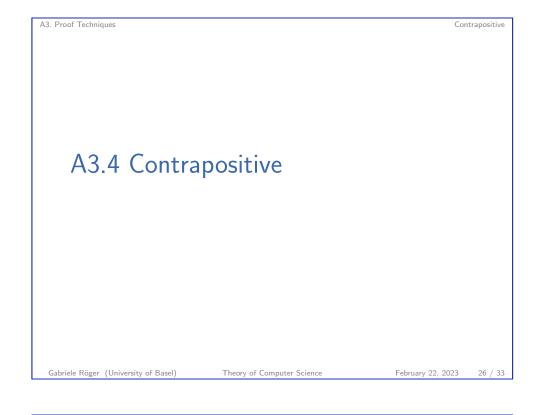
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A3. Proof Techniques Proof by Contrapositive

Proof by Contrapositive Prove "If A, then B" by proving "If not B, then not A."

### Examples:

- ▶ Prove "For all  $n \in \mathbb{N}_0$ : if  $n^2$  is odd, then *n* is odd" by proving "For all  $n \in \mathbb{N}_0$ , if *n* is even, then  $n^2$  is even."
- ▶ Prove "For all  $n \in \mathbb{N}_0$ : if *n* is not a square number, then  $\sqrt{n}$  is irrational" by proving "For all  $n \in \mathbb{N}_0$ : if  $\sqrt{n}$  is rational, then *n* is a square number."





A3. Proof Techniques		Mathematical	Induction	
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#### A3. Proof Techniques

Mathematical Induction: Example

Theorem

For all  $n \in \mathbb{N}_0$  with  $n \ge 1$ :  $\sum_{k=1}^n (2k-1) = n^2$ 

### Proof.

Mathematical induction over *n*:

basis 
$$n = 1$$
:  $\sum_{k=1}^{1} (2k - 1) = 2 - 1 = 1 = 1^2$   
IH:  $\sum_{k=1}^{m} (2k - 1) = m^2$  for all  $1 \le m \le n$   
inductive step  $n \to n + 1$ :

$$\sum_{k=1}^{n+1} (2k-1) = \left(\sum_{k=1}^{n} (2k-1)\right) + 2(n+1) - 1$$
$$\stackrel{\text{IH}}{=} n^2 + 2(n+1) - 1$$
$$= n^2 + 2n + 1 = (n+1)^2$$

A3. Proof Techniques

### Mathematical Induction

### Mathematical Induction

Proof of a statement for all natural numbers n with  $n \ge m$ 

- **basis**: proof of the statement for n = m
- induction hypothesis (IH): suppose that the statement is true for all k with m ≤ k ≤ n
   inductive step: proof of the statement for n + 1 using the induction hypothesis

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A3. Proof Techniques

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Summary

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Mathematical Induction

# A3.6 Summary

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Mathematical Induction



Summary

### Summary

- ► A proof is based on axioms and previously proven statements.
- Individual proof steps must be obvious derivations.
- direct proof: sequence of derivations or rewriting
- indirect proof: refute the negated statement
- contrapositive: prove " $A \Rightarrow B$ " as "not  $B \Rightarrow \text{not } A$ "
- mathematical induction: prove statement for a starting point and show that it always carries over to the next number

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