

# Theory of Computer Science

## D6. Beyond NP

Gabriele Röger

University of Basel

May 17, 2021

# Theory of Computer Science

May 17, 2021 — D6. Beyond NP

## D6.1 coNP

## D6.2 Time and Space Complexity

## D6.3 Polynomial Hierarchy

## D6.4 Counting

## Complexity Theory: What we already have seen

- ▶ **Complexity theory** investigates which problems are “easy” to solve and which ones are “hard”.
- ▶ two important problem classes:
  - ▶ **P**: problems that are solvable in **polynomial time** by “**normal**” **computation mechanisms**
  - ▶ **NP**: problems that are solvable in **polynomial time** with the help of **nondeterminism**
- ▶ We know that  $P \subseteq NP$ , but we do not know whether  $P = NP$ .
- ▶ Many practically relevant problems are **NP-complete**:
  - ▶ They belong to NP.
  - ▶ All problems in NP can be polynomially reduced to them.
- ▶ If there is an efficient algorithm for **one** NP-complete problem, then there are efficient algorithms for **all** problems in NP.

D6. Beyond NP

coNP

## D6.1 coNP

## Complexity Class coNP

### Definition (coNP)

**coNP** is the set of all languages  $L$  for which  $\bar{L} \in \text{NP}$ .

**Example:** The complement of SAT is in coNP.

## Hardness and Completeness

### Definition (Hardness and Completeness)

Let  $C$  be a complexity class.

A problem  $Y$  is called **C-hard** if  $X \leq_p Y$  for **all** problems  $X \in C$ .

$Y$  is called **C-complete** if  $Y \in C$  and  $Y$  is C-hard.

### Example (TAUTOLOGY)

The following problem **TAUTOLOGY** is coNP-complete:

**Given:** a propositional logic formula  $\varphi$

**Question:** Is  $\varphi$  valid, i.e. is it true under all variable assignments?

## Known Results and Open Questions

### Open

- ▶  $\text{NP} \stackrel{?}{=} \text{coNP}$

### Known

- ▶  $P \subseteq \text{coNP}$
- ▶ If  $X$  is NP-complete then  $\bar{L}$  is coNP-complete.
- ▶ If  $\text{NP} \neq \text{coNP}$  then  $P \neq \text{NP}$ .
- ▶ If a coNP-complete problem is in NP, then  $\text{NP} = \text{coNP}$ .
- ▶ If a coNP-complete problem is in P, then  $P = \text{coNP} = \text{NP}$ .

## D6.2 Time and Space Complexity

## Reminder: Time Complexity Classes

### Definition (Time Complexity Classes TIME and NTIME)

Let  $t : \mathbb{N} \rightarrow \mathbb{R}^+$  be a function.

The **time complexity class**  $\text{TIME}(t(n))$  is the collection of all languages that are decidable by an  $O(t)$  **time Turing machine**, and  $\text{NTIME}(t(n))$  is the collection of all languages that are decidable by an  $O(t)$  **time nondeterministic Turing machine**.

- ▶  $\text{TIME}(f)$ : all languages accepted by a **DTM** in time  $f$ .
- ▶  $\text{NTIME}(f)$ : all languages accepted by a **NTM** in time  $f$ .
- ▶  $P = \bigcup_{k \in \mathbb{N}} \text{TIME}(n^k)$
- ▶  $NP = \bigcup_{k \in \mathbb{N}} \text{NTIME}(n^k)$

## Space

- ▶ **Analogously**: A TM decides a language  $L$  in **space**  $f$  if the computation on every input visits at most  $f(|w|)$  tape cells besides its input on the tape.
- ▶  $\text{SPACE}(f)$ : all languages decided by a **DTM** in space  $f$ .
- ▶  $\text{NSPACE}(f)$ : all languages decided by a **NTM** in space  $f$ .

## Important Complexity Classes Beyond NP

- ▶  $\text{PSPACE} = \bigcup_{k \in \mathbb{N}} \text{SPACE}(n^k)$
- ▶  $\text{NPSPACE} = \bigcup_{k \in \mathbb{N}} \text{NSPACE}(n^k)$
- ▶  $\text{EXPTIME} = \bigcup_{k \in \mathbb{N}} \text{TIME}(2^{n^k})$
- ▶  $\text{EXPSPACE} = \bigcup_{k \in \mathbb{N}} \text{SPACE}(2^{n^k})$

Some known results:

- ▶  $\text{PSPACE} = \text{NPSPACE}$  (from Savitch's theorem)
- ▶  $\text{PSPACE} \subseteq \text{EXPTIME} \subseteq \text{EXPSPACE}$   
(at least one relationship strict)
- ▶  $P \neq \text{EXPTIME}$ ,  $\text{PSPACE} \neq \text{EXPSPACE}$
- ▶  $P \subseteq NP \subseteq \text{PSPACE}$

## D6.3 Polynomial Hierarchy

## Oracle Machines

An **oracle machine** is like a Turing machine that has access to an **oracle** which can solve some decision problem in constant time.

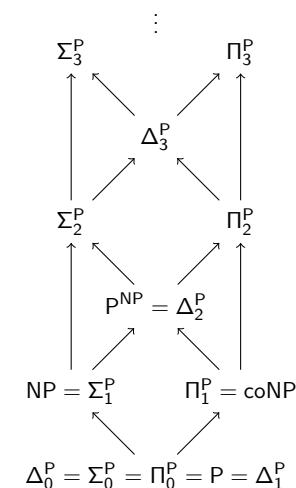
Example oracle classes:

- ▶  $P^{NP} = \{L \mid L \text{ can get decided in polynomial time by a DTM with an oracle that decides some problem in NP}\}$
- ▶  $NP^{NP} = \{L \mid L \text{ can get decided in pol. time by a NTM with an oracle deciding some problem in NP}\}$

## Polynomial Hierarchy

Inductively defined:

- ▶  $\Delta_0^P := \Sigma_0^P := \Pi_0^P := P$
- ▶  $\Delta_{i+1}^P := P^{\Sigma_i^P}$
- ▶  $\Sigma_{i+1}^P := NP^{\Sigma_i^P}$
- ▶  $\Pi_{i+1}^P := coNP^{\Sigma_i^P}$
- ▶  $PH := \bigcup_k \Sigma_k^P$



## Polynomial Hierarchy: Results

- ▶  $PH \subseteq PSPACE$  ( $PH \stackrel{?}{=} PSPACE$  is open)
- ▶ There are complete problems for each level.
- ▶ If there is a PH-complete problem, then the polynomial hierarchy collapses to some finite level.
- ▶ If  $P = NP$ , the polynomial hierarchy collapses to the first level.

## D6.4 Counting

# #P

## Complexity class #P

- Set of functions  $f : \{0, 1\}^* \rightarrow \mathbb{N}_0$ , where  $f(n)$  is the number of accepting paths of a polynomial-time NTM

### Example (#SAT)

The following problem **#SAT** is #P-complete:

**Given:** a propositional logic formula  $\varphi$

**Question:** Under how many variable assignments is  $\varphi$  true?

## What's Next?

contents of this course:

- A. **background** ✓
  - ▷ mathematical foundations and proof techniques
- B. **automata theory and formal languages** ✓
  - ▷ What is a computation?
- C. **Turing computability** ✓
  - ▷ What can be computed at all?
- D. **complexity theory** ✓
  - ▷ What can be computed efficiently?
- E. ~~more computability theory~~
  - ▷ Other models of computability