

Theory of Computer Science

D4. Halting Problem Variants & Rice's Theorem

Gabriele Röger

University of Basel

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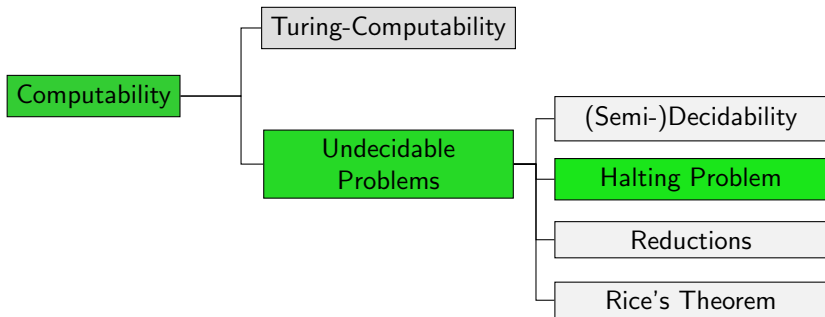
April 29, 2020 — D4. Halting Problem Variants & Rice's Theorem

D4.1 Other Halting Problem Variants

D4.2 Rice's Theorem

D4.1 Other Halting Problem Variants

Overview: Computability Theory



Reminder: Special Halting Problem

Definition (Special Halting Problem)

The **special halting problem** or **self-application problem** is the language

$$K = \{w \in \{0, 1\}^* \mid M_w \text{ started on } w \text{ terminates}\}.$$

German: spezielles Halteproblem, Selbstanwendbarkeitsproblem

General Halting Problem (1)

Definition (General Halting Problem)

The **general halting problem** or **halting problem** is the language

$$H = \{w\#x \in \{0, 1, \#\}^* \mid w, x \in \{0, 1\}^*, \\ M_w \text{ started on } x \text{ terminates}\}$$

German: allgemeines Halteproblem, Halteproblem

Note: H is semi-decidable. (Why?)

Theorem (Undecidability of General Halting Problem)

The general halting problem is undecidable.

Intuition: if the special case K is not decidable, then the more general problem H definitely cannot be decidable.

General Halting Problem (2)

Proof.

We show $K \leq H$.

We define $f : \{0, 1\}^* \rightarrow \{0, 1, \#\}^*$ as $f(w) := w\#w$.

f is clearly total and computable, and

$$w \in K$$

iff M_w started on w terminates

iff $w\#w \in H$

iff $f(w) \in H$.

Therefore f is a reduction from K to H .

Because K is undecidable, H is also undecidable. □

Halting Problem on Empty Tape (1)

Definition (Halting Problem on the Empty Tape)

The **halting problem on the empty tape** is the language

$$H_0 = \{w \in \{0, 1\}^* \mid M_w \text{ started on } \varepsilon \text{ terminates}\}.$$

German: Halteproblem auf leerem Band

Note: H_0 is semi-decidable. (Why?)

Theorem (Undecidability of Halting Problem on Empty Tape)

The halting problem on the empty tape is undecidable.

Halting Problem on Empty Tape (2)

Proof.

We show $H \leq H_0$.

Consider the function $f : \{0, 1, \#\}^* \rightarrow \{0, 1\}^*$

that computes the word $f(z)$ for a given $z \in \{0, 1, \#\}^*$ as follows:

- ▶ Test if z has the form $w\#x$ with $w, x \in \{0, 1\}^*$.
- ▶ If not, return any word that is not in H_0
(e. g., encoding of a TM that instantly starts an endless loop).
- ▶ If yes, split z into w and x .
- ▶ Decode w to a TM M_2 .

...

Halting Problem on Empty Tape (3)

Proof (continued).

- ▶ Construct a TM M_1 that behaves as follows:
 - ▶ If the input is empty: write x onto the tape and move the head to the first symbol of x (if $x \neq \varepsilon$); then stop
 - ▶ otherwise, stop immediately
- ▶ Construct TM M that first runs M_1 and then M_2 .
- ▶ Return the encoding of M .

f is total and (with some effort) computable. Also:

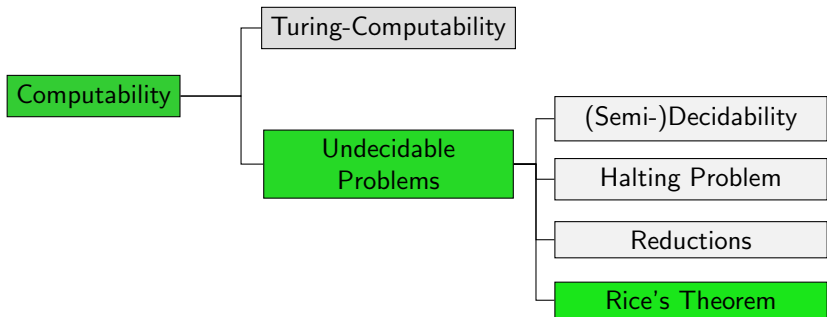
$$\begin{aligned}
 z \in H &\text{ iff } z = w\#x \text{ and } M_w \text{ run on } x \text{ terminates} \\
 &\text{ iff } M_{f(z)} \text{ started on empty tape terminates} \\
 &\text{ iff } f(z) \in H_0
 \end{aligned}$$

$\rightsquigarrow H \leq H_0 \rightsquigarrow H_0$ undecidable



D4.2 Rice's Theorem

Overview: Computability Theory



Rice's Theorem (1)

- ▶ We have shown that a number of (related) problems are undecidable:
 - ▶ special halting problem K
 - ▶ general halting problem H
 - ▶ halting problem on empty tape H_0
- ▶ Many more results of this type could be shown.
- ▶ Instead, we prove a much more general result, **Rice's theorem**, which shows that a very large class of different problems are undecidable.
- ▶ Rice's theorem can be summarized informally as: **every** non-trivial question about **what** a given Turing machine computes is undecidable.

Rice's Theorem (2)

Theorem (Rice's Theorem)

Let \mathcal{R} be the class of all computable functions.

Let \mathcal{S} be an *arbitrary* subset of \mathcal{R} except $\mathcal{S} = \emptyset$ or $\mathcal{S} = \mathcal{R}$.

Then the language

$$C(\mathcal{S}) = \{w \in \{0, 1\}^* \mid \text{the function computed by } M_w \text{ is in } \mathcal{S}\}$$

is undecidable.

German: Satz von Rice

Question: why the restriction to $\mathcal{S} \neq \emptyset$ and $\mathcal{S} \neq \mathcal{R}$?

Extension (without proof): in most cases neither $C(\mathcal{S})$ nor $\overline{C(\mathcal{S})}$ is semi-decidable. (But there are sets \mathcal{S} for which one of the two languages is semi-decidable.)

Rice's Theorem (3)

Proof.

Let Ω be the function that is undefined everywhere.

Case distinction:

Case 1: $\Omega \in \mathcal{S}$

Let $q \in \mathcal{R} \setminus \mathcal{S}$ be an arbitrary computable function outside of \mathcal{S} (exists because $\mathcal{S} \subseteq \mathcal{R}$ and $\mathcal{S} \neq \mathcal{R}$).

Let Q be a Turing machine that computes q .

...

Rice's Theorem (4)

Proof (continued).

We show that $\bar{H}_0 \leq C(S)$.

Consider function $f : \{0, 1\}^* \rightarrow \{0, 1\}^*$,
where $f(w)$ is defined as follows:

- ▶ Construct TM M that first behaves on input y like M_w on the empty tape (independently of what y is).
- ▶ Afterwards (if that computation terminates!) M clears the tape, creates the start configuration of Q for input y and then simulates Q .
- ▶ $f(w)$ is the encoding of this TM M

f is total and computable.

...

Rice's Theorem (5)

Proof (continued).

Which function is computed by the TM encoded by $f(w)$?

$M_{f(w)}$ computes $\begin{cases} \Omega & \text{if } M_w \text{ does not terminate on } \varepsilon \\ q & \text{otherwise} \end{cases}$

For all words $w \in \{0, 1\}^*$:

$w \in H_0 \implies M_w$ terminates on ε

$\implies M_{f(w)}$ computes the function q

\implies the function computed by $M_{f(w)}$ is not in \mathcal{S}

$\implies f(w) \notin C(\mathcal{S})$

...

Rice's Theorem (6)

Proof (continued).

Further:

$$\begin{aligned}w \notin H_0 &\implies M_w \text{ does not terminate on } \varepsilon \\ &\implies M_{f(w)} \text{ computes the function } \Omega \\ &\implies \text{the function computed by } M_{f(w)} \text{ is in } \mathcal{S} \\ &\implies f(w) \in C(\mathcal{S})\end{aligned}$$

Together this means: $w \notin H_0$ iff $f(w) \in C(\mathcal{S})$,
thus $w \in \bar{H}_0$ iff $f(w) \in C(\mathcal{S})$.

Therefore, f is a reduction of \bar{H}_0 to $C(\mathcal{S})$.

Since H_0 is undecidable, \bar{H}_0 is also undecidable.

We can conclude that $C(\mathcal{S})$ is undecidable.

...

Rice's Theorem (7)

Proof (continued).

Case 2: $\Omega \notin \mathcal{S}$

Analogous to Case 1 but this time choose $q \in \mathcal{S}$.

The corresponding function f then reduces H_0 to $C(\mathcal{S})$.

Thus, it also follows in this case that $C(\mathcal{S})$ is undecidable. □

Rice's Theorem: Consequences

Was it worth it?

We can now conclude immediately that (for example) the following informally specified problems are all undecidable:

- ▶ Does a given TM compute a constant function?
- ▶ Does a given TM compute a total function (i. e. will it always terminate, and in particular terminate in a “correct” configuration)?
- ▶ Is the output of a given TM always longer than its input?
- ▶ Does a given TM compute the identity function?
- ▶ Does a given TM compute the computable function f ?
- ▶ ...

Rice's Theorem: Examples

- ▶ Does a given TM compute a constant function?
 $\mathcal{S} = \{f \mid f \text{ is total and computable and}$
for all x, y in the domain of $f : f(x) = f(y)\}$
- ▶ Does a given TM compute a total function?
 $\mathcal{S} = \{f \mid f \text{ is total and computable}\}$
- ▶ Does a given TM compute the identity function?
 $\mathcal{S} = \{f \mid f(x) = x \text{ for all } x\}$
- ▶ Does a given TM add two natural numbers?
 $\mathcal{S} = \{f : \mathbb{N}_0^2 \rightarrow \mathbb{N}_0 \mid f(x, y) = x + y\}$
- ▶ Does a given TM compute the computable function f ?
 $\mathcal{S} = \{f\}$
(full automatization of software verification is impossible)

Rice's Theorem: Pitfalls

- ▶ $\mathcal{S} = \{f \mid f \text{ can be computed by a DTM with an even number of states}\}$
 Rice's theorem not applicable because $\mathcal{S} = \mathcal{R}$
- ▶ $\mathcal{S} = \{f : \{0, 1\}^* \rightarrow_p \{0, 1\} \mid f(w) = 1 \text{ iff } M_w \text{ does not terminate on } \epsilon\}$?
 Rice's theorem not applicable because $\mathcal{S} \not\subseteq \mathcal{R}$
- ▶ Show that $\{w \mid M_w \text{ traverses all states on every input}\}$ is undecidable.
 Rice's theorem not directly applicable because not a semantic property (the function computed by M_w can also be computed by a TM that does not traverse all states)

Rice's Theorem: Practical Applications

Undecidable due to Rice's theorem + a small reduction:

- ▶ **automated debugging:**
 - ▶ Can a given variable ever receive a `null` value?
 - ▶ Can a given assertion in a program ever trigger?
 - ▶ Can a given buffer ever overflow?
- ▶ **virus scanners and other software security analysis:**
 - ▶ Can this code do something harmful?
 - ▶ Is this program vulnerable to SQL injections?
 - ▶ Can this program lead to a privilege escalation?
- ▶ **optimizing compilers:**
 - ▶ Is this dead code?
 - ▶ Is this a constant expression?
 - ▶ Can pointer aliasing happen here?
 - ▶ Is it safe to parallelize this code path?
- ▶ **parallel program analysis:**
 - ▶ Is a deadlock possible here?
 - ▶ Can a race condition happen here?

Summary

undecidable but semi-decidable problems:

- ▶ **special halting problem** a.k.a. self-application problem (from previous chapter)
- ▶ **general halting problem**
- ▶ **halting problem on empty tape**

Rice's theorem:

- ▶ “In general one cannot determine algorithmically what a given program (or Turing machine) computes.”