

Theory of Computer Science

D1. Turing-Computability

Gabriele Röger

University of Basel

April 20, 2020

Theory of Computer Science

April 20, 2020 — D1. Turing-Computability

D1.1 Turing-Computable Functions

D1.2 Summary

Overview: Course

contents of this course:

A. background ✓

▷ mathematical foundations and proof techniques

B. logic ✓

▷ How can knowledge be represented?
How can reasoning be automated?

C. automata theory and formal languages ✓

▷ What is a computation?

D. Turing computability

▷ What can be computed at all?

E. complexity theory

▷ What can be computed efficiently?

F. more computability theory

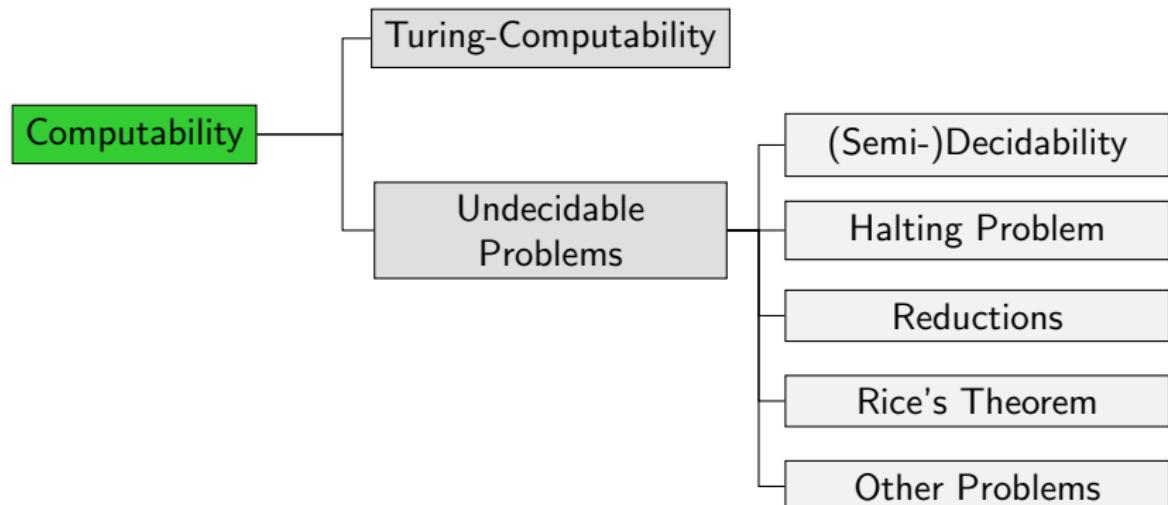
▷ Other models of computability

Main Question

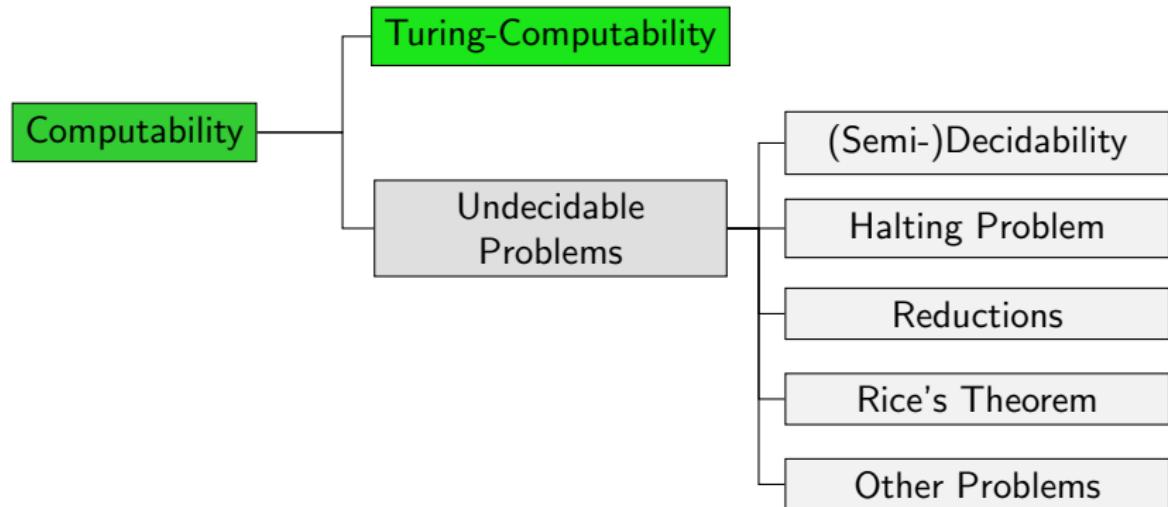
Main question in this part of the course:

What can be computed
by a computer?

Overview: Computability Theory



Overview: Computability Theory



D1.1 Turing-Computable Functions

Computation

What is a computation?

- ▶ **intuitive model of computation** (pen and paper)
- ▶ vs. **computation on physical computers**
- ▶ vs. **formal mathematical models**

In the following chapters we investigate
models of computation for **partial functions** $f : \mathbb{N}_0^k \rightarrow_p \mathbb{N}_0$.

- ▶ no real limitation: arbitrary information
can be encoded as numbers

German: Berechnungsmodelle

Church-Turing Thesis

Church-Turing Thesis

All functions that can be **computed in the intuitive sense**
can be computed by a **Turing machine**.

German: Church-Turing-These

- ▶ cannot be proven ([why not?](#))
- ▶ but we will collect evidence for it (\rightsquigarrow [part F](#))

Reminder: Deterministic Turing Machine (DTM)

Definition (Deterministic Turing Machine)

A **deterministic Turing machine (DTM)** is given by a 7-tuple

$M = \langle Q, \Sigma, \Gamma, \delta, q_0, \square, E \rangle$ with:

- ▶ Q finite, non-empty set of **states**
- ▶ $\Sigma \neq \emptyset$ finite **input alphabet**
- ▶ $\Gamma \supset \Sigma$ finite **tape alphabet**
- ▶ $\delta : (Q \setminus E) \times \Gamma \rightarrow Q \times \Gamma \times \{L, R, N\}$ **transition function**
- ▶ $q_0 \in Q$ **start state**
- ▶ $\square \in \Gamma \setminus \Sigma$ **blank symbol**
- ▶ $E \subseteq Q$ **end states**

Computation of Functions?

How can a DTM compute a function?

- ▶ “Input” x is the initial tape content
- ▶ “Output” $f(x)$ is the tape content (ignoring blanks at the left and right) when reaching an end state
- ▶ If the TM does not stop for the given input, $f(x)$ is undefined for this input.

Which kinds of functions can be computed this way?

- ▶ directly, only functions on **words**: $f : \Sigma^* \rightarrow_p \Sigma^*$
- ▶ interpretation as functions on **numbers** $f : \mathbb{N}_0^k \rightarrow_p \mathbb{N}_0$: encode numbers as words

Turing Machines: Computed Function

Definition (Function Computed by a Turing Machine)

A DTM $M = \langle Q, \Sigma, \Gamma, \delta, q_0, \square, E \rangle$ **computes** the (partial) function $f : \Sigma^* \rightarrow_p \Sigma^*$ for which:

for all $x, y \in \Sigma^*$: $f(x) = y$ iff $\langle \varepsilon, q_0, x \rangle \vdash^* \langle \square \dots \square, q_e, y \square \dots \square \rangle$

with $q_e \in E$. (special case: initial configuration $\langle \varepsilon, q_0, \square \rangle$ if $x = \varepsilon$)

German: DTM berechnet f

- ▶ What happens if symbols from $\Gamma \setminus \Sigma$ (e. g., \square) occur in y ?
- ▶ What happens if the read-write head is not on the first symbol of y at the end?
- ▶ Is f uniquely defined by this definition? Why?

Turing-Computable Functions on Words

Definition (Turing-Computable, $f : \Sigma^* \rightarrow_p \Sigma^*$)

A (partial) function $f : \Sigma^* \rightarrow_p \Sigma^*$ is called **Turing-computable** if a DTM that computes f exists.

German: Turing-berechenbar

Example: Turing-Computable Functions on Words

Example

Let $\Sigma = \{a, b, \#\}$.

The function $f : \Sigma^* \rightarrow_p \Sigma^*$ with $f(w) = w\#w$ for all $w \in \Sigma^*$ is Turing-computable.

\rightsquigarrow blackboard

Encoding Numbers as Words

Definition (Encoded Function)

Let $f : \mathbb{N}_0^k \rightarrow_p \mathbb{N}_0$ be a (partial) function.

The **encoded function** f^{code} of f is the partial function

$f^{\text{code}} : \Sigma^* \rightarrow_p \Sigma^*$ with $\Sigma = \{0, 1, \#\}$ and $f^{\text{code}}(w) = w'$ iff

- ▶ there are $n_1, \dots, n_k, n' \in \mathbb{N}_0$ such that
- ▶ $f(n_1, \dots, n_k) = n'$,
- ▶ $w = \text{bin}(n_1)\# \dots \#\text{bin}(n_k)$ and
- ▶ $w' = \text{bin}(n')$.

Here $\text{bin} : \mathbb{N}_0 \rightarrow \{0, 1\}^*$ is the binary encoding
(e.g., $\text{bin}(5) = 101$).

German: kodierte Funktion

Example: $f(5, 2, 3) = 4$ corresponds to $f^{\text{code}}(101\#10\#11) = 100$.

Turing-Computable Numerical Functions

Definition (Turing-Computable, $f : \mathbb{N}_0^k \rightarrow_p \mathbb{N}_0$)

A (partial) function $f : \mathbb{N}_0^k \rightarrow_p \mathbb{N}_0$ is called **Turing-computable** if a DTM that computes f^{code} exists.

German: Turing-berechenbar

Example: Turing-Computable Numerical Function

Example

The following numerical functions are Turing-computable:

- ▶ $\text{succ} : \mathbb{N}_0 \rightarrow_p \mathbb{N}_0$ with $\text{succ}(n) := n + 1$
- ▶ $\text{pred}_1 : \mathbb{N}_0 \rightarrow_p \mathbb{N}_0$ with $\text{pred}_1(n) := \begin{cases} n - 1 & \text{if } n \geq 1 \\ 0 & \text{if } n = 0 \end{cases}$
- ▶ $\text{pred}_2 : \mathbb{N}_0 \rightarrow_p \mathbb{N}_0$ with $\text{pred}_2(n) := \begin{cases} n - 1 & \text{if } n \geq 1 \\ \text{undefined} & \text{if } n = 0 \end{cases}$

↔ blackboard

More Turing-Computable Numerical Functions

Example

The following numerical functions are Turing-computable:

- ▶ $add : \mathbb{N}_0^2 \rightarrow_p \mathbb{N}_0$ with $add(n_1, n_2) := n_1 + n_2$
- ▶ $sub : \mathbb{N}_0^2 \rightarrow_p \mathbb{N}_0$ with $sub(n_1, n_2) := \max\{n_1 - n_2, 0\}$
- ▶ $mul : \mathbb{N}_0^2 \rightarrow_p \mathbb{N}_0$ with $mul(n_1, n_2) := n_1 \cdot n_2$
- ▶ $div : \mathbb{N}_0^2 \rightarrow_p \mathbb{N}_0$ with $div(n_1, n_2) := \begin{cases} \left\lceil \frac{n_1}{n_2} \right\rceil & \text{if } n_2 \neq 0 \\ \text{undefined} & \text{if } n_2 = 0 \end{cases}$

~~ sketch?

D1.2 Summary

Summary

- ▶ main question: what can a computer compute?
- ▶ approach: investigate formal models of computation
- ▶ here: deterministic Turing machines
- ▶ **Turing-computable** function $f : \Sigma^* \rightarrow_p \Sigma^*$:
there is a DTM that transforms every input $w \in \Sigma^*$
into the output $f(w)$ (undefined if DTM does not stop
or stops in invalid configuration)
- ▶ **Turing-computable** function $f : \mathbb{N}_0^k \rightarrow_p \mathbb{N}_0$:
ditto; numbers encoded in binary and separated by #