

# Theory of Computer Science

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## Exercise Sheet 11

Due: Wednesday, May 13, 2020

### Exercise 11.1 (Undecidability; 1+3+2 marks)

- Show that for every alphabet  $\Sigma$  there is a context-free grammar  $G_{\text{palindrome}}(\Sigma)$  that generates exactly the non-empty palindromes over  $\Sigma$ .
- Consider a PCP-instance  $\mathcal{I} = \langle (x_1, y_1), \dots, (x_k, y_k) \rangle$  with  $x_i, y_i \in \Sigma^+$ . Let  $\Sigma' = \Sigma \cup \{\#\}$  where  $\#$  is a symbol that does not occur in  $\Sigma$ . Specify a context-free grammar  $G_{\mathcal{I}}$  over  $\Sigma'$ , that generates the following language (with  $^{-1}$  we denote the reverse of a word here):

$$\mathcal{L}(G_{\mathcal{I}}) = \{x_{i_1} \dots x_{i_n} \# y_{i_n}^{-1} \dots y_{i_1}^{-1} \mid n \geq 1 \text{ and } i_1, \dots, i_n \in \{1, \dots, k\}\}$$

Moreover, show that  $\mathcal{I}$  has a solution iff  $\mathcal{L}(G_{\mathcal{I}})$  contains a palindrome.

- Use the results from parts (a) and (b) to prove that the intersection problem of *context-free* languages is undecidable.

INTERSECTION<sub>CF</sub> : Given two context-free grammars  $G_1$  and  $G_2$ , is  $\mathcal{L}(G_1) \cap \mathcal{L}(G_2) = \emptyset$ ?

*Note: of course you can use the statements of parts (a) and (b) even if you did not solve the exercises.*

### Exercise 11.2 (Non-deterministic Algorithms; 2+2 marks)

Specify a non-deterministic polynomial algorithm for each of the following problems. This shows that the problems are in NP, a complexity class we will introduce next week.

- HITTINGSET:

- Given:* finite set  $M$ , set of sets  $\mathcal{S} = \{S_1, \dots, S_n\}$  with  $S_i \subseteq M$  for all  $i \in \{1, \dots, n\}$ , natural number  $k \in \mathbb{N}_0$
- Question:* Is there a set  $H$  with at most  $k$  elements, which contains at least one element from each set in  $\mathcal{S}$ .

Formally: Is there a set  $H$  with  $|H| \leq k$  and  $H \cap S_i \neq \emptyset$  for all  $i \in \{1, \dots, n\}$ ?

- SETPACKING:

- Given:* finite set  $M$ , set of sets  $\mathcal{S} = \{S_1, \dots, S_n\}$  with  $S_i \subseteq M$  for all  $i \in \{1, \dots, n\}$ , natural number  $k \in \mathbb{N}_0$
- Question:* Is there a set  $\mathcal{S}' \subseteq \mathcal{S}$  with  $|\mathcal{S}'| \geq k$ , such that all sets in  $\mathcal{S}'$  are pairwise disjoint, i.e., for all  $S_i, S_j \in \mathcal{S}'$  with  $S_i \neq S_j$  it holds that  $S_i \cap S_j = \emptyset$ ?