

Theory of Computer Science

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Exercise Sheet 4

Due: Wednesday, March 25, 2020

Exercise 4.1 (Predicate logic, 2 + 2 marks)

(a) Show for arbitrary predicate logic formulas φ and ψ that

$$(\forall x\varphi \vee \forall x\psi) \models \forall x(\varphi \vee \psi).$$

(b) Show that in general it is *not* the case that

$$(\forall x\varphi \vee \forall x\psi) \equiv \forall x(\varphi \vee \psi)$$

Specify a counter example with the following signature: $\mathcal{S} = \langle \{x\}, \{\}, \{\}, \{P, Q\} \rangle$, where $ar(P) = ar(Q) = 1$.

Exercise 4.2 (Predicate logic, 1 mark)

Use equivalence transformations to bring the following formula in negation normal form. For this purpose, move negation symbols inwards by using DeMorgan's rule or equivalence $\neg\forall x\varphi \equiv \exists x\neg\varphi$ and $\neg\exists x\varphi \equiv \forall x\neg\varphi$, or eliminate them with double negation.

$$\varphi = \neg\forall x((P(x) \vee \neg Q(x, c)) \wedge \exists y(P(x) \rightarrow Q(y, x)))$$

Exercise 4.3 (Formal languages and grammars, 1+3+1 marks)

Consider the following formal language over $\{a, b, c\}$:

$$L = \{a^n b^m c^{2n} \mid n \geq 0, m \geq 0\}$$

(a) Is ε an element of L ? Justify your answer.

(b) Specify a *complete description* of a formal grammar G that generates L (i.e., $\mathcal{L}(G) = L$). A formal grammar is a four tuple $G = \langle \Sigma, V, P, S \rangle$, remember to define all components of this tuple.

(c) Which types (in the Chomsky-Hierarchy) is your formal grammar part of? You don't have to prove your answers.