Theory of Computer Science
E4. Some NP-Complete Problems, Part I

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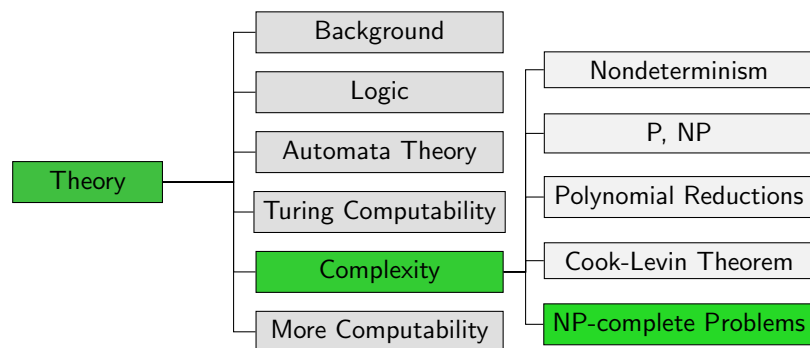

E4.1 Graph Problems

E4.2 Routing Problems

Course Overview

E4.1 Graph Problems

$3\text{SAT} \leq_p \text{CLIQUE}$

SAT

$\downarrow$

3SAT

CLIQUE   DIRHAMILTONCYCLE   SUBSETSUM

$\downarrow$ $\downarrow$ $\downarrow$

INDSET   HAMILTONCYCLE   PARTITION

$\downarrow$ $\downarrow$ $\downarrow$

VERTEXCOVER   TSP   BINPACKING
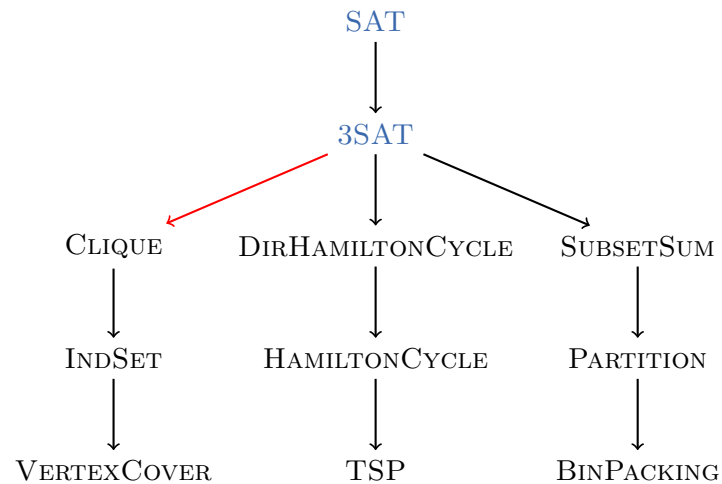
CLIQUE

Definition (CLIQUE)

The problem CLIQUE is defined as follows:

Given: undirected graph $G = \langle V, E \rangle$, number $K \in \mathbb{N}_0$

Question: Does G have a clique of size at least K,
i. e., a set of vertices $C \subseteq V$ with $|C| \geq K$
and $\{u, v\} \in E$ for all $u, v \in C$ with $u \neq v$?

German: Clique

CLIQUE is NP-Complete (1)

Theorem (CLIQUE is NP-Complete)

CLIQUE *is NP-complete.*

CLIQUE is NP-Complete (2)

Proof.

CLIQUE $\in$ NP: guess and check.

CLIQUE is NP-hard: We show $3\text{SAT} \leq_p \text{CLIQUE}$.

▶ We are given a 3-CNF formula φ, and we may assume
  that each clause has exactly three literals.

▶ In polynomial time, we must construct
  a graph $G = \langle V, E \rangle$ and a number K such that:
  G has a clique of size at least K iff φ is satisfiable.

⇝ construction of V, E, K on the following slides.

. . .

CLIQUE is NP-Complete (3)

Proof (continued).

Let m be the number of clauses in φ.

Let ℓ_{ij} the j-th literal in clause i.

Define V, E, K as follows:

- $V = \{\langle i,j\rangle \mid 1 \le i \le m, 1 \le j \le 3\}$
  $\rightsquigarrow$ a vertex for every literal of every clause
- E contains edge between $\langle i,j\rangle$ and $\langle i',j'\rangle$ if and only if
  - $i \ne i'$ $\rightsquigarrow$ belong to different clauses, and
  - ℓ_{ij} and $\ell_{i'j'}$ are not complementary literals
- $K = m$

$\rightsquigarrow$ obviously polynomially computable

to show: reduction property   . . .

CLIQUE is NP-Complete (4)

Proof (continued).

($\Rightarrow$): If φ is satisfiable, then $\langle V, E\rangle$ has clique of size at least K:

- Given a satisfying variable assignment choose a vertex corresponding to a satisfied literal in each clause.
- The chosen K vertices are all connected with each other and hence form a clique of size K.

. . .
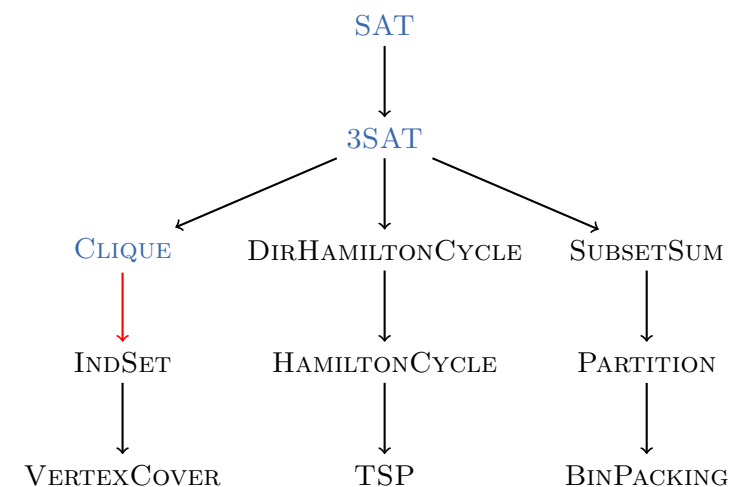
CLIQUE is NP-Complete (5)

Proof (continued).

($\Leftarrow$): If $\langle V, E\rangle$ has a clique of size at least K, then φ is satisfiable:

- Consider a given clique C of size at least K.
- The vertices in C must all correspond to different clauses (vertices in the same clause are not connected by edges).
- $\rightsquigarrow$ exactly one vertex per clause is included in C
- Two vertices in C never correspond to complementary literals X and $\neg X$ (due to the way we defined the edges).
- If a vertex corresp. to X was chosen, map X to 1 (true).
- If a vertex corresp. to $\neg X$ was chosen, map X to 0 (false).
- If neither was chosen, arbitrarily map X to 0 or 1.

$\rightsquigarrow$ satisfying assignment

$\square$

CLIQUE $\le_p$ INDSET

SAT

3SAT

CLIQUE   DIRHAMILTONCYCLE   SUBSETSUM

INDSET   HAMILTONCYCLE   PARTITION

VERTEXCOVER   TSP   BINPACKING

INDSET

> **Definition (INDSET)**
>
> The problem INDSET is defined as follows:
>
> Given: undirected graph $G = \langle V, E \rangle$, number $K \in \mathbb{N}_0$
>
> Question: Does G have an independent set of size at least K,
> i.e., a set of vertices $I \subseteq V$ with $|I| \geq K$
> and $\{u, v\} \notin E$ for all $u, v \in I$ with $u \neq v$?
>
> German: unabhängige Menge

INDSET is NP-Complete (1)

> **Theorem (INDSET is NP-Complete)**
>
> INDSET *is NP-complete.*

INDSET is NP-Complete (2)

> **Proof.**
>
> INDSET $\in$ NP: guess and check.
>
> INDSET is NP-hard: We show CLIQUE $\leq_p$ INDSET.
>
> We describe a polynomial reduction f.
> Let $\langle G, K \rangle$ with $G = \langle V, E \rangle$ be the given input for CLIQUE.
> Then $f(\langle G, K \rangle)$ is the INDSET instance $\langle \overline{G}, K \rangle$, where
> $\overline{G} := \langle V, \overline{E} \rangle$ and $\overline{E} := \{\{u, v\} \subseteq V \mid u \neq v, \{u, v\} \notin E\}$.
> (This graph $\overline{G}$ is called the complement graph of G.)
> Clearly f can be computed in polynomial time.   . . .
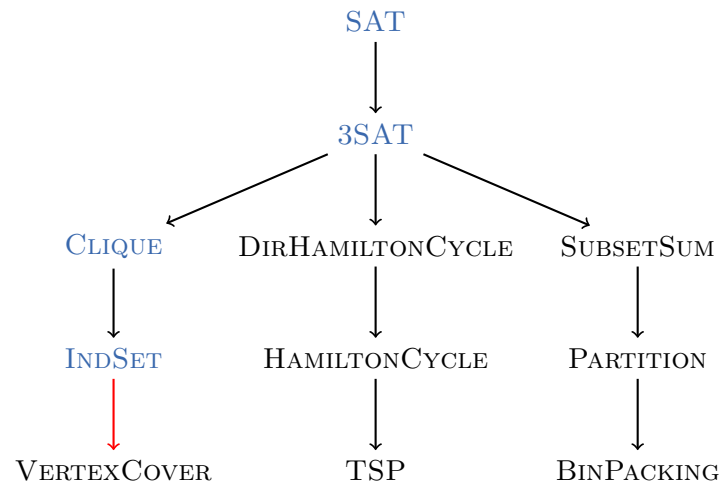
INDSET is NP-Complete (3)

> **Proof (continued).**
>
> We have:
>
> $$\langle \langle V, E \rangle, K \rangle \in \text{CLIQUE}$$
> iff   there exists a set $V' \subseteq V$ with $|V'| \geq K$
>   and $\{u, v\} \in E$ for all $u, v \in V'$ with $u \neq v$
> iff   there exists a set $V' \subseteq V$ with $|V'| \geq K$
>   and $\{u, v\} \notin \overline{E}$ for all $u, v \in V'$ with $u \neq v$
> iff $\langle \langle V, \overline{E} \rangle, K \rangle \in \text{INDSET}$
> iff $f(\langle \langle V, E \rangle, K \rangle) \in \text{INDSET}$
>
> and hence f is a reduction.   □

INDSET $\leq_p$ VERTEXCOVER

SAT

$\downarrow$

3SAT

CLIQUE   DIRHAMILTONCYCLE   SUBSETSUM

$\downarrow$ $\downarrow$ $\downarrow$

INDSET   HAMILTONCYCLE   PARTITION

$\downarrow$ $\downarrow$ $\downarrow$

VERTEXCOVER   TSP   BINPACKING

VERTEXCOVER

Definition (VERTEXCOVER)

The problem VERTEXCOVER is defined as follows:

Given: undirected graph $G = \langle V, E \rangle$, number $K \in \mathbb{N}_0$

Question: Does G have a vertex cover of size at most K,
i.e., a set of vertices $C \subseteq V$ with $|C| \leq K$ and $\{u, v\} \cap C \neq \emptyset$
for all $\{u, v\} \in E$?

German: Knotenüberdeckung

VERTEXCOVER is NP-Complete (1)

Theorem (VERTEXCOVER is NP-Complete)

VERTEXCOVER *is NP-complete.*

VERTEXCOVER is NP-Complete (2)

Proof.

VERTEXCOVER $\in$ NP: guess and check.

VERTEXCOVER is NP-hard:

We show INDSET $\leq_p$ VERTEXCOVER.

We describe a polynomial reduction f.
Let $\langle G, K \rangle$ with $G = \langle V, E \rangle$ be the given input for INDSET.
Then $f(\langle G, K \rangle) := \langle G, |V| - K \rangle$.
This can clearly be computed in polynomial time.   . . .

VERTEXCOVER is NP-Complete (3)

Proof (continued).

For vertex set $V' \subseteq V$, we write $\overline{V'}$ for its complement $V \setminus V'$.

Observation: a set of vertices is a vertex cover
iff its complement is an independent set.

We thus have:
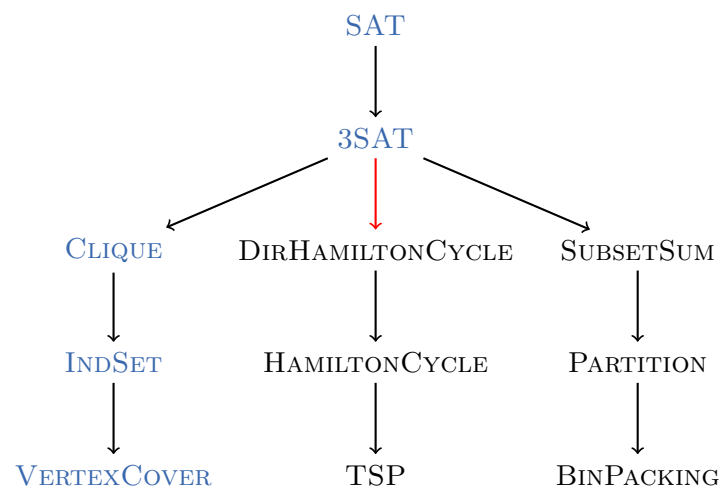
$$\langle\langle V, E\rangle, K\rangle \in \text{INDSET}$$
iff $\langle V, E\rangle$ has an independent set I with $|I| \geq K$
iff $\langle V, E\rangle$ has a vertex cover C with $|\overline{C}| \geq K$
iff $\langle V, E\rangle$ has a vertex cover C with $|C| \leq |V| - K$
iff $\langle\langle V, E\rangle, |V| - K\rangle \in \text{VERTEXCOVER}$
iff $f(\langle\langle V, E\rangle, K\rangle) \in \text{VERTEXCOVER}$

and hence f is a reduction.   □

E4.2 Routing Problems

3SAT $\leq_p$ DIRHAMILTONCYCLE

SAT

3SAT

CLIQUE   DIRHAMILTONCYCLE   SUBSETSUM

INDSET   HAMILTONCYCLE   PARTITION

VERTEXCOVER   TSP   BINPACKING

DIRHAMILTONCYCLE is NP-Complete (1)

Definition (Reminder: DIRHAMILTONCYCLE)
The problem DIRHAMILTONCYCLE is defined as follows:

Given: directed graph $G = \langle V, E\rangle$
Question: Does G contain a Hamilton cycle?

Theorem
DIRHAMILTONCYCLE is NP-complete.

DIRHAMILTONCYCLE is NP-Complete (2)

Proof.

$\text{DIRHAMILTONCYCLE} \in \text{NP}$: guess and check.

DIRHAMILTONCYCLE is NP-hard:

We show $3\text{SAT} \leq_{\mathsf{p}} \text{DIRHAMILTONCYCLE}$.

- ▶ We are given a 3-CNF formula φ where each clause contains exactly three literals and no clause contains duplicated literals.
- ▶ We must, in polynomial time, construct a directed graph $G = \langle V, E \rangle$ such that: G contains a Hamilton cycle iff φ is satisfiable.
- ▶ construction of $\langle V, E \rangle$ on the following slides

. . .

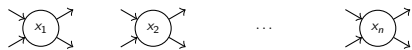## DIRHAMILTONCYCLE is NP-Complete (3)

Proof (continued).

- ▶ Let $X_1, \ldots, X_n$ be the propositional variables in φ.
- ▶ Let $c_1, \ldots, c_m$ be the clauses of φ with $c_i = (\ell_{i1} \vee \ell_{i2} \vee \ell_{i3})$.
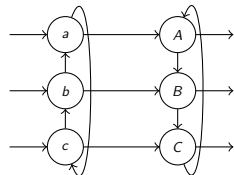- ▶ Construct a graph with $6m + n$ vertices (described on the following slides).

. . .

DIRHAMILTONCYCLE is NP-Complete (4)

Proof (continued).

- ▶ For every variable X_i, add vertex x_i with 2 incoming and 2 outgoing edges:

- ▶ For every clause c_j, add the subgraph C_j with 6 vertices:

- ▶ We describe later how to connect these parts.

. . .

DIRHAMILTONCYCLE is NP-Complete (5)

Proof (continued).

Let π be a Hamilton cycle of the total graph.

- ▶ Whenever π enters subgraph C_j from one of its "entrances", it must leave via the corresponding "exit": ($a \longrightarrow A$, $b \longrightarrow B$, $c \longrightarrow C$). Otherwise, π cannot be a Hamilton cycle.
- ▶ Hamilton cycles can behave in the following ways with regard to C_j:
  - ▶ π passes through C_j once (from any entrance)
  - ▶ π passes through C_j twice (from any two entrances)
  - ▶ π passes through C_j three times (once from every entrance)

. . .

DIRHAMILTONCYCLE is NP-Complete (6)

Proof (continued).

Connect the "open ends" in the graph as follows:

- Identify entrances/exits of the clause subgraph C_j
  with the three literals in clause c_j.
- One exit of x_i is positive, the other one is negative.
- For the positive exit, determine the clauses
  in which the positive literal X_i occurs:
  - Connect the positive exit of x_i with the X_i-entrance
    of the first such clause graph.
  - Connect the X_i-exit of this clause graph with the X_i-entrance
    of the second such clause graph, and so on.
  - Connect the X_i-exit of the last such clause graph
    with the positive entrance of x_{i+1} (or x_1 if $i = n$).
- analogously for the negative exit of x_i and the literal $\neg X_i$

. . .

DIRHAMILTONCYCLE is NP-Complete (7)

Proof (continued).

The construction is polynomial and is a reduction:

($\Rightarrow$): construct a Hamilton cycle from a satisfying assignment

- Given a satisfying assignment $\mathcal{I}$, construct a Hamilton cycle
  that leaves x_i through the positive exit if $\mathcal{I}(X_i)$ is true
  and by the negative exit if $\mathcal{I}(X_i)$ is false.
- Afterwards, we visit all C_j-subgraphs for clauses
  that are satisfied by this literal.
- In total, we visit each C_j-subgraph 1–3 times.

. . .

DIRHAMILTONCYCLE is NP-Complete (8)

Proof (continued).

($\Leftarrow$): construct a satisfying assignment from a Hamilton cycle

- A Hamilton cycle visits every vertex x_i
  and leaves it by the positive or negative exit.
- Map X_i to true or false depending on which exit
  is used to leave x_i.
- Because the cycle must traverse each C_j-subgraph
  at least once (otherwise it is not a Hamilton cycle),
  this results in a satisfying assignment. (Details omitted.)

$\square$

DIRHAMILTONCYCLE $\leq_p$ HAMILTONCYCLE
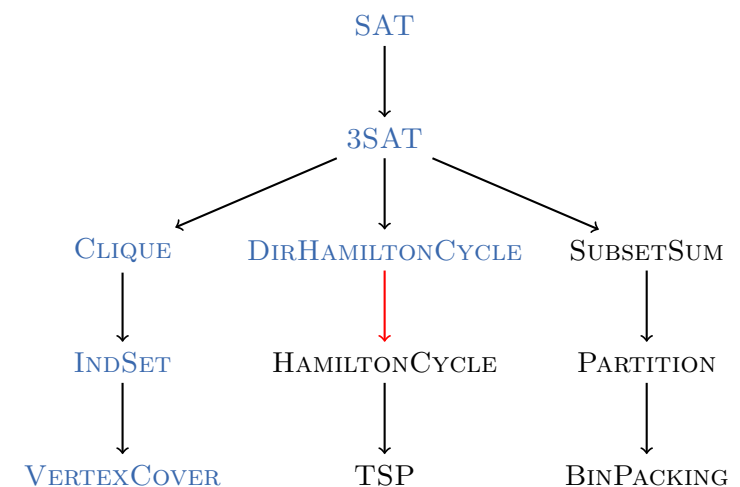
HAMILTONCYCLE is NP-Complete (1)

Definition (Reminder: HAMILTONCYCLE)

The problem HAMILTONCYCLE is defined as follows:

Given: undirected graph $G = \langle V, E \rangle$

Question: Does G contain a Hamilton cycle?

Theorem

HAMILTONCYCLE *is NP-complete.*

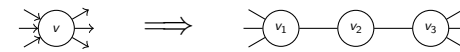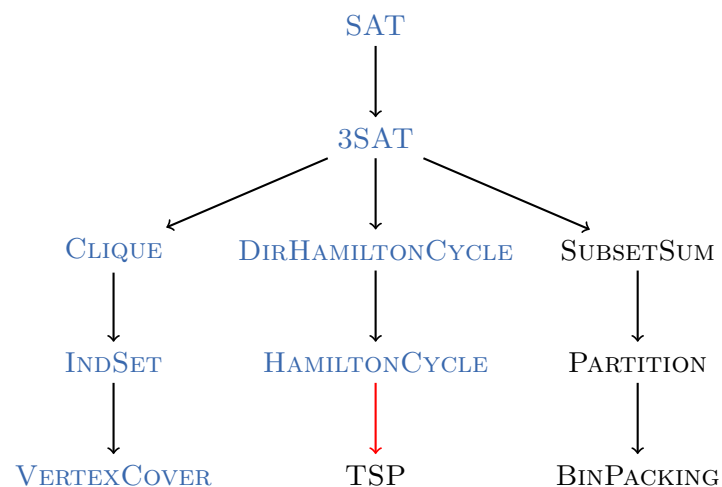# HAMILTONCYCLE is NP-Complete (2)

Proof sketch.

HAMILTONCYCLE $\in$ NP: guess and check.

HAMILTONCYCLE is NP-hard: We show
DIRHAMILTONCYCLE $\leq_p$ HAMILTONCYCLE.

Basic building block of the reduction:

HAMILTONCYCLE $\leq_p$ TSP

TSP is NP-Complete (1)

Definition (Reminder: TSP)

TSP (traveling salesperson problem) is the following
decision problem:

▶ Given: finite set $S \neq \emptyset$ of cities, symmetric cost function
$cost : S \times S \to \mathbb{N}_0$, cost bound $K \in \mathbb{N}_0$

▶ Question: Is there a tour with total cost at most K, i. e.,
a permutation $\langle s_1, \ldots, s_n \rangle$ of the cities with
$\sum_{i=1}^{n-1} cost(s_i, s_{i+1}) + cost(s_n, s_1) \leq K$?

German: Problem der/des Handlungsreisenden

Theorem

TSP *is NP-complete.*

TSP is NP-Complete (2)

Proof.

$\text{TSP} \in \text{NP}$: guess and check.

TSP is NP-hard: We showed HAMILTONCYCLE $\leq_p$ TSP
in Chapter E2. ☐

Summary

► In this chapter we showed NP-completeness of
  ► three classical graph problems:
    CLIQUE, INDSET, VERTEXCOVER
  ► three classical routing problems:
    DIRHAMILTONCYCLE, HAMILTONCYCLE, TSP