

# Theory of Computer Science

## D1. Turing-Computability

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# Overview: Course

contents of this course:

A. background ✓

- ▷ mathematical foundations and proof techniques

B. logic ✓

- ▷ How can knowledge be represented?  
How can reasoning be automated?

C. automata theory and formal languages ✓

- ▷ What is a computation?

D. Turing computability

- ▷ What can be computed at all?

E. complexity theory

- ▷ What can be computed efficiently?

F. more computability theory

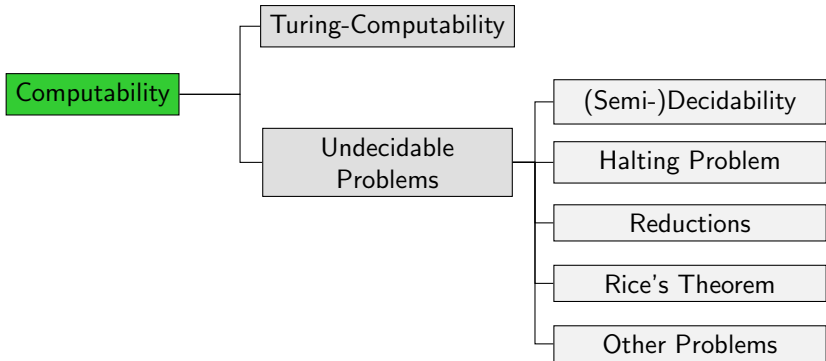
- ▷ Other models of computability

# Main Question

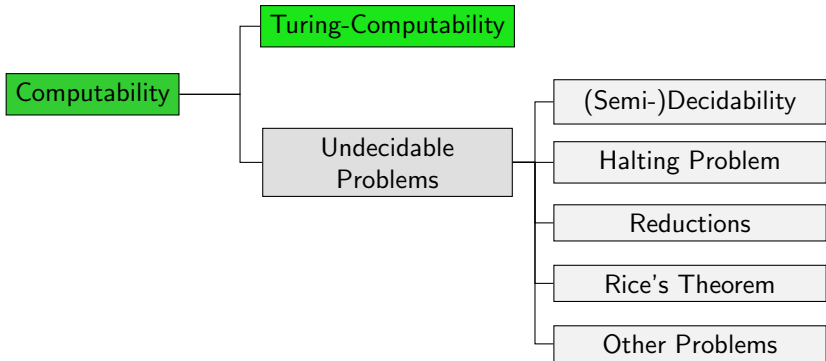
Main question in this part of the course:

**What can be computed  
by a computer?**

# Overview: Computability Theory



# Overview: Computability Theory



# Turing-Computable Functions

# Computation

What is a computation?

- intuitive model of computation (pen and paper)
- vs. computation on physical computers
- vs. formal mathematical models

In the following chapters we investigate

models of computation for partial functions  $f : \mathbb{N}_0^k \rightarrow_p \mathbb{N}_0$ .

- no real limitation: arbitrary information can be encoded as numbers

German: Berechnungsmodelle

# Church-Turing Thesis

## Church-Turing Thesis

All functions that can be **computed in the intuitive sense** can be computed by a **Turing machine**.

**German:** Church-Turing-These

- cannot be proven (**why not?**)
- but we will collect evidence for it (**↔ part F**)



# Reminder: Deterministic Turing Machine (DTM)

## Definition (Deterministic Turing Machine)

A **deterministic Turing machine (DTM)** is given by a 7-tuple  $M = \langle Q, \Sigma, \Gamma, \delta, q_0, \square, E \rangle$  with:

- $Q$  finite, non-empty set of **states**
- $\Sigma \neq \emptyset$  finite **input alphabet**
- $\Gamma \supset \Sigma$  finite **tape alphabet**
- $\delta : (Q \setminus E) \times \Gamma \rightarrow Q \times \Gamma \times \{L, R, N\}$  **transition function**
- $q_0 \in Q$  **start state**
- $\square \in \Gamma \setminus \Sigma$  **blank symbol**
- $E \subseteq Q$  **end states**

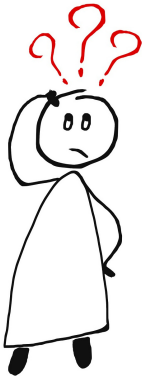
# Reminder: Configurations and Computation Steps

## How do Turing Machines Work?

- **configuration:**  $\langle \alpha, q, \beta \rangle$  with  $\alpha \in \Gamma^*$ ,  $q \in Q$ ,  $\beta \in \Gamma^+$
- **one computation step:**  $c \vdash c'$  if one computation step can turn configuration  $c$  into configuration  $c'$
- **multiple computation steps:**  $c \vdash^* c'$  if 0 or more computation steps can turn configuration  $c$  into configuration  $c'$   
( $c = c_0 \vdash c_1 \vdash c_2 \vdash \dots \vdash c_{n-1} \vdash c_n = c'$ ,  $n \geq 0$ )

(Definition of  $\vdash$ , i.e., how a computation step changes the configuration, is not repeated here.  $\rightsquigarrow$  [Chapter C7](#))

# Questions



Questions?

# Computation of Functions?

## How can a DTM compute a function?

- “Input”  $x$  is the initial tape content
- “Output”  $f(x)$  is the tape content (ignoring blanks at the left and right) when reaching an end state
- If the TM does not stop for the given input,  $f(x)$  is undefined for this input.

## Which kinds of functions can be computed this way?

- directly, only functions on **words**:  $f : \Sigma^* \rightarrow_p \Sigma^*$
- interpretation as functions on **numbers**  $f : \mathbb{N}_0^k \rightarrow_p \mathbb{N}_0$ :  
encode numbers as words

# Turing Machines: Computed Function

## Definition (Function Computed by a Turing Machine)

A DTM  $M = \langle Q, \Sigma, \Gamma, \delta, q_0, \square, E \rangle$  **computes** the (partial) function  $f : \Sigma^* \rightarrow_p \Sigma^*$  for which:

for all  $x, y \in \Sigma^*$ :  $f(x) = y$  iff  $\langle \varepsilon, q_0, x \rangle \vdash^* \langle \square \dots \square, q_e, y \square \dots \square \rangle$

with  $q_e \in E$ . (special case: initial configuration  $\langle \varepsilon, q_0, \square \rangle$  if  $x = \varepsilon$ )

**German:** DTM berechnet  $f$

- What happens if symbols from  $\Gamma \setminus \Sigma$  (e. g.,  $\square$ ) occur in  $y$ ?
- What happens if the read-write head is not on the first symbol of  $y$  at the end?
- Is  $f$  uniquely defined by this definition? Why?

# Turing-Computable Functions on Words

Definition (Turing-Computable,  $f : \Sigma^* \rightarrow_p \Sigma^*$ )

A (partial) function  $f : \Sigma^* \rightarrow_p \Sigma^*$  is called **Turing-computable** if a DTM that computes  $f$  exists.

German: Turing-berechenbar

## Example: Turing-Computable Functions on Words

### Example

Let  $\Sigma = \{a, b, \#\}$ .

The function  $f : \Sigma^* \rightarrow_p \Sigma^*$  with  $f(w) = w\#w$  for all  $w \in \Sigma^*$  is Turing-computable.

↪ blackboard

# Encoding Numbers as Words

## Definition (Encoded Function)

Let  $f : \mathbb{N}_0^k \rightarrow_p \mathbb{N}_0$  be a (partial) function.

The **encoded function**  $f^{\text{code}}$  of  $f$  is the partial function  $f^{\text{code}} : \Sigma^* \rightarrow_p \Sigma^*$  with  $\Sigma = \{0, 1, \#\}$  and  $f^{\text{code}}(w) = w'$  iff

- there are  $n_1, \dots, n_k, n' \in \mathbb{N}_0$  such that
- $f(n_1, \dots, n_k) = n'$ ,
- $w = \text{bin}(n_1)\#\dots\#\text{bin}(n_k)$  and
- $w' = \text{bin}(n')$ .

Here  $\text{bin} : \mathbb{N}_0 \rightarrow \{0, 1\}^*$  is the binary encoding (e. g.,  $\text{bin}(5) = 101$ ).

**German:** kodierte Funktion

**Example:**  $f(5, 2, 3) = 4$  corresponds to  $f^{\text{code}}(101\#10\#11) = 100$ .



# Turing-Computable Numerical Functions

Definition (Turing-Computable,  $f : \mathbb{N}_0^k \rightarrow_p \mathbb{N}_0$ )

A (partial) function  $f : \mathbb{N}_0^k \rightarrow_p \mathbb{N}_0$  is called **Turing-computable** if a DTM that computes  $f^{\text{code}}$  exists.

German: Turing-berechenbar

# Example: Turing-Computable Numerical Function

## Example

The following numerical functions are Turing-computable:

■  $\text{succ} : \mathbb{N}_0 \rightarrow_p \mathbb{N}_0$  with  $\text{succ}(n) := n + 1$

■  $\text{pred}_1 : \mathbb{N}_0 \rightarrow_p \mathbb{N}_0$  with  $\text{pred}_1(n) := \begin{cases} n - 1 & \text{if } n \geq 1 \\ 0 & \text{if } n = 0 \end{cases}$

■  $\text{pred}_2 : \mathbb{N}_0 \rightarrow_p \mathbb{N}_0$  with  $\text{pred}_2(n) := \begin{cases} n - 1 & \text{if } n \geq 1 \\ \text{undefined} & \text{if } n = 0 \end{cases}$

↪ blackboard/exercises

# More Turing-Computable Numerical Functions

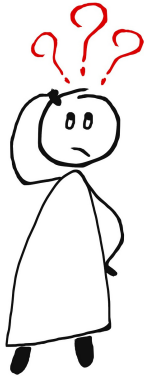
## Example

The following numerical functions are Turing-computable:

- $add : \mathbb{N}_0^2 \rightarrow_p \mathbb{N}_0$  with  $add(n_1, n_2) := n_1 + n_2$
- $sub : \mathbb{N}_0^2 \rightarrow_p \mathbb{N}_0$  with  $sub(n_1, n_2) := \max\{n_1 - n_2, 0\}$
- $mul : \mathbb{N}_0^2 \rightarrow_p \mathbb{N}_0$  with  $mul(n_1, n_2) := n_1 \cdot n_2$
- $div : \mathbb{N}_0^2 \rightarrow_p \mathbb{N}_0$  with  $div(n_1, n_2) := \begin{cases} \lceil \frac{n_1}{n_2} \rceil & \text{if } n_2 \neq 0 \\ \text{undefined} & \text{if } n_2 = 0 \end{cases}$

↪ sketch?

# Questions



Questions?

# Summary

# Summary

- main question: **what can a computer compute?**
- approach: investigate **formal models of computation**
- here: deterministic Turing machines
- **Turing-computable** function  $f : \Sigma^* \rightarrow_p \Sigma^*$ :  
there is a DTM that transforms every input  $w \in \Sigma^*$  into the output  $f(w)$  (undefined if DTM does not stop or stops in invalid configuration)
- **Turing-computable** function  $f : \mathbb{N}_0^k \rightarrow_p \mathbb{N}_0$ :  
ditto; numbers encoded in binary and separated by #