

Theory of Computer Science

C2. Regular Languages: Finite Automata

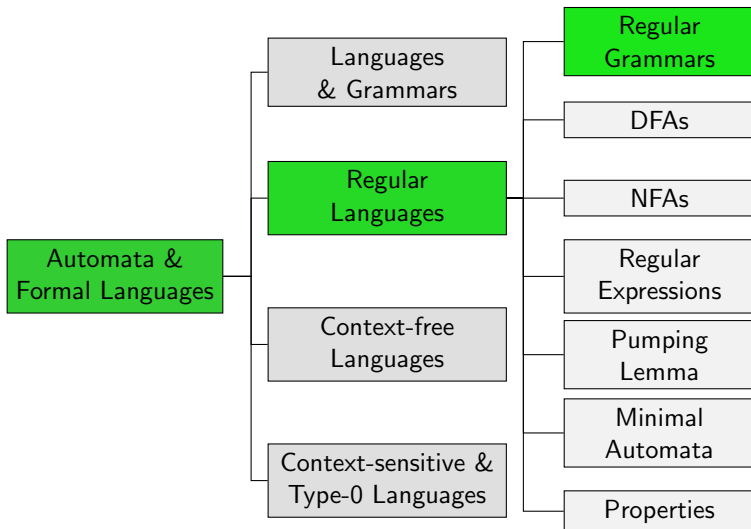
Gabriele Röger

University of Basel

March 20, 2019

Regular Grammars

Overview



Repetition: Regular Grammars

Definition (Regular Grammars)

A regular **grammar** is a 4-tuple $\langle \Sigma, V, P, S \rangle$ with

- 1 Σ finite alphabet of terminals
- 2 V finite set of variables (with $V \cap \Sigma = \emptyset$)
- 3 $P \subseteq (V \times (\Sigma \cup \Sigma V)) \cup \{\langle S, \varepsilon \rangle\}$ finite set of rules
- 4 if $S \rightarrow \varepsilon \in P$, there is no $X \in V, y \in \Sigma$ with $X \rightarrow yS \in P$
- 5 $S \in V$ start variable.

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Rule $X \rightarrow \varepsilon$ is only allowed if $X = S$ and S never occurs in the right-hand side of a rule.

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How restrictive is this?

Epsilon Rules

Theorem

For every grammar G with rules $P \subseteq V \times (\Sigma \cup \Sigma V \cup \{\epsilon\})$ there is a regular grammar G' with $\mathcal{L}(G) = \mathcal{L}(G')$.

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Let $G = \langle \Sigma, V, P, S \rangle$ be a grammar s.t. $P \subseteq V \times (\Sigma \cup \Sigma V \cup \{\varepsilon\})$.
Let $V_\varepsilon = \{A \in V \mid A \rightarrow \varepsilon \in P\}$.

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Let $V_\varepsilon = \{A \in V \mid A \rightarrow \varepsilon \in P\}$.

Let P' be the rule set that is created from P by removing all rules of the form $A \rightarrow \varepsilon$ ($A \neq S$). Additionally, for every rule of the form $B \rightarrow xA$ with $A \in V_\varepsilon, B \in V, x \in \Sigma$ we add a rule $B \rightarrow x$ to P' .

...

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Proof (continued).

Then $\mathcal{L}(G) = \mathcal{L}(\langle \Sigma, V, P', S \rangle)$ and P' contains no rule $A \rightarrow \varepsilon$ with $A \neq S$.
If $S \rightarrow \varepsilon \notin P$, we are done.



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For every grammar G with rules $P \subseteq V \times (\Sigma \cup \Sigma V \cup \{\varepsilon\})$ there is a regular grammar G' with $\mathcal{L}(G) = \mathcal{L}(G')$.

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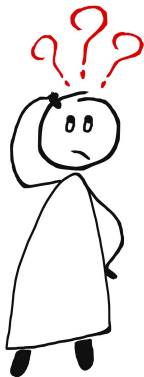
If $S \rightarrow \varepsilon \notin P$, we are done.

Otherwise, let S' be a new variable and construct P'' from P' by

- 1 replacing rules $X \rightarrow aS$ where $X \in V, a \in \Sigma$ with $X \rightarrow aS'$,
- 2 for every rule $S \rightarrow aX$ where $X \in V, a \in \Sigma$ adding the rule $S' \rightarrow aX$, and
- 3 for every rule $S \rightarrow a$ where $a \in \Sigma$ adding the rule $S' \rightarrow a$.

Then $\mathcal{L}(G) = \mathcal{L}(\langle \Sigma, V \cup \{S'\}, P'', S \rangle)$. □

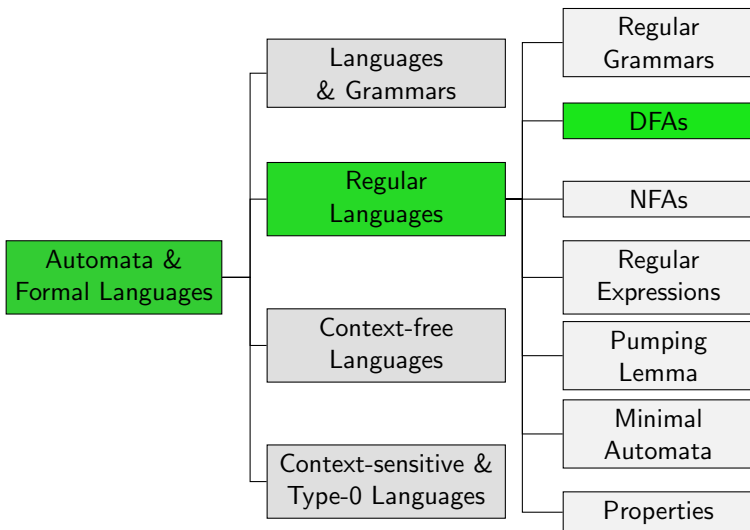
Questions



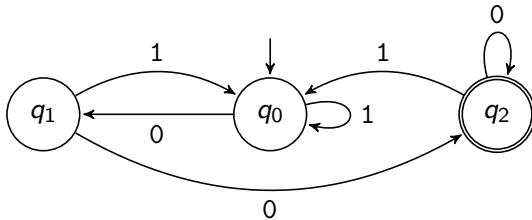
Questions?

DFAs

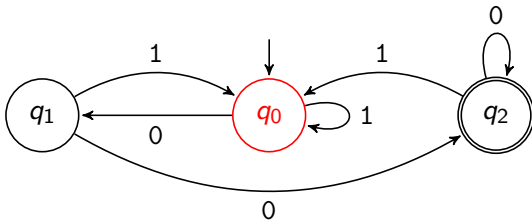
Overview



Finite Automata: Example



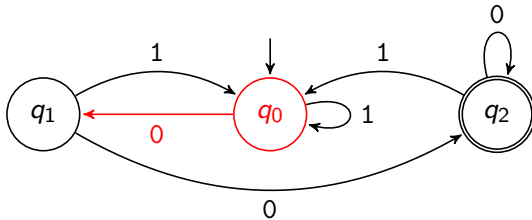
Finite Automata: Example



When reading the input 01100 the automaton visits the states

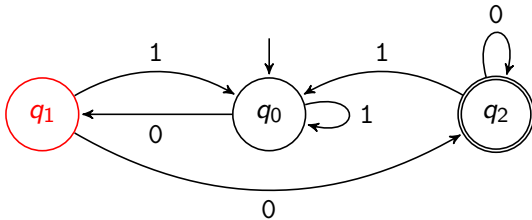
q0,

Finite Automata: Example



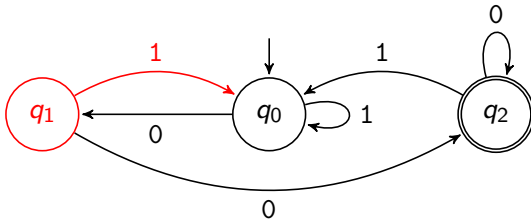
When reading the input **0**1100 the automaton visits the states q_0 ,

Finite Automata: Example



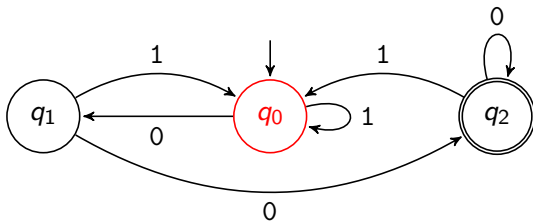
When reading the input 01100 the automaton visits the states q_0 , q_1 ,

Finite Automata: Example



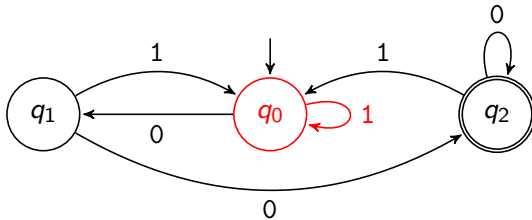
When reading the input 0**1**100 the automaton visits the states $q_0, q_1,$

Finite Automata: Example



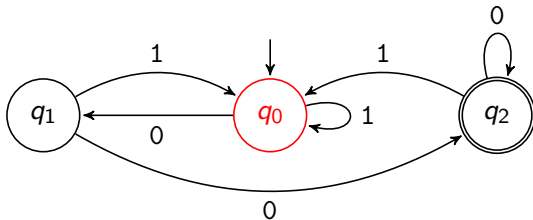
When reading the input 01100 the automaton visits the states
 q_0 , q_1 , q_0 ,

Finite Automata: Example



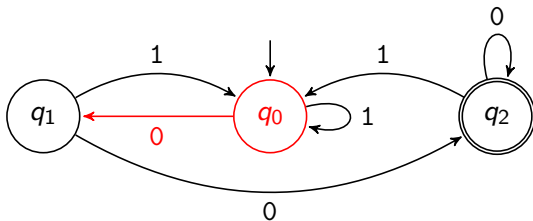
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Finite Automata: Example



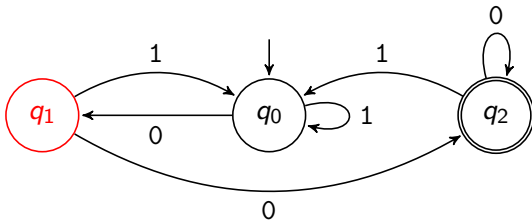
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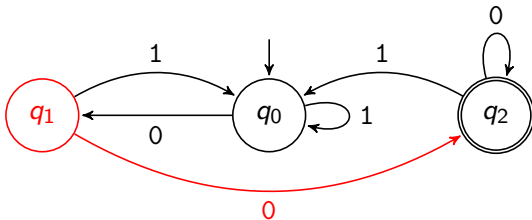
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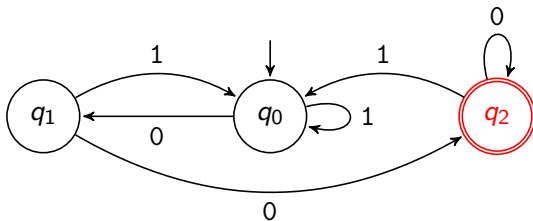
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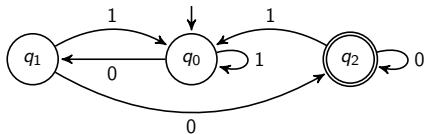
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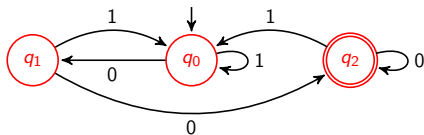


When reading the input 01100 the automaton visits the states $q_0, q_1, q_0, q_0, q_1, q_2$.

Finite Automata: Terminology and Notation

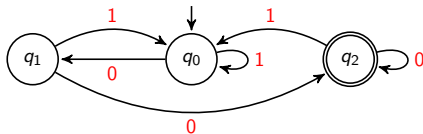


Finite Automata: Terminology and Notation



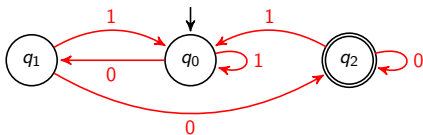
- states $Q = \{q_0, q_1, q_2\}$

Finite Automata: Terminology and Notation



- states $Q = \{q_0, q_1, q_2\}$
- input alphabet $\Sigma = \{0, 1\}$

Finite Automata: Terminology and Notation



- states $Q = \{q_0, q_1, q_2\}$
- input alphabet $\Sigma = \{0, 1\}$
- transition function δ

$$\delta(q_0, 0) = q_1$$

$$\delta(q_0, 1) = q_0$$

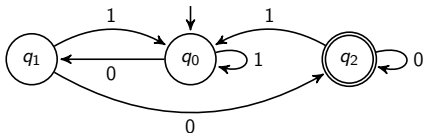
$$\delta(q_1, 0) = q_2$$

$$\delta(q_1, 1) = q_0$$

$$\delta(q_2, 0) = q_2$$

$$\delta(q_2, 1) = q_0$$

Finite Automata: Terminology and Notation



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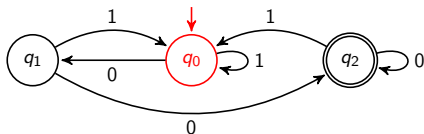
$$\delta(q_2, 0) = q_2$$

$$\delta(q_2, 1) = q_0$$

δ	0	1
q_0	q_1	q_0
q_1	q_2	q_0
q_2	q_2	q_0

table form of δ

Finite Automata: Terminology and Notation



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- transition function δ
- start state q_0

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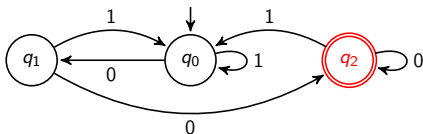
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Finite Automata: Terminology and Notation



- states $Q = \{q_0, q_1, q_2\}$
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- end states $\{q_2\}$

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Deterministic Finite Automaton: Definition

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- $\delta : Q \times \Sigma \rightarrow Q$ is the **transition function**
- $q_0 \in Q$ is the **start state**
- $E \subseteq Q$ is the set of **end states**

German: deterministischer endlicher Automat, Zustände, Eingabealphabet, Überführungs-/Übergangsfunktion, Startzustand, Endzustände

DFA: Recognized Words

Definition (Words Recognized by a DFA)

DFA $M = \langle Q, \Sigma, \delta, q_0, E \rangle$ **recognizes the word** $w = a_1 \dots a_n$ if there is a sequence of states $q'_0, \dots, q'_n \in Q$ with

- 1 $q'_0 = q_0$,
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German: DFA erkennt das Wort

DFA: Recognized Words

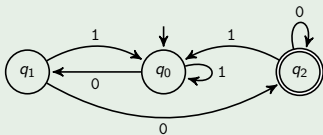
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Example



recognizes:

00
10010100
01000

does not recognize:

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DFA: Accepted Language

Definition (Language Accepted by a DFA)

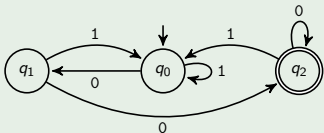
Let M be a deterministic finite automaton.
The **language accepted by M** is defined as
 $\mathcal{L}(M) = \{w \in \Sigma^* \mid w \text{ is recognized by } M\}$.

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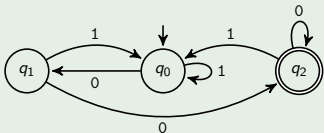


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Example



The DFA accepts the language
 $\{w \in \{0, 1\}^* \mid w \text{ ends with } 00\}$.

Languages Accepted by DFAs are Regular

Theorem

Every language accepted by a DFA is regular (type 3).

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Every language accepted by a DFA is regular (type 3).

Proof.

Let $M = \langle Q, \Sigma, \delta, q_0, E \rangle$ be a DFA.

We define a regular grammar G with $\mathcal{L}(G) = \mathcal{L}(M)$.

Define $G = \langle \Sigma, Q, P, q_0 \rangle$ where P contains

- a rule $q \rightarrow aq'$ for every $\delta(q, a) = q'$, and
- a rule $q \rightarrow \varepsilon$ for every $q \in E$.

(We can eliminate forbidden epsilon rules as described at the start of the chapter.)

...

Languages Accepted by DFAs are Regular

Theorem

Every language accepted by a DFA is regular (type 3).

Proof (continued).

For every $w = a_1 a_2 \dots a_n \in \Sigma^*$:

$w \in \mathcal{L}(M)$

iff there is a sequence of states q'_0, q'_1, \dots, q'_n with
 $q'_0 = q_0, q'_n \in E$ and $\delta(q'_{i-1}, a_i) = q'_i$ for all $i \in \{1, \dots, n\}$

iff there is a sequence of variables q'_0, q'_1, \dots, q'_n with
 q'_0 is start variable and we have $q'_0 \Rightarrow a_1 q'_1 \Rightarrow a_1 a_2 q'_2 \Rightarrow$
 $\dots \Rightarrow a_1 a_2 \dots a_n q'_n \Rightarrow a_1 a_2 \dots a_n$.

iff $w \in \mathcal{L}(G)$



Languages Accepted by DFAs are Regular

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iff $w \in \mathcal{L}(G)$



Example: blackboard

Question



Is the inverse true as well:
for every regular language, is there a
DFA that accepts it? That is, are the
languages accepted by DFAs **exactly** the
regular languages?

Question

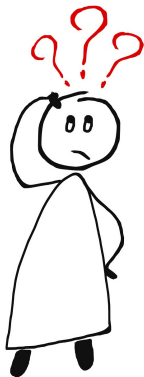


Is the inverse true as well:
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Yes!

We will prove this later (via a detour).

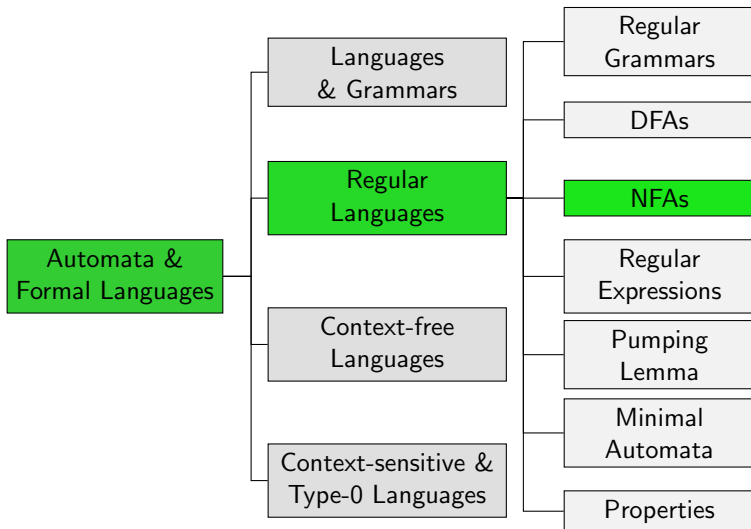
Questions



Questions?

NFAs

Overview

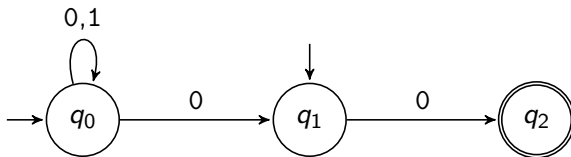


Nondeterministic Finite Automata

Why are DFAs called **deterministic** automata? What are **nondeterministic** automata, then?

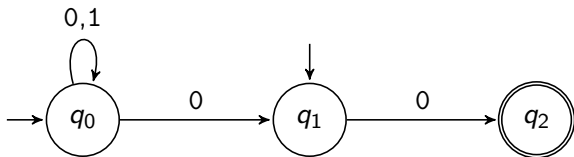


Nondeterministic Finite Automata: Example



differences to DFAs:

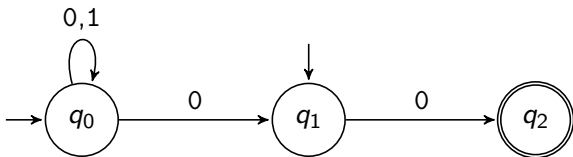
Nondeterministic Finite Automata: Example



differences to DFAs:

- **multiple** start states possible

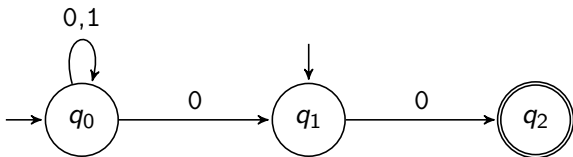
Nondeterministic Finite Automata: Example



differences to DFAs:

- **multiple** start states possible
- transition function δ can lead to **zero** or **more** successor states for the **same** $a \in \Sigma$

Nondeterministic Finite Automata: Example



differences to DFAs:

- **multiple** start states possible
- transition function δ can lead to **zero** or **more** successor states for the **same** $a \in \Sigma$
- automaton recognizes a word if there is **at least one** accepting sequence of states

Nondeterministic Finite Automaton: Definition

Definition (Nondeterministic Finite Automata)

A **nondeterministic finite automaton (NFA)** is a 5-tuple $M = \langle Q, \Sigma, \delta, S, E \rangle$ where

- Q is the finite set of **states**
- Σ is the **input alphabet** (with $Q \cap \Sigma = \emptyset$)
- $\delta : Q \times \Sigma \rightarrow \mathcal{P}(Q)$ is the transition function (mapping to the **power set** of Q)
- $S \subseteq Q$ is the set of **start states**
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German: nichtdeterministischer endlicher Automat

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DFAs are (essentially) a special case of NFAs.

NFA: Recognized Words

Definition (Words Recognized by an NFA)

NFA $M = \langle Q, \Sigma, \delta, S, E \rangle$ **recognizes the word** $w = a_1 \dots a_n$ if there is a sequence of states $q'_0, \dots, q'_n \in Q$ with

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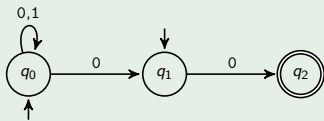
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NFA: Accepted Language

Definition (Language Accepted by an NFA)

Let $M = \langle Q, \Sigma, \delta, S, E \rangle$ be a nondeterministic finite automaton.

The **language accepted by M** is defined as

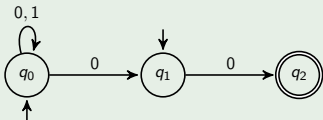
$\mathcal{L}(M) = \{w \in \Sigma^* \mid w \text{ is recognized by } M\}$.

NFA: Accepted Language

Definition (Language Accepted by an NFA)

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Example



NFA: Accepted Language

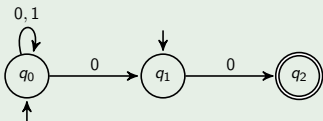
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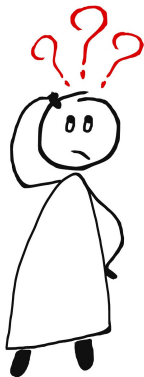
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Example



The NFA accepts the language
 $\{w \in \{0, 1\}^* \mid w = 0 \text{ or } w \text{ ends with } 00\}$.

Questions



Questions?

NFAs are No More Powerful than DFAs

Theorem (Rabin, Scott)

Every language accepted by an NFA is also accepted by a DFA.

NFAs are No More Powerful than DFAs

Theorem (Rabin, Scott)

Every language accepted by an NFA is also accepted by a DFA.

Proof.

For every NFA $M = \langle Q, \Sigma, \delta, S, E \rangle$ we can construct a DFA $M' = \langle Q', \Sigma, \delta', q'_0, E' \rangle$ with $\mathcal{L}(M) = \mathcal{L}(M')$. Here M' is defined as follows:

- $Q' := \mathcal{P}(Q)$ (the power set of Q)
- $q'_0 := S$
- $E' := \{Q \subseteq Q \mid Q \cap E \neq \emptyset\}$
- For all $Q \in Q'$: $\delta'(Q, a) := \bigcup_{q \in Q} \delta(q, a)$

NFAs are No More Powerful than DFAs

Theorem (Rabin, Scott)

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Proof (continued).

For every $w = a_1 a_2 \dots a_n \in \Sigma^*$:

$w \in \mathcal{L}(M)$

iff there is a sequence of states q_0, q_1, \dots, q_n with

$q_0 \in S$, $q_n \in E$ and $q_i \in \delta(q_{i-1}, a_i)$ for all $i \in \{1, \dots, n\}$

iff there is a sequence of subsets Q_0, Q_1, \dots, Q_n with

$Q_0 = q'_0$, $Q_n \in E'$ and $\delta'(Q_{i-1}, a_i) = Q_i$ for all $i \in \{1, \dots, n\}$

iff $w \in \mathcal{L}(M')$



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Example: blackboard

NFAs are More Compact than DFAs

Example

For $k \geq 1$ consider the language

$L_k = \{w \in \{0, 1\}^* \mid |w| \geq k \text{ and the } k\text{-th last symbol of } w \text{ is } 0\}$.

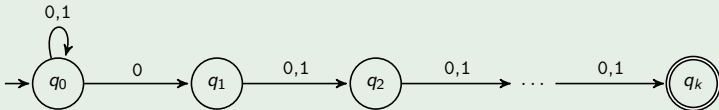
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The language L_k can be accepted by an NFA with $k + 1$ states:



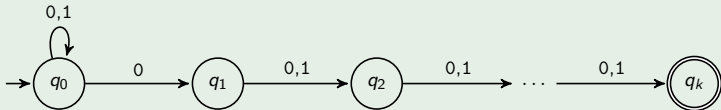
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There is no DFA with less than 2^k states that accepts L_k (without proof).

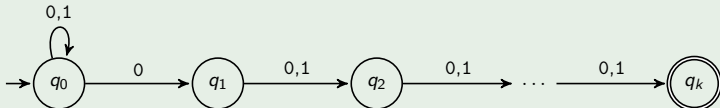
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There is no DFA with less than 2^k states that accepts L_k
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NFAs can often represent languages more compactly than DFAs.

Regular Grammars are No More Powerful than NFAs

Theorem

For every regular grammar G there is an NFA M with $\mathcal{L}(G) = \mathcal{L}(M)$.

Regular Grammars are No More Powerful than NFAs

Theorem

For every regular grammar G there is an NFA M with $\mathcal{L}(G) = \mathcal{L}(M)$.

Proof.

Let $G = \langle \Sigma, V, P, S \rangle$ be a regular grammar.

Define NFA $M = \langle Q, \Sigma, \delta, S', E \rangle$ with

$$Q = V \cup \{X\}, \quad X \notin V$$

$$S' = \{S\}$$

$$E = \begin{cases} \{S, X\} & \text{if } S \rightarrow \varepsilon \in P \\ \{X\} & \text{if } S \rightarrow \varepsilon \notin P \end{cases}$$

$$B \in \delta(A, a) \text{ if } A \rightarrow aB \in P$$

$$X \in \delta(A, a) \text{ if } A \rightarrow a \in P$$

Regular Grammars are No More Powerful than NFAs

Theorem

For every regular grammar G there is an NFA M with $\mathcal{L}(G) = \mathcal{L}(M)$.

Proof (continued).

For every $w = a_1a_2 \dots a_n \in \Sigma^*$ with $n \geq 1$:

$w \in \mathcal{L}(G)$

iff there is a sequence on variables A_1, A_2, \dots, A_{n-1} with

$$S \Rightarrow a_1A_1 \Rightarrow a_1a_2A_2 \Rightarrow \dots \Rightarrow a_1a_2 \dots a_{n-1}A_{n-1} \Rightarrow a_1a_2 \dots a_n.$$

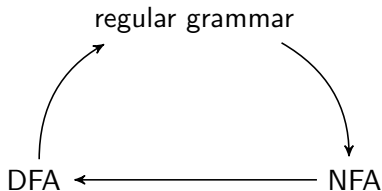
iff there is a sequence of variables A_1, A_2, \dots, A_{n-1} with

$$A_1 \in \delta(S, a_1), A_2 \in \delta(A_1, a_2), \dots, X \in \delta(A_{n-1}, a_n).$$

iff $w \in \mathcal{L}(M)$.

Case $w = \varepsilon$ is also covered because $S \in E$ iff $S \rightarrow \varepsilon \in P$. □

Finite Automata and Regular Languages



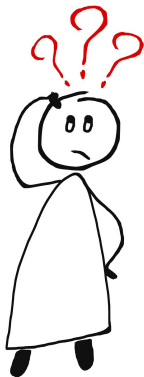
In particular, this implies:

Corollary

\mathcal{L} regular \iff \mathcal{L} is accepted by a DFA.

\mathcal{L} regular \iff \mathcal{L} is accepted by an NFA.

Questions



Questions?

Summary

Summary

- We now know **three formalisms** that all **describe exactly the regular languages**: regular grammars, DFAs and NFAs
- We will get to know a fourth formalism in the next chapter.
- **DFAs** are automata where **every state transition is uniquely determined**.
- **NFAs** recognize a word if there is **at least one accepting sequence of states**.