Foundations of Artificial Intelligence
29. Propositional Logic: Basics

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Classification

classification:

Propositional Logic

environment:

- **static** vs. dynamic
- **deterministic** vs. non-deterministic vs. stochastic
- **fully** vs. partially vs. not **observable**
- **discrete** vs. continuous
- **single-agent** vs. multi-agent

problem solving method:

- problem-specific vs. **general** vs. learning

(applications also in more complex environments)

Propositional Logic: Overview

Chapter overview: propositional logic

- 29. Basics
- 30. Reasoning and Resolution
- 31. DPLL Algorithm
- 32. Local Search and Outlook

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Motivation

Propositional Logic: Motivation

propositional logic

- modeling and representing problems and knowledge
- basics for general problem descriptions and solving strategies
  ($\rightsquigarrow$ automated planning $\rightsquigarrow$ later in this course)
- allows for automated reasoning

German: Aussagenlogik, automatisches Schliessen

Relationship to CSPs

- previous topic: constraint satisfaction problems
- satisfiability problem in propositional logic can be viewed as non-binary CSP over $\{\mathbf{F}, \mathbf{T}\}$
- formula encodes constraints
- solution: satisfying assignment of values to variables
- SAT algorithms for this problem: ⇝ DPLL (Wednesday)

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Propositional Logic: Description of State Spaces

propositional variables for missionaries and cannibals problem:

```
two-missionaries-on-left-shore
one-cannibal-on-left-shore
boat-on-left-shore
...
```

- problem description for general problem solvers
- states represented as truth values of atomic propositions

German: Aussagenvariablen

Propositional Logic: Intuition

propositions: atomic statements over the world
that cannot be divided further

Propositions with logical connectives like
"and", "or" and "not" form the propositional formulas.

German: logische Verknüpfungen

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Syntax

Syntax

Σ alphabet of propositions
(e.g., $\{P, Q, R, \dots\}$ or $\{X_1, X_2, X_3, \dots\}$).

Definition (propositional formula)

- $\top$ and $\bot$ are formulas.
- Every proposition in Σ is an (atomic) formula.
- If φ is a formula, then $\neg\varphi$ is a formula (negation).
- If φ and ψ are formulas, then so are
  - $(\varphi \wedge \psi)$ (conjunction)
  - $(\varphi \vee \psi)$ (disjunction)
  - $(\varphi \to \psi)$ (implication)

German: aussagenlogische Formel, atomare Formel, Konjunktion, Disjunktion, Implikation

binding strength: $(\neg) > (\wedge) > (\vee) > (\to)$
(may omit redundant parentheses)

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Semantics

A formula is true or false,
depending on the interpretation of the propositions.

Semantics: Intuition

- A proposition p is either true or false.
  The truth value of p is determined by an interpretation.

- The truth value of a formula follows from
  the truth values of the propositions.

Example

$\varphi = (P \vee Q) \wedge R$

- If P and Q are false, then φ is false
  (independent of the truth value of R).

- If P and R are true, then φ is true
  (independent of the truth value of Q).

Semantics: Formally

- defined over interpretation $I : \Sigma \rightarrow \{\mathbf{T}, \mathbf{F}\}$

- interpretation I: assignment of propositions in Σ

- When is a formula φ true under interpretation I?
  symbolically: When does $I \models \varphi$ hold?

German: Interpretation, Belegung

Semantics: Formally

Definition ($I \models \varphi$)

- $I \models \top$ and $I \not\models \bot$
- $I \models P$ iff $I(P) = \mathbf{T}$ for $P \in \Sigma$
- $I \models \neg\varphi$ iff $I \not\models \varphi$
- $I \models (\varphi \wedge \psi)$ iff $I \models \varphi$ and $I \models \psi$
- $I \models (\varphi \vee \psi)$ iff $I \models \varphi$ or $I \models \psi$
- $I \models (\varphi \rightarrow \psi)$ iff $I \not\models \varphi$ or $I \models \psi$
- $I \models \Phi$ for a set of formulas Φ iff $I \models \varphi$ for all $\varphi \in \Phi$

German: I erfüllt φ, φ gilt unter I

Examples

Example (Interpretation I)

$I = \{P \mapsto \textbf{T}, Q \mapsto \textbf{T}, R \mapsto \textbf{F}, S \mapsto \textbf{F}\}$

Which formulas are true under I?

- $\varphi_1 = \neg(P \wedge Q) \wedge (R \wedge \neg S)$. Does $I \models \varphi_1$ hold?
- $\varphi_2 = (P \wedge Q) \wedge \neg(R \wedge \neg S)$. Does $I \models \varphi_2$ hold?
- $\varphi_3 = (R \rightarrow P)$. Does $I \models \varphi_3$ hold?

Terminology

> **Definition (model)**
>
> An interpretation I is called a model of φ if $I \models \varphi$.

German: Modell

> **Definition (satisfiable etc.)**
>
> A formula φ is called
>
> - satisfiable if there is an interpretation I such that $I \models \varphi$.
> - unsatisfiable if φ is not satisfiable.
> - falsifiable if there is an interpretation I such that $I \not\models \varphi$.
> - valid (= a tautology) if $I \models \varphi$ for all interpretations I.

German: erfüllbar, unerfüllbar, falsifizierbar,
allgemeingültig (gültig, Tautologie)

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Terminology

Definition (logical equivalence)

Formulas φ and ψ are called logically equivalent ($\varphi \equiv \psi$)
if for all interpretations I: $I \models \varphi$ iff $I \models \psi$.

German: logisch äquivalent

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Truth Tables

Truth Tables

How to determine automatically if a given formula is (un)satisfiable, falsifiable, valid?

⤳ simple method: truth tables

example: Is $\varphi = ((P \vee H) \wedge \neg H) \to P$ valid?

P	H	$P \vee H$	$((P \vee H) \wedge \neg H)$	$((P \vee H) \wedge \neg H) \to P$
F	**F**	**F**	**F**	**T**
F	**T**	**T**	**F**	**T**
T	**F**	**T**	**T**	**T**
T	**T**	**T**	**F**	**T**

$I \models \varphi$ for all interpretations I ⤳ φ is valid.

- satisfiability, falsifiability, unsatisfiability?

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Normal Forms

Normal Forms: Terminology

Definition (literal)

If $P \in \Sigma$, then the formulas P and $\neg P$ are called literals.

P is called positive literal, $\neg P$ is called negative literal.

The complementary literal to P is $\neg P$ and vice versa.
For a literal ℓ, the complementary literal to ℓ is denoted with $\bar{\ell}$.

German: Literal, positives/negatives/komplementäres Literal

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Normal Forms: Terminology

Definition (clause)

A disjunction of 0 or more literals is called a clause.
The empty clause ⊥ is also written as □.
Clauses consisting of only one literal are called unit clauses.

German: Klausel

Definition (monomial)

A conjunction of 0 or more literals is called a monomial.

German: Monom

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Normal Forms

Definition (normal forms)

A formula φ is in conjunctive normal form (CNF, clause form)
if φ is a conjunction of 0 or more clauses:

$$\varphi = \bigwedge_{i=1}^{n} \left(\bigvee_{j=1}^{m_i} \ell_{i,j} \right)$$

A formula φ is in disjunctive normal form (DNF)
if φ is a disjunction of 0 or more monomials:

$$\varphi = \bigvee_{i=1}^{n} \left(\bigwedge_{j=1}^{m_i} \ell_{i,j} \right)$$

German: konjunktive Normalform, disjunktive Normalform

Normal Forms

For every propositional formula, there exists
a logically equivalent propositional formula in CNF and in DNF.

Conversion to CNF

important rules for conversion to CNF:

- $(\varphi \to \psi) \equiv (\neg\varphi \vee \psi)$ $((\to)\text{-elimination})$
- $\neg(\varphi \wedge \psi) \equiv (\neg\varphi \vee \neg\psi)$ (De Morgan)
- $\neg(\varphi \vee \psi) \equiv (\neg\varphi \wedge \neg\psi)$ (De Morgan)
- $\neg\neg\varphi \equiv \varphi$ (double negation)
- $((\varphi \wedge \psi) \vee \eta) \equiv ((\varphi \vee \eta) \wedge (\psi \vee \eta))$ (distributivity)

There are formulas φ for which every logically equivalent formula
in CNF and DNF is exponentially longer than φ.

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Summary

Summary (1)

- **Propositional logic** forms the basis for a general representation of problems and knowledge.
- **Propositions** (atomic formulas) are statements over the world which cannot be divided further.
- **Propositional formulas** combine atomic formulas with ¬, ∧, ∨, → to more complex statements.
- **Interpretations** determine which atomic formulas are true and which ones are false.

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Summary (2)

- important terminology:
  - model
  - satisfiable, unsatisfiable, falsifiable, valid
  - logically equivalent
- different kinds of formulas:
  - atomic formulas and literals
  - clauses and monomials
  - conjunctive normal form and disjunctive normal form