

Planning and Optimization

E7. Linear & Integer Programming

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Examples

Linear Program: Example Maximization Problem

Example

maximize $2x - 3y + z$ subject to

$$x + 2y + z \leq 10$$

$$x - z \leq 0$$

$$x \geq 0, \quad y \geq 0, \quad z \geq 0$$

↪ unique optimal solution:

$x = 5, y = 0, z = 5$ (objective value 15)

Example: Diet Problem

- n different types of food F_1, \dots, F_n
- m different nutrients N_1, \dots, N_m
- The minimum daily requirement of nutrient N_j is r_j .
- The amount of nutrient N_j in one unit of food F_i is a_{ij} .
- One unit of food F_i costs c_i .

How to supply the required nutrients at minimum cost?

Example: Diet Problem

- Use LP variable x_i for the number of units of food F_i purchased per day.
- The cost per day is $\sum_{i=1}^n c_i x_i$.
- The amount of nutrient N_j in this diet is $\sum_{i=1}^n a_{ij} x_i$.
- The minimum daily requirement for each nutrient N_j must be met: $\sum_{i=1}^n a_{ij} x_i \geq r_j$ for $1 \leq j \leq m$
- We can't buy negative amounts of food: $x_i \geq 0$ for $1 \leq i \leq n$
- We want to minimize the cost of food.

Diet Problem: Linear Program

Example (Linear Program for Diet Problem)

minimize $\sum_{i=1}^n c_i x_i$ subject to

$$\sum_{i=1}^n a_{ij} x_i \geq r_j \quad \text{for } 1 \leq j \leq m$$

$$x_i \geq 0 \quad \text{for } 1 \leq i \leq n$$

Linear Programs

Linear Programs and Integer Programs

Linear Program

A **linear program** (LP) consists of:

- a finite set of **real-valued variables** V
- a finite set of **linear inequalities** (constraints) over V
- an **objective function**, which is a linear combination of V
- which should be **minimized** or **maximized**.

Integer program (IP): ditto, but with some **integer-valued** variables

Standard Maximum Problem

Normal form for maximization problems:

Definition (Standard Maximum Problem)

Find values for x_1, \dots, x_n , to maximize

$$c_1x_1 + c_2x_2 + \cdots + c_nx_n$$

subject to the constraints

$$a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n \leq b_1$$

$$a_{21}x_1 + a_{22}x_2 + \cdots + a_{2n}x_n \leq b_2$$

$$\vdots$$

$$a_{m1}x_1 + a_{m2}x_2 + \cdots + a_{mn}x_n \leq b_m$$

and $x_1 \geq 0, x_2 \geq 0, \dots, x_n \geq 0$.

Standard Maximum Problem: Matrix and Vectors

A standard maximum problem is often given by

- an m -vector $\mathbf{b} = \langle b_1, \dots, b_m \rangle^T$,
- an n -vector $\mathbf{c} = \langle c_1, \dots, c_n \rangle^T$,
- and an $m \times n$ matrix

$$\mathbf{A} = \begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{pmatrix}$$

- Then the problem is to find a vector $\mathbf{x} = \langle x_1, \dots, x_n \rangle^T$ to maximize $\mathbf{c}^T \mathbf{x}$ subject to $\mathbf{A}\mathbf{x} \leq \mathbf{b}$ and $\mathbf{x} \geq \mathbf{0}$.

Standard Minimum Problem

Normal form for minimization problems:

Definition (Standard Minimum Problem)

Find values for y_1, \dots, y_m , to minimize

$$b_1 y_1 + b_2 y_2 + \cdots + b_m y_m$$

subject to the constraints

$$y_1 a_{11} + y_2 a_{21} + \cdots + y_m a_{m1} \geq c_1$$

$$y_1 a_{12} + y_2 a_{22} + \cdots + y_m a_{m2} \geq c_2$$

$$\vdots$$

$$y_1 a_{1n} + y_2 a_{2n} + \cdots + y_m a_{mn} \geq c_n$$

and $y_1 \geq 0, y_2 \geq 0, \dots, y_m \geq 0$.

Standard Minimum Problem: Matrix and Vectors

- A standard minimum problem is defined by the same matrix \mathbf{A} and vectors \mathbf{b}, \mathbf{c} as a maximum problem.
- The problem is to find a vector $\mathbf{y} = \langle y_1, \dots, y_m \rangle^T$ to minimize $\mathbf{y}^T \mathbf{b}$ subject to $\mathbf{y}^T \mathbf{A} \geq \mathbf{c}$ and $\mathbf{y} \geq \mathbf{0}$.

Terminology

- A vector \mathbf{x} for a maximum problem or \mathbf{y} for a minimum problem is **feasible** if it satisfies the constraints.
- A linear program is **feasible** if there is such a feasible vector. Otherwise it is **infeasible**.
- A feasible maximum (resp. minimum) problem is **unbounded** if the objective function can assume arbitrarily large positive (resp. negative) values at feasible vectors. Otherwise it is **bounded**.
- The **objective value** of a bounded feasible maximum (resp. minimum) problem is the maximum (resp. minimum) value of the objective function with a feasible vector.

Standard Problems are a Normal Form

All linear programs can be converted into a standard maximum problem:

- To transform a **minimum problem** into a maximum problem, multiply the objective function by -1 .
- Transform constraints $a_{i1}x_1 + \cdots + a_{in}x_n \geq b_i$ to $(-a_{i1})x_1 + \cdots + (-a_{in})x_n \leq -b_i$.
- Solve **equality constraints** $a_{i1}x_1 + \cdots + a_{in}x_n = b_i$ for some x_j with $a_{ij} \neq 0$ and substitute this solution wherever x_j appears.
- If a variable x can be **negative**, introduce variables $x' \geq 0$ and $x'' \geq 0$ and replace x everywhere with $x' - x''$.

Solving Linear Programs and Integer Programs

Complexity:

- LP solving is a **polynomial-time** problem.
- Finding solutions for IPs is **NP-complete**.

Common idea:

- Approximate IP solution with corresponding LP (**LP relaxation**).

LP Relaxation

Theorem (LP Relaxation)

*The **LP relaxation** of an integer program is the problem that arises by removing the requirement that variables are integer-valued.*

*For a **maximization** (resp. minimization) problem, the objective value of the LP relaxation is an **upper** (resp. lower) **bound** on the value of the IP.*

Proof idea.

Every feasible vector for the IP is also feasible for the LP.



Duality

Some LP Theory: Duality

Some LP theory: Every LP has an alternative view (its **dual**).

- roughly: variables and constraints swap roles
- dual of maximization LP is minimization LP and vice versa
- dual of dual: original LP

Dual Problem

Definition (Dual Problem)

The **dual** of the standard maximum problem

$$\text{maximize } \mathbf{c}^T \mathbf{x} \text{ subject to } \mathbf{A}\mathbf{x} \leq \mathbf{b} \text{ and } \mathbf{x} \geq \mathbf{0}$$

is the minimum problem

$$\text{minimize } \mathbf{y}^T \mathbf{b} \text{ subject to } \mathbf{y}^T \mathbf{A} \geq \mathbf{c}^T \text{ and } \mathbf{y} \geq \mathbf{0}$$

Dual for Diet Problem

Example (Dual of Linear Program for Diet Problem)

maximize $\sum_{j=1}^m y_j r_j$ subject to

$$\sum_{j=1}^m a_{ij} y_j \leq c_i \quad \text{for } 1 \leq i \leq n$$

$$y_j \geq 0 \quad \text{for } 1 \leq j \leq m$$

Duality Theorem

Theorem (Duality Theorem)

*If a standard linear program is **bounded feasible**, then so is its dual, and their **objective values are equal**.*

(Proof omitted.)

The dual provides a different perspective on a problem.

Dual for Diet Problem: Interpretation

Example (Dual of Linear Program for Diet Problem)

maximize $\sum_{j=1}^m y_j r_j$ subject to

$$\sum_{j=1}^m a_{ij} y_j \leq c_i \quad \text{for } 1 \leq i \leq n$$

$$y_j \geq 0 \quad \text{for } 1 \leq j \leq m$$

Find nutrient prices that maximize total worth of daily nutrients.
The value of nutrients in food F_i may not exceed the cost of F_i .

Summary

Summary

- Linear (and integer) programs consist of an objective function that should be maximized or minimized subject to a set of given linear constraints.
- Finding solutions for integer programs is NP-complete.
- LP solving is a polynomial time problem.
- The dual of a maximization LP is a minimization LP and vice versa.
- The dual of a bounded feasible LP has the same objective value.

Further Reading

The slides in this chapter are based on the following excellent tutorial on LP solving:



[Thomas S. Ferguson.](#)

Linear Programming – A Concise Introduction.

[UCLA, unpublished document available online.](#)