

# Planning and Optimization

## E7. Linear & Integer Programming

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# Examples

## Linear Program: Example Maximization Problem

## Example

maximize  $2x - 3y + z$  subject to

$$\begin{aligned} x + 2y + z &\leq 10 \\ x - z &\leq 0 \end{aligned}$$

$$x \geq 0, \quad y \geq 0, \quad z \geq 0$$

~~&gt; unique optimal solution:

 $x = 5, y = 0, z = 5$  (objective value 15)

## Example: Diet Problem

- $n$  different types of food  $F_1, \dots, F_n$
- $m$  different nutrients  $N_1, \dots, N_m$
- The minimum daily requirement of nutrient  $N_j$  is  $r_j$ .
- The amount of nutrient  $N_j$  in one unit of food  $F_i$  is  $a_{ij}$ .
- One unit of food  $F_i$  costs  $c_i$ .

How to supply the required nutrients at minimum cost?

## Example: Diet Problem

- Use LP variable  $x_i$  for the number of units of food  $F_i$  purchased per day.
- The cost per day is  $\sum_{i=1}^n c_i x_i$ .
- The amount of nutrient  $N_j$  in this diet is  $\sum_{i=1}^n a_{ij} x_i$ .
- The minimum daily requirement for each nutrient  $N_j$  must be met:  $\sum_{i=1}^n a_{ij} x_i \geq r_j$  for  $1 \leq j \leq m$
- We can't buy negative amounts of food:  $x_i \geq 0$  for  $1 \leq i \leq n$
- We want to minimize the cost of food.

## Diet Problem: Linear Program

## Example (Linear Program for Diet Problem)

$$\text{minimize} \quad \sum_{i=1}^n c_i x_i \quad \text{subject to}$$

$$\sum_{i=1}^n a_{ij} x_i \geq r_j \quad \text{for } 1 \leq j \leq m$$

$$x_i \geq 0 \quad \text{for } 1 \leq i \leq n$$

# Linear Programs

# Linear Programs and Integer Programs

## Linear Program

A **linear program (LP)** consists of:

- a finite set of **real-valued variables**  $V$
- a finite set of **linear inequalities** (constraints) over  $V$
- an **objective function**, which is a linear combination of  $V$
- which should be **minimized** or **maximized**.

**Integer program (IP):** ditto, but with some **integer-valued** variables

# Standard Maximum Problem

Normal form for maximization problems:

## Definition (Standard Maximum Problem)

Find values for  $x_1, \dots, x_n$ , to maximize

$$c_1 x_1 + c_2 x_2 + \dots + c_n x_n$$

subject to the constraints

$$a_{11} x_1 + a_{12} x_2 + \dots + a_{1n} x_n \leq b_1$$

$$a_{21} x_1 + a_{22} x_2 + \dots + a_{2n} x_n \leq b_2$$

⋮

$$a_{m1} x_1 + a_{m2} x_2 + \dots + a_{mn} x_n \leq b_m$$

and  $x_1 \geq 0, x_2 \geq 0, \dots, x_n \geq 0$ .

# Standard Maximum Problem: Matrix and Vectors

A standard maximum problem is often given by

- an  $m$ -vector  $\mathbf{b} = \langle b_1, \dots, b_m \rangle^T$ ,
- an  $n$ -vector  $\mathbf{c} = \langle c_1, \dots, c_n \rangle^T$ ,
- and an  $m \times n$  matrix

$$\mathbf{A} = \begin{pmatrix} a_{11} & a_{12}x_2 & \dots & a_{1n} \\ a_{21} & a_{22}x_2 & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{pmatrix}$$

- Then the problem is to find a vector  $\mathbf{x} = \langle x_1, \dots, x_n \rangle^T$  to maximize  $\mathbf{c}^T \mathbf{x}$  subject to  $\mathbf{A}\mathbf{x} \leq \mathbf{b}$  and  $\mathbf{x} \geq \mathbf{0}$ .

# Standard Minimum Problem

Normal form for minimization problems:

## Definition (Standard Minimum Problem)

Find values for  $y_1, \dots, y_m$ , to minimize

$$b_1 y_1 + b_2 y_2 + \dots + b_m y_m$$

subject to the constraints

$$y_1 a_{11} + y_2 a_{21} + \dots + y_m a_{m1} \geq c_1$$

$$y_1 a_{12} + y_2 a_{22} + \dots + y_m a_{m2} \geq c_2$$

⋮

$$y_1 a_{1n} + y_2 a_{2n} + \dots + y_m a_{mn} \geq c_n$$

and  $y_1 \geq 0, y_2 \geq 0, \dots, y_m \geq 0$ .

# Standard Minimum Problem: Matrix and Vectors

- A standard minimum problem is defined by the same matrix  $\mathbf{A}$  and vectors  $\mathbf{b}, \mathbf{c}$  as a maximum problem.
- The problem is to find a vector  $\mathbf{y} = \langle y_1, \dots, y_m \rangle^T$  to minimize  $\mathbf{y}^T \mathbf{b}$  subject to  $\mathbf{y}^T \mathbf{A} \geq \mathbf{c}$  and  $\mathbf{y} \geq \mathbf{0}$ .

# Terminology

- A vector  $x$  for a maximum problem or  $y$  for a minimum problem is **feasible** if it satisfies the constraints.
- A linear program is **feasible** if there is such a feasible vector. Otherwise it is **infeasible**.
- A feasible maximum (resp. minimum) problem is **unbounded** if the objective function can assume arbitrarily large positive (resp. negative) values at feasible vectors. Otherwise it is **bounded**.
- The **objective value** of a bounded feasible maximum (resp. minimum) problem is the maximum (resp. minimum) value of the objective function with a feasible vector.

# Standard Problems are a Normal Form

All linear programs can be converted into a standard maximum problem:

- To transform a **minimum problem** into a maximum problem, multiply the objective function by  $-1$ .
- Transform constraints  $a_{i1}x_1 + \dots + a_{in}x_n \geq b_i$  to  $(-a_{i1})x_1 + \dots + (-a_{in})x_n \leq -b_i$ .
- Solve **equality constraints**  $a_{i1}x_1 + \dots + a_{in}x_n = b_i$  for some  $x_j$  with  $a_{ij} \neq 0$  and substitute this solution wherever  $x_j$  appears.
- If a variable  $x$  can be **negative**, introduce variables  $x' \geq 0$  and  $x'' \geq 0$  and replace  $x$  everywhere with  $x' - x''$ .

# Solving Linear Programs and Integer Programs

## Complexity:

- LP solving is a **polynomial-time** problem.
- Finding solutions for IPs is **NP-complete**.

## Common idea:

- Approximate IP solution with corresponding LP  
(**LP relaxation**).

# LP Relaxation

## Theorem (LP Relaxation)

*The LP relaxation of an integer program is the problem that arises by removing the requirement that variables are integer-valued.*

*For a maximization (resp. minimization) problem, the objective value of the LP relaxation is an upper (resp. lower) bound on the value of the IP.*

## Proof idea.

Every feasible vector for the IP is also feasible for the LP.



# Duality

# Some LP Theory: Duality

**Some LP theory:** Every LP has an alternative view (its **dual**).

- roughly: variables and constraints swap roles
- dual of maximization LP is minimization LP and vice versa
- dual of dual: original LP

# Dual Problem

## Definition (Dual Problem)

The **dual** of the standard maximum problem

$$\text{maximize } \mathbf{c}^T \mathbf{x} \text{ subject to } \mathbf{A}\mathbf{x} \leq \mathbf{b} \text{ and } \mathbf{x} \geq \mathbf{0}$$

is the minimum problem

$$\text{minimize } \mathbf{y}^T \mathbf{b} \text{ subject to } \mathbf{y}^T \mathbf{A} \geq \mathbf{c}^T \text{ and } \mathbf{y} \geq \mathbf{0}$$

# Dual for Diet Problem

## Example (Dual of Linear Program for Diet Problem)

$$\text{maximize} \quad \sum_{j=1}^m y_j r_j \quad \text{subject to}$$

$$\sum_{j=1}^m a_{ij} y_j \leq c_i \quad \text{for } 1 \leq i \leq n$$

$$y_j \geq 0 \quad \text{for } 1 \leq j \leq m$$

# Duality Theorem

## Theorem (Duality Theorem)

*If a standard linear program is **bounded feasible**, then so is its dual, and their **objective values** are equal.*

(Proof omitted.)

The dual provides a different perspective on a problem.

# Dual for Diet Problem: Interpretation

## Example (Dual of Linear Program for Diet Problem)

$$\text{maximize} \quad \sum_{j=1}^m y_j r_j \quad \text{subject to}$$

$$\sum_{j=1}^m a_{ij} y_j \leq c_i \quad \text{for } 1 \leq i \leq n$$

$$y_j \geq 0 \quad \text{for } 1 \leq j \leq m$$

Find nutrient prices that maximize total worth of daily nutrients.  
The value of nutrients in food  $F_i$  may not exceed the cost of  $F_i$ .

# Summary

# Summary

- Linear (and integer) programs consist of an **objective function** that should be **maximized or minimized** subject to a set of given **linear constraints**.
- Finding solutions for **integer** programs is **NP-complete**.
- **LP solving** is a **polynomial time** problem.
- The dual of a maximization LP is a minimization LP and vice versa.
- The **dual** of a bounded feasible LP has the **same objective value**.

# Further Reading

The slides in this chapter are based on the following excellent tutorial on LP solving:

-  [Thomas S. Ferguson.](#)  
Linear Programming – A Concise Introduction.  
UCLA, unpublished document available online.